

Tables

Wolfram Kahl

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1 Utilities

theory *Utils* = *Main*:

The following variant of modus ponens is useful as erule, for example for preparing local inductions.

lemma *mp1*: $\llbracket P; P \longrightarrow Q \rrbracket \Longrightarrow Q$
by *simp*

1.1 Functions

1.1.1 Combinators

consts *const* :: 'a \Rightarrow 'b \Rightarrow 'a
defs *const-def*: *const* == $\% x y . x$

lemma *const[simp]*: *const* *x y* = *x*
by (*simp add: const-def*)

lemma *const-expand*: *const* *c* = ($\% x . c$)
by (*rule ext, simp*)

consts *flip* :: ('a \Rightarrow 'b \Rightarrow 'c) \Rightarrow ('b \Rightarrow 'a \Rightarrow 'c)
defs *flip-def[simp]*: *flip* *f x y* == *f y x*

1.1.2 Currying

constdefs

curry :: (('a * 'b) \Rightarrow 'c) \Rightarrow 'a \Rightarrow 'b \Rightarrow 'c
curry *f* == $\% a b . f (a,b)$
uncurry :: ('a \Rightarrow 'b \Rightarrow 'c) \Rightarrow ('a * 'b) \Rightarrow 'c
uncurry *g* == $\% p . g (fst p) (snd p)$

lemma *curry[simp]*: *curry* *f a b* = *f (a,b)*
by (*simp add: curry-def*)

lemma *uncurry[simp]*: *uncurry* *g (a,b)* = *g a b*
by (*simp add: uncurry-def*)

lemma *uncurry-lambda2[simp]*: *uncurry* ($\lambda h t . F h t$) = ($\lambda p . F (fst p) (snd p)$)
by (*simp add: uncurry-def*)

lemma *uncurry-K2[simp]*: *uncurry* ($\lambda h t . F h$) = ($\lambda p . F (fst p)$)
by *simp*

lemma *uncurry-K[simp]*: *uncurry* ($\lambda h . F$) = ($\lambda p . F (snd p)$)
by *simp*

1.1.3 Pair Manipulation

constdefs

```
pupd1 :: ('a ⇒ 'b) ⇒ ('a * 'c) ⇒ ('b * 'c)
pupd1 f p == (f (fst p), snd p)
pupd2 :: ('a ⇒ 'b) ⇒ ('c * 'a) ⇒ ('c * 'b)
pupd2 g p == (fst p, g (snd p))
```

lemma *pupd1[simp]*: $pupd1\ f\ (x,y) = (f\ x,\ y)$
by (*simp add: pupd1-def*)

lemma *pupd2[simp]*: $pupd2\ g\ (x,y) = (x,\ g\ y)$
by (*simp add: pupd2-def*)

1.1.4 Extensionality

I had some problems applying *HOL.ext* directly.

lemma *f-ext*:
assumes *eq*: $\bigwedge x . f\ x = g\ x$
shows $f = g$
by (*rule HOL.ext*)

1.2 Function Properties

1.2.1 Associativity

constdefs

```
assoc :: ('a ⇒ 'a ⇒ 'a) ⇒ bool
assoc f == (ALL x y z . f (f x y) z = f x (f y z))
```

lemma *assoc[simp]*:
 $assoc\ f \implies f\ (f\ x\ y)\ z = f\ x\ (f\ y\ z)$
by (*unfold assoc-def, auto*)

lemma *assoc-intro[intro?]*:
 $\llbracket \bigwedge x\ y\ z . f\ (f\ x\ y)\ z = f\ x\ (f\ y\ z) \rrbracket \implies assoc\ f$
by (*unfold assoc-def, auto*)

lemma *const-assoc[simp]*: $assoc\ const$
by (*rule assoc-intro, simp*)

1.2.2 Idempotence

constdefs

```
idempotent :: ('a ⇒ 'a ⇒ 'a) ⇒ bool
idempotent f == (ALL x . f x x = x)
```

lemma *idempotent[simp]*:

idempotent f $\implies f x x = x$
by (*unfold idempotent-def*, *auto*)

lemma *idempotent-intro*[*intro?*]:
[[$\bigwedge x . f x x = x$]] \implies *idempotent f*
by (*unfold idempotent-def*, *auto*)

1.2.3 Commutativity

constdefs

commutative :: ('a \Rightarrow 'a \Rightarrow 'a) \Rightarrow *bool*
commutative f == (ALL x y . f x y = f y x)

lemma *commutative*:
commutative f $\implies f x y = f y x$
by (*unfold commutative-def*, *auto*)

lemma *commutative-intro*[*intro?*]:
[[$\bigwedge x y . f x y = f y x$]] \implies *commutative f*
by (*unfold commutative-def*, *auto*)

1.2.4 Units

consts *LRunit* :: ('a \Rightarrow 'a \Rightarrow 'a) \Rightarrow 'a \Rightarrow 'bool

defs *LRunit-def*: *LRunit f u* == ALL y . f u y = y & f y u = y

lemma *LRunit-left*: *LRunit f u* $\implies f u x = x$
by (*unfold LRunit-def*, *simp*)

lemma *LRunit-right*: *LRunit f u* $\implies f x u = x$
by (*unfold LRunit-def*, *simp*)

lemma *LRunit-equal*:

assumes *u*[*intro,simp*]: *LRunit f u*
assumes *v*[*intro,simp*]: *LRunit f v*
shows *u = v*

proof –

have *u = f u v* **by** (*rule sym*, *rule LRunit-right*, *simp*)

also have *f u v = v* **by** (*rule LRunit-left*, *simp*)

finally show *?thesis* .

qed

lemma *LRunit-const*:

ALL x . *LRunit c* (*F x*) $\implies F x = F$ *arbitrary*
apply (*subgoal-tac* ALL x y . *F x = F y*, *simp*)
apply (*intro strip*)
apply (*rule LRunit-equal* [*of c*])
apply *simp-all*
done

1.3 Additional Material for Lists

lemma *foldr-append*:

$foldr\ f\ (xs\ @\ ys)\ z = foldr\ f\ xs\ (foldr\ f\ ys\ z)$

by (*induct-tac xs, auto*)

consts

$foldr-1 :: ('a \Rightarrow 'a \Rightarrow 'a) * 'a\ list \Rightarrow 'a$

recdef *foldr-1 measure* (% (f,xs) . length xs)

foldr-1-sing[simp]: $foldr-1\ (f,\ x\ \#\ []) = x$

foldr-1-cons0[simp]: $foldr-1\ (f,\ x\ \#\ y\ \#\ ys) = f\ x\ (foldr-1\ (f,\ y\ \#\ ys))$

lemma *foldr-1-cons*[simp]:

$\llbracket xs \sim = [] \rrbracket \Longrightarrow foldr-1\ (f,\ x\ \#\ xs) = f\ x\ (foldr-1\ (f,\ xs))$

by (*cases xs, auto*)

lemma *foldr-1-append-distr-0*:

assoc f \Longrightarrow

$(ALL\ x\ ys\ .\ ys\ \sim = [] \longrightarrow foldr-1\ (f,\ x\ \#\ xs\ @\ ys) = f\ (foldr-1\ (f,\ x\ \#\ xs))\ (foldr-1\ (f,\ ys)))$

by (*induct-tac xs, auto*)

lemma *foldr-1-append-distr*:

assumes *xs*[simp]: $xs\ \sim = []$

assumes *ys*[simp]: $ys\ \sim = []$

assumes *f*[simp]: *assoc f*

shows $foldr-1\ (f,\ xs\ @\ ys) = f\ (foldr-1\ (f,\ xs))\ (foldr-1\ (f,\ ys))$

proof (*cases xs*)

case *Nil*

thus *?thesis* **by** *simp*

next

case *Cons*

thus *?thesis* **by** (*insert foldr-1-append-distr-0 [of f tl xs], simp*)

qed

lemma *mem-append*[simp]:

$x\ mem\ (xs\ @\ ys) = (x\ mem\ xs\ | x\ mem\ ys)$

by (*induct-tac xs, simp-all*)

constdefs

zipWith :: $('a \Rightarrow 'b \Rightarrow 'c) \Rightarrow 'a\ list \Rightarrow 'b\ list \Rightarrow 'c\ list$

$zipWith\ f\ xs\ ys == map\ (uncurry\ f)\ (zip\ xs\ ys)$

lemma *zipWith-Cons*[simp]:

$zipWith\ f\ (x\ \#\ xs)\ (y\ \#\ ys) = f\ x\ y\ \#\ zipWith\ f\ xs\ ys$

by (*simp add: zipWith-def*)

lemma *zipWith-Nil-1*[simp]:

$zipWith\ f\ []\ ys = []$

by (*simp add: zipWith-def*)

lemma *zipWith-Nil-2*[*simp*]:

zipWith f xs [] = []

by (*simp add: zipWith-def*)

lemma *zipWith-assoc-0*:

assoc f \implies ALL ys zs . zipWith f (zipWith f xs ys) zs = zipWith f xs (zipWith f ys zs)

apply (*induct-tac xs, simp*)

apply (*intro strip*)

apply (*induct-tac ys, simp*)

apply (*induct-tac zs, simp*)

apply *simp*

done

lemma *zipWith-assoc*[*simp*]:

assoc f \implies zipWith f (zipWith f xs ys) zs = zipWith f xs (zipWith f ys zs)

by (*insert zipWith-assoc-0, auto*)

lemma *assoc-zipWith*[*simp*]: *assoc f \implies assoc (zipWith f)*

by (*rule assoc-intro, simp*)

1.4 Option Utilities

consts

option :: 'b \Rightarrow ('a \Rightarrow 'b) \Rightarrow 'a option \Rightarrow 'b

primrec

option-N: option r f None = r

option-S: option r f (Some a) = f a

1.4.1 The Option Monad

consts

optThen :: 'a option \Rightarrow ('a \Rightarrow 'b option) \Rightarrow 'b option

primrec

optThen-N: optThen None = (% f . None)

optThen-S: optThen (Some x) = (% f . f x)

declare *optThen-N*[*simp del*]

declare *optThen-S*[*simp del*]

lemma *optThen-N-f*[*simp*]: *optThen None f = None*

by (*simp add: optThen-N*)

lemma *optThen-S-f*[*simp*]: *optThen (Some x) f = f x*

by (*simp add: optThen-S*)

lemma *optThen-result-Some*[simp]:
 $optThen\ m\ f = Some\ x \implies \exists y . m = Some\ y \ \& \ f\ y = Some\ x$
by (*cases m, simp-all*)

lemma *optThen-option-map*[simp]:
 $optThen\ (option\ map\ f\ x)\ g = optThen\ x\ (g\ o\ f)$
by (*cases x, simp-all*)

lemma *optThen-assoc*[simp]:
 $optThen\ (optThen\ x\ f)\ g = optThen\ x\ (\% x . optThen\ (f\ x)\ g)$
by (*cases x, simp-all*)

1.4.2 Equality Option Collapse

consts
 $optEq :: 'a\ option \Rightarrow 'a\ option \Rightarrow 'a\ option$

primrec
 $optEq\ N: optEq\ None = (\% my . None)$
 $optEq\ S: optEq\ (Some\ x) = option\ None\ (\% y . if\ (x=y)\ then\ Some\ x\ else\ None)$

declare *optEq-N*[simp del]
declare *optEq-S*[simp del]

lemma *optEq-S-S*[simp]:
 $optEq\ (Some\ x)\ (Some\ y) = (if\ (x=y)\ then\ Some\ x\ else\ None)$
by (*simp add: optEq-S*)

lemma *optEq-S-N*[simp]: $optEq\ (Some\ x)\ None = None$
by (*simp add: optEq-S*)

lemma *optEq-N-S*[simp]: $optEq\ None\ (Some\ x) = None$
by (*simp add: optEq-N*)

lemma *optEq-N-N*[simp]: $optEq\ None\ None = None$
by (*simp add: optEq-N*)

lemma *optEq-assoc*[simp]: $optEq\ (optEq\ a\ b)\ c = optEq\ a\ (optEq\ b\ c)$
apply (*cases a, simp add: optEq-N*)
apply (*cases b, simp add: optEq-N*)
apply (*cases c, simp add: optEq-N*)
apply *simp*
done

lemma *assoc-optEq*[simp]: $assoc\ optEq$
by (*rule assoc-intro, simp*)

lemma *optEq-idempotent*[simp]: $idempotent\ optEq$
by (*rule idempotent-intro, case-tac x, simp-all*)


```

lemma optEq-commutative[simp]: commutative optEq
apply (rule commutative-intro)
apply (case-tac x, case-tac y, simp-all add: optEq-N optEq-S)
apply (case-tac y, simp-all add: optEq-N optEq-S)
done

end

```

2 Non-Empty Lists

theory *NEList = Utils:*

2.1 Datatype Definition

```

typedef 'a neList = { xs :: 'a list . xs  $\sim$  [] }
by auto

```

```

lemma Rep-neList-non-Nil[simp]: Rep-neList x  $\sim$  []
by (insert Rep-neList [of x], simp add: neList-def)

```

```

lemma neList-append[simp,intro?]:
  [| xs : neList; ys : neList |]  $\implies$  (xs @ ys) : neList
by (simp add: neList-def)

```

```

lemma neList-map[simp,intro?]:
  [| xs : neList |]  $\implies$  (map f xs) : neList
by (simp add: neList-def)

```

```

lemma neList-zipWith[simp,intro?]:
  [| xs : neList; ys : neList |]  $\implies$  (zipWith f xs ys) : neList
apply (auto simp add: neList-def)
apply (cases xs, simp)
apply (cases ys, simp-all)
done

```

2.2 Construction

```

constdefs
  singleton :: 'a  $\Rightarrow$  'a neList
  singleton x == Abs-neList [x]
  append :: 'a neList  $\Rightarrow$  'a neList  $\Rightarrow$  'a neList
  append xs ys == Abs-neList (Rep-neList xs @ Rep-neList ys)
  cons1 :: 'a  $\Rightarrow$  'a neList  $\Rightarrow$  'a neList
  cons1 x xs == Abs-neList (x # Rep-neList xs)

```

```

lemma singleton-inj: singleton x = singleton y  $\implies$  x = y
apply (simp add: singleton-def Abs-neList-inverse)
apply (subgoal-tac [x] = [y])
  prefer 2
  apply (subst Abs-neList-inject [THEN sym])
  apply (simp add: neList-def)
  apply (simp add: neList-def)
  apply assumption
apply auto
done

```

```

lemma cons1:
  cons1 x xs = append (singleton x) xs
apply (unfold cons1-def)
apply (unfold append-def)
apply (unfold singleton-def)
apply (simp-all add: Abs-neList-inverse Rep-neList-inverse neList-def)
done

```

```

lemma cons1-inj: cons1 x xs = cons1 y ys  $\implies$  x = y & xs = ys
apply (simp add: cons1-def Abs-neList-inverse)
apply (subgoal-tac x # Rep-neList xs = y # Rep-neList ys)
  prefer 2
  apply (subst Abs-neList-inject [THEN sym])
  apply (simp add: neList-def)
  apply (simp add: neList-def)
  apply assumption
apply auto
apply (subst Rep-neList-inject [THEN sym])
apply assumption
done

```

```

lemma cons1-neq-tl[simp]: cons1 x xs  $\sim$  xs
apply (simp add: cons1-def singleton-def Abs-neList-inverse)
apply (cases Rep-neList xs)
apply (auto simp add: neList-def Abs-neList-inverse)
done

```

```

lemma cons1-neq-singleton[simp]: cons1 x xs  $\sim$  singleton y
apply (simp add: cons1-def singleton-def Abs-neList-inverse)
apply (subst Abs-neList-inject)
apply (auto simp add: neList-def)
done

```

```

lemma singleton-neq-cons1[simp]: singleton y  $\sim$  cons1 x xs
apply (simp add: cons1-def singleton-def Abs-neList-inverse)
apply (subst Abs-neList-inject)
apply (auto simp add: neList-def)
apply (cases Rep-neList xs)

```

apply (*auto simp add: neList-def*)
done

lemma *append-neq-singleton*[*simp*]: *append xs ys* \sim *singleton x*
apply (*simp add: append-def singleton-def Abs-neList-inverse*)
apply (*subst Abs-neList-inject*)
apply (*auto simp add: neList-def*)
apply (*cases xs*)
apply (*cases ys*)
apply (*simp add: neList-def Abs-neList-inverse*)
apply (*case-tac ya, simp*)
apply (*case-tac y, simp-all*)
done

lemma *singleton-neq-append*[*simp*]: *singleton x* \sim *append xs ys*
apply (*simp add: append-def singleton-def Abs-neList-inverse*)
apply (*subst Abs-neList-inject*)
apply (*auto simp add: neList-def*)
apply (*cases xs*)
apply (*cases ys*)
apply (*simp add: neList-def Abs-neList-inverse*)
apply (*case-tac ya, simp*)
apply (*case-tac y, simp-all*)
done

lemma *append-inj2-0*: (*append xs ys = append xs zs*) = (*ys = zs*)
apply (*simp add: append-def Abs-neList-inverse*)
apply (*subst Abs-neList-inject*)
apply (*auto simp add: neList-def Rep-neList-inject [THEN iffD1]*)
done

lemma *append-inj2*: *append xs ys = append xs zs* \implies *ys = zs*
by (*rule append-inj2-0 [THEN iffD1]*)

lemma *append-inj2-neq*: *append xs ys* \sim *append xs zs* \implies *ys* \sim *zs*
by (*insert append-inj2-0, auto*)

lemma *append-inj2-conv*: *ys = zs* \implies *append xs ys = append xs zs*
by (*insert append-inj2-0, auto*)

lemma *append-inj2-neq-conv*: *ys* \sim *zs* \implies *append xs ys* \sim *append xs zs*
by (*insert append-inj2-0, auto*)

lemma *append-inj1-0*: (*append xs zs = append ys zs*) = (*xs = ys*)
apply (*simp add: append-def Abs-neList-inverse*)
apply (*subst Abs-neList-inject*)
apply (*auto simp add: neList-def Rep-neList-inject [THEN iffD1]*)
done

lemma *append-inj1*: $append\ xs\ zs = append\ ys\ zs \implies xs = ys$
by (*rule append-inj1-0 [THEN iffD1]*)

lemma *append-inj1-neq*: $append\ xs\ zs \sim = append\ ys\ zs \implies xs \sim = ys$
by (*insert append-inj1-0, auto*)

lemma *append-assoc[simp]*:
 $append\ (append\ xs\ ys)\ zs = append\ xs\ (append\ ys\ zs)$
apply (*unfold append-def*)
apply (*cases xs*)
apply (*cases ys*)
apply (*cases zs*)
apply (*subgoal-tac y @ ya : neList*)
apply (*subgoal-tac ya @ yb : neList*)
apply (*simp add: Abs-neList-inverse*)
apply (*unfold neList-def, simp-all*)
done

lemma *assoc-append[simp]*: *assoc append*
by (*rule assoc-intro, simp*)

2.3 foldr1

constdefs
 $foldr1 :: ('a \Rightarrow 'a \Rightarrow 'a) \Rightarrow 'a\ neList \Rightarrow 'a$
 $foldr1\ f\ xs == foldr1\ (f,\ Rep-neList\ xs)$

lemma *foldr1-singleton[simp]*:
 $foldr1\ f\ (singleton\ x) = x$
apply (*unfold singleton-def*)
apply (*unfold foldr1-def*)
apply (*simp-all add: Abs-neList-inverse Rep-neList-inverse neList-def*)
done

lemma *foldr1-cons1[simp]*:
 $foldr1\ f\ (cons1\ x\ xs) = f\ x\ (foldr1\ f\ xs)$
apply (*unfold cons1-def*)
apply (*unfold foldr1-def*)
apply (*simp-all add: Abs-neList-inverse Rep-neList-inverse neList-def*)
done

lemma *foldr1-append-distr[simp]*:
 $assoc\ f \implies foldr1\ f\ (append\ xs\ ys) = f\ (foldr1\ f\ xs)\ (foldr1\ f\ ys)$
apply (*cases xs*)
apply (*cases ys*)
apply (*unfold append-def*)
apply (*unfold foldr1-def*)
apply (*cases Rep-neList xs*)
prefer 2

```

apply (cases Rep-neList ys)
apply (simp-all add: Abs-neList-inverse Rep-neList-inverse neList-def
        del: foldr-1-cons)
apply (insert foldr-1-append-distr [of Rep-neList xs Rep-neList ys f, THEN sym])
apply (simp-all add: Abs-neList-inverse Rep-neList-inverse neList-def
        del: foldr-1-cons)
done

```

2.4 map

constdefs

```

  map1 :: ('a ⇒ 'b) ⇒ 'a neList ⇒ 'b neList
  map1 f xs == Abs-neList (map f (Rep-neList xs))

```

lemma map1-singleton[simp]:

```

  map1 f (singleton x) = singleton (f x)
apply (unfold singleton-def)
apply (unfold map1-def)
apply (simp-all add: Abs-neList-inverse Rep-neList-inverse neList-def)
done

```

lemma map1-cons1[simp]:

```

  map1 f (cons1 x xs) = cons1 (f x) (map1 f xs)
apply (unfold cons1-def)
apply (unfold map1-def)
apply (simp-all add: Abs-neList-inverse Rep-neList-inverse neList-def)
done

```

lemma map1-append-distr[simp]:

```

  map1 f (append xs ys) = append (map1 f xs) (map1 f ys)
apply (cases xs)
apply (cases ys)
apply (unfold append-def)
apply (unfold map1-def)
apply (cases Rep-neList xs)
prefer 2
apply (cases Rep-neList ys)
apply (simp-all add: Abs-neList-inverse Rep-neList-inverse neList-def
        del: foldr-1-cons)
done

```

lemma map1-contract-comp:

```

  map1 f (map1 g xs) = map1 (f o g) xs
apply (unfold map1-def)
apply (cases xs)
apply (cases Rep-neList xs)
apply (simp-all add: Abs-neList-inverse Rep-neList-inverse neList-def map-compose
        del: foldr-1-cons)
done

```

lemma *map1-comp*:
 $(\text{map1 } f) \circ (\text{map1 } g) = \text{map1 } (f \circ g)$
by (*simp add: comp-def map1-contract-comp*)

2.5 Induction

lemma *neList-patterns*:
 $(\exists x . xs = \text{singleton } x) \mid (\exists y \ ys . xs = \text{append } (\text{singleton } y) \ ys)$
apply (*cases xs*)
apply (*cases Rep-neList xs*)
apply (*simp-all add: Abs-neList-inverse Rep-neList-inverse neList-def*
del: foldr-1-cons)
apply (*cases tl (Rep-neList xs)*)
apply (*unfold singleton-def*)
apply (*rule disjI1*)
apply (*rule exI*)
apply (*simp add: Abs-neList-inverse Rep-neList-inverse neList-def*
del: foldr-1-cons)
apply (*rule disjI2*)
apply (*rule-tac x=a in exI*)
apply (*rule-tac x=Abs-neList list in exI*)
apply (*unfold append-def*)
apply (*simp add: Abs-neList-inverse Rep-neList-inverse neList-def*
del: foldr-1-cons)
done

lemma *neList-list-induct*:
 $(\forall x . P (\text{singleton } x)) \longrightarrow$
 $(\forall x \ y \ ys . P (\text{Abs-neList } (y \# \ ys))) \longrightarrow P (\text{Abs-neList } (x \# \ y \# \ ys)) \longrightarrow$
 $(\forall z . P (\text{Abs-neList } (z \# \ zs)))$
apply (*induct-tac zs*)
apply (*simp add: Abs-neList-inverse Rep-neList-inverse neList-def singleton-def*
del: foldr-1-cons)
apply (*intro strip*)
apply *simp*
done

lemma *neList-cons-induct-0*:
 $\llbracket \bigwedge x . P (\text{singleton } x); \bigwedge x \ xs . P \ xs \implies P (\text{cons1 } x \ xs) \rrbracket$
 $\implies P \ zs$
apply (*insert neList-patterns [of zs]*)
apply (*simp only: cons1*)
apply (*erule disjE*)
apply (*erule exE*)
apply *simp*
apply (*erule exE*)
apply (*erule exE*)
apply (*insert neList-list-induct [of P tl (Rep-neList zs)]*)

```

apply simp
apply (subgoal-tac  $\forall x y ys. P (Abs-neList (y \# ys)) \longrightarrow P (Abs-neList (x \# y \# ys))$ )
  prefer 2
  apply (intro strip)
  apply (subgoal-tac  $P (singleton\ x)$ )
    prefer 2
    apply simp
    apply (subgoal-tac  $P (append (singleton\ x) (Abs-neList (ya \# ysa)))$ )
      prefer 2
      apply simp
      apply (subgoal-tac  $Abs-neList (x \# ya \# ysa) = append (singleton\ x) (Abs-neList (ya \# ysa))$ )
        prefer 2
        apply (simp (no-asm-simp) add: Abs-neList-inverse Rep-neList-inverse neList-def singleton-def append-def del: foldr-1-cons)
          apply simp
          apply simp
          apply (drule-tac  $x=y$  in spec)
          apply (simp add: Abs-neList-inverse Rep-neList-inverse neList-def singleton-def append-def del: foldr-1-cons)
            done

lemma neList-cons-induct:
   $\bigwedge zs. \llbracket \bigwedge x. P (singleton\ x); \bigwedge x xs. P\ xs \implies P (cons1\ x\ xs) \rrbracket \implies P\ zs$ 
by (rule neList-cons-induct-0, auto)

lemma neList-cons-inductA:
   $\llbracket \bigwedge x. P (singleton\ x); \bigwedge x xs. P\ xs \implies P (cons1\ x\ xs) \rrbracket \implies ALL\ zs. P\ zs$ 
by (intro strip, rule neList-cons-induct, simp-all)

lemma neList-append-induct-0:
   $\llbracket \bigwedge x. P (singleton\ x); \bigwedge xs ys. P\ xs \implies P\ ys \implies P (append\ xs\ ys) \rrbracket \implies P\ zs$ 
apply (insert neList-patterns [of zs])
apply (erule disjE)
apply (erule exE)
apply simp
apply (erule exE)
apply (erule exE)
apply (insert neList-list-induct [of P tl (Rep-neList zs)])
apply simp
apply (subgoal-tac  $\forall x y ys. P (Abs-neList (y \# ys)) \longrightarrow P (Abs-neList (x \# y \# ys))$ )
  prefer 2

```

```

apply (intro strip)
apply (subgoal-tac P (singleton x))
  prefer 2
  apply simp
apply (subgoal-tac P (append (singleton x) (Abs-neList (ya # ysa))))
  prefer 2
  apply simp
apply (subgoal-tac Abs-neList (x # ya # ysa) = append (singleton x) (Abs-neList
(ya # ysa))
  prefer 2
  apply (simp (no-asm-simp) add: Abs-neList-inverse Rep-neList-inverse neList-def
singleton-def append-def
    del: foldr-1-cons)
  apply simp
apply simp
apply (drule-tac x=y in spec)
apply (simp add: Abs-neList-inverse Rep-neList-inverse neList-def singleton-def
append-def
  del: foldr-1-cons)
done

```

```

lemma neList-append-induct:
 $\bigwedge zs . \llbracket \bigwedge x . P (\text{singleton } x); \bigwedge xs \ ys . P \ xs \implies P \ ys \implies P (\text{append } xs \ ys) \rrbracket$ 
 $\implies P \ zs$ 
by (rule neList-append-induct-0, auto)

```

2.6 Elements

```

constdefs
  elem :: 'a  $\Rightarrow$  'a neList  $\Rightarrow$  bool
  elem x xs == x mem (Rep-neList xs)

```

```

lemma elem-singleton[simp]:
  elem x (singleton y) = (x = y)
apply (simp add: Abs-neList-inverse Rep-neList-inverse neList-def singleton-def
elem-def
  del: foldr-1-cons)
apply fast
done

```

```

lemma elem-append[simp]:
  elem x (append xs ys) = (elem x xs | elem x ys)
by (simp add: Abs-neList-inverse Rep-neList-inverse neList-def elem-def append-def
  del: foldr-1-cons)

```

```

constdefs
  set1 :: 'a neList  $\Rightarrow$  'a set
  set1 xs == set (Rep-neList xs)

```


lemma *set1-singleton*[simp]:
 $set1 (singleton x) = \{ x \}$
by (simp add: Abs-neList-inverse Rep-neList-inverse neList-def singleton-def set1-def
del: foldr-1-cons)

lemma *set1-append*[simp]:
 $set1 (append xs ys) = set1 xs \cup set1 ys$
apply (subgoal-tac (Rep-neList xs @ Rep-neList ys) : neList)
apply (simp add: Abs-neList-inverse Rep-neList-inverse neList-def append-def set1-def
del: foldr-1-cons)
apply (cases xs)
apply (simp add: Abs-neList-inverse Rep-neList-inverse neList-def elem-def append-def
del: foldr-1-cons)
done

2.7 neFold

constdefs
 $neFold :: ('a \Rightarrow 'b) \Rightarrow ('b \Rightarrow 'b \Rightarrow 'b) \Rightarrow 'a\ neList \Rightarrow 'b$
 $neFold\ s\ c\ xs == foldr1\ c\ (map1\ s\ xs)$

lemma *neFold-singleton*[simp]:
 $neFold\ s\ c\ (singleton\ x) = s\ x$
by (simp add: neFold-def)

lemma *neFold-cons1*[simp]:
 $neFold\ s\ c\ (cons1\ x\ xs) = c\ (s\ x)\ (neFold\ s\ c\ xs)$
by (simp add: neFold-def)

lemma *neFold-append*[simp]:
 $assoc\ c \Longrightarrow neFold\ s\ c\ (append\ xs\ ys) = c\ (neFold\ s\ c\ xs)\ (neFold\ s\ c\ ys)$
by (simp add: neFold-def)

lemma *neFold-const-idempotent*[simp]:
 $\llbracket assoc\ c; idempotent\ c \rrbracket \Longrightarrow neFold\ (\% h . a)\ c\ t = a$
by (induct-tac t rule: neList-append-induct, simp-all)

2.8 ZipWith

constdefs
 $zipWith1 :: ('a \Rightarrow 'b \Rightarrow 'c) \Rightarrow 'a\ neList \Rightarrow 'b\ neList \Rightarrow 'c\ neList$
 $zipWith1\ f\ xs\ ys == Abs-neList\ (zipWith\ f\ (Rep-neList\ xs)\ (Rep-neList\ ys))$

lemma *zipWith1-S-S*[simp]:
 $zipWith1\ f\ (singleton\ x)\ (singleton\ y) = singleton\ (f\ x\ y)$
by (simp add: Abs-neList-inverse Rep-neList-inverse neList-def singleton-def zipWith1-def)

lemma *zipWith1-S-C*[simp]:
 $zipWith1\ f\ (singleton\ x)\ (cons1\ y\ ys) = singleton\ (f\ x\ y)$

by (*simp add: Abs-neList-inverse Rep-neList-inverse neList-def singleton-def cons1-def append-def zipWith1-def*)

lemma *zipWith1-C-S*[*simp*]:

zipWith1 f (cons1 x xs) (singleton y) = singleton (f x y)

by (*simp add: Abs-neList-inverse Rep-neList-inverse neList-def singleton-def cons1-def append-def zipWith1-def*)

lemma *zipWith1-C-C*[*simp*]:

zipWith1 f (cons1 x xs) (cons1 y ys) = cons1 (f x y) (zipWith1 f xs ys)

apply (*subgoal-tac zipWith f (Rep-neList xs) (Rep-neList ys) : neList*)

apply (*simp add: Abs-neList-inverse Rep-neList-inverse neList-def singleton-def cons1-def append-def zipWith1-def*)

apply (*cases xs, simp*)

apply (*cases ys, simp*)

apply (*simp-all add: Abs-neList-inverse*)

done

lemma *zipWith1-assoc*[*simp*]:

assoc f \implies zipWith1 f (zipWith1 f xs ys) zs = zipWith1 f xs (zipWith1 f ys zs)

apply (*simp add: zipWith1-def*)

apply (*cases xs, simp*)

apply (*cases ys, simp*)

apply (*cases zs, simp*)

apply (*simp-all add: Abs-neList-inverse*)

done

lemma *assoc-zipWith1*[*simp*]: *assoc f \implies assoc (zipWith1 f)*

by (*rule assoc-intro, simp*)

Attempting a direct proof of this associativity, not using associativity of zipWith:

lemma *zipWith1-assoc-2*[*simp*]:

assoc f \implies

ALL xs ys zs . zipWith1 f (zipWith1 f xs ys) zs = zipWith1 f xs (zipWith1 f ys zs)

apply (*rule neList-cons-inductA*)

apply (*rule neList-cons-inductA*)

apply *simp-all*

apply (*rule neList-cons-inductA*)

apply *simp-all*

apply (*rule neList-cons-inductA*)

apply *simp-all*

done

lemma *zipWith1-assoc-1*[*simp*]:

assoc f \implies

ALL ys zs . zipWith1 f (zipWith1 f xs ys) zs = zipWith1 f xs (zipWith1 f ys zs)

apply (*induct-tac xs rule: neList-cons-induct*)

```

apply simp
apply (rule allI)
apply (induct-tac ys rule: neList-cons-induct)
apply (simp add: cons1)
apply (rule allI)
apply (induct-tac zs rule: neList-cons-induct)
apply (simp-all add: cons1)
done

```

```

lemma zipWith1-assoc-3[simp]:
  assoc f  $\implies$ 
  zipWith1 f (zipWith1 f xs ys) zs = zipWith1 f xs (zipWith1 f ys zs)
apply (induct-tac xs rule: neList-cons-induct)
apply simp
apply (induct-tac ys rule: neList-cons-induct)
apply (simp add: cons1)
apply (induct-tac zs rule: neList-cons-induct)
apply (simp-all add: cons1)
done

```

2.9 Other Properties

```

lemma foldr1-map1-LRunit[simp]:
  [  $\bigwedge x . LRunit f (u x); assoc f$  ]  $\implies$ 
  foldr1 f (map1 u xs) = u arbitrary
apply (induct-tac xs rule: neList-append-induct, simp-all)
apply (rule LRunit-equal [of f], simp-all add: LRunit-left LRunit-right)
done

```

end

3 The Table Datatype Definition

theory Table = NEList:

3.1 Datatype and Primitive Operations

datatype ('h,'t) T = T ('h * 't) neList

consts

unT :: ('h,'t) T \Rightarrow ('h * 't) neList

primrec

unT-def[simp]: *unT* (T ps) = ps

constdefs

addH :: 'h \Rightarrow 't \Rightarrow ('h,'t) T
addH h t == T (singleton (h,t))

constdefs

$hConc :: ('h, 't) T \Rightarrow ('h, 't) T \Rightarrow ('h, 't) T$
 $hConc\ t1\ t2 == T\ (append\ (unT\ t1)\ (unT\ t2))$

lemma $hConc\text{-}assoc[simp]$: $hConc\ (hConc\ t1\ t2)\ t3 = hConc\ t1\ (hConc\ t2\ t3)$
by ($unfold\ hConc\text{-}def, simp$)

lemma $assoc\text{-}hConc[simp]$: $assoc\ hConc$
by ($rule\ assoc\text{-}intro, simp$)

constdefs

$cell :: 'c \Rightarrow ('c, unit) T$
 $cell\ c == addH\ c\ ()$

3.2 Folding and Induction via hConc

constdefs

$tFold :: ('h \Rightarrow 't \Rightarrow 'r) \Rightarrow ('r \Rightarrow 'r \Rightarrow 'r) \Rightarrow ('h, 't) T \Rightarrow 'r$
 $tFold\ h\ c\ t == neFold\ (uncurry\ h)\ c\ (unT\ t)$

lemma $tFold\text{-}addH[simp]$:
 $tFold\ ah\ c\ (addH\ h\ t) = ah\ h\ t$
by ($simp\ add: tFold\text{-}def\ addH\text{-}def$)

lemma $tFold\text{-}hConc[simp]$:
 $assoc\ c \Longrightarrow tFold\ h\ c\ (hConc\ t1\ t2) = c\ (tFold\ h\ c\ t1)\ (tFold\ h\ c\ t2)$
by ($simp\ add: tFold\text{-}def\ hConc\text{-}def$)

lemma $T\text{-}induct\text{-}0$:

$\llbracket \bigwedge h\ t0 . P\ (addH\ h\ t0); \bigwedge t1\ t2 . P\ t1 \Longrightarrow P\ t2 \Longrightarrow P\ (hConc\ t1\ t2) \rrbracket$
 $\Longrightarrow P\ t$

apply ($cases\ t, simp$)

apply ($rule\ \text{-}tac\ zs=neList\ \mathbf{in}\ neList\text{-}append\text{-}induct$)

apply ($subst\ surjective\text{-}pairing$)

apply ($simp\ only: addH\text{-}def$)

apply ($simp\ add: hConc\text{-}def$)

apply ($subgoal\ \text{-}tac\ append\ xs\ ys = append\ (unT\ (T\ xs))\ (unT\ (T\ ys))$)

apply ($simp\ (no\ \text{-}asm\text{-}simp)\ del: unT\text{-}def$)

apply $simp$

done

lemma $T\text{-}induct$:

$\bigwedge t . \llbracket \bigwedge h\ t0 . P\ (addH\ h\ t0); \bigwedge t1\ t2 . P\ t1 \Longrightarrow P\ t2 \Longrightarrow P\ (hConc\ t1\ t2) \rrbracket$
 $\Longrightarrow P\ t$

by ($rule\ T\text{-}induct\text{-}0, auto$)

3.3 hCons

constdefs

$hCons :: ('h * 't) \Rightarrow ('h, 't) T \Rightarrow ('h, 't) T$
 $hCons\ p\ t1 == T\ (cons1\ p\ (unT\ t1))$

lemma *hCons*:

$hCons\ p\ t1 = hConc\ (uncurry\ addH\ p)\ t1$

proof –

have $hCons\ p\ t1 = T\ (cons1\ p\ (unT\ t1))$ **by** (*simp add: hCons-def*)
also have $\dots = T\ (append\ (singleton\ p)\ (unT\ t1))$ **by** (*simp add: cons1*)
also have $\dots = T\ (append\ (unT\ (T\ (singleton\ (fst\ p,\ snd\ p))))\ (unT\ t1))$ **by**
simp
also have $\dots = T\ (append\ (unT\ (addH\ (fst\ p)\ (snd\ p)))\ (unT\ t1))$ **by** (*simp*
add: addH-def)
also have $\dots = hConc\ (addH\ (fst\ p)\ (snd\ p))\ t1$ **by** (*simp add: hConc-def*)
also have $\dots = hConc\ (uncurry\ addH\ (fst\ p,\ snd\ p))\ t1$ **by** (*simp del: surjective-pairing*
[*THEN sym*])
also have $\dots = hConc\ (uncurry\ addH\ p)\ t1$ **by** *simp*
finally show *?thesis* .
qed

lemma *hCons-p*:

$hCons\ (h,t)\ t1 = hConc\ (addH\ h\ t)\ t1$

by (*simp add: hCons*)

lemma *tFold-hCons[simp]*:

$tFold\ ah\ c\ (hCons\ p\ t) = c\ (uncurry\ ah\ p)\ (tFold\ ah\ c\ t)$

by (*simp add: tFold-def hCons-def*)

lemma *T-cons-induct-0*:

$\llbracket \bigwedge h\ t0 . P\ (addH\ h\ t0); \bigwedge p\ t2 . P\ t2 \implies P\ (hCons\ p\ t2) \rrbracket$
 $\implies P\ t$

apply (*cases t, simp*)

apply (*rule-tac zs=neList in neList-cons-induct*)

apply (*subst surjective-pairing*)

apply (*simp only: addH-def*)

apply (*subst surjective-pairing*)

apply (*simp add: hConc-def hCons-def addH-def*)

apply (*subgoal-tac cons1 x xs = cons1 x (unT (T xs))*)

apply (*simp (no-asm-simp) del: unT-def*)

apply *simp*

done

lemma *T-cons-induct*:

$\bigwedge t . (\llbracket \bigwedge h\ t0 . P\ (addH\ h\ t0); \bigwedge p\ t2 . P\ t2 \implies P\ (hCons\ p\ t2) \rrbracket \implies P\ t)$

by (*rule T-cons-induct-0, auto*)

lemma *T-cons-inductA*:

$\llbracket \bigwedge h\ t0 . P\ (addH\ h\ t0); \bigwedge p\ t2 . P\ t2 \implies P\ (hCons\ p\ t2) \rrbracket \implies ALL\ t . P\ t$

by (*intro strip, rule T-cons-induct, simp-all*)

lemma *hConc-hCons[simp]*: $hConc (hCons p t1) t2 = hCons p (hConc t1 t2)$
by (*induct-tac t1 rule: T-cons-induct, simp-all add: hCons*)

end

4 Table Utilities and Properties

theory *Tables = Table:*

4.1 Additional Properties of *tFold*

lemma *tFold-const2-idempotent[simp]*:
 $\llbracket \text{assoc } c; \text{idempotent } c \rrbracket \implies tFold (\% h t . a) c t = a$
by (*induct-tac t rule: T-induct, simp-all*)

lemma *tFold-tFold-join*:
 $\llbracket \text{assoc } c2 \rrbracket \implies$
 $tFold a2 c2 (tFold (\lambda h0 t0. \text{addH } (HH h h0 t0) (TT h h0 t0)) hConc t) =$
 $tFold (\lambda h0 t0 . a2 (HH h h0 t0) (TT h h0 t0)) c2 t$
by (*induct-tac t rule: T-induct, simp-all*)

lemma *tFold-tFold-hConc*:
 $\llbracket \text{assoc } c2 \rrbracket \implies$
 $tFold a2 c2 (tFold a1 hConc t) =$
 $tFold (\lambda h0 t0 . tFold a2 c2 (a1 h0 t0)) c2 t$
by (*induct-tac t rule: T-induct, simp-all*)

lemma *tFold-tFold-const*:
 $\llbracket \text{assoc } c2 \rrbracket \implies$
 $tFold a2 c2 (tFold a1 const t) =$
 $tFold (\lambda h0 t0 . tFold a2 c2 (a1 h0 t0)) const t$
apply (*induct-tac t rule: T-induct, simp-all del: const*)
apply (*subst const*)
apply (*simp del: const*)
done

lemma *f-tFold-const*:
 $f (tFold a const t) = tFold (\% h0 t0 . f (a h0 t0)) const t$
apply (*induct-tac t rule: T-induct, simp-all del: const*)
apply (*subst const*)
apply (*simp del: const*)
done

lemma *tFold-const-const[simp]*: $tFold (\% h t . r) const t = r$
by (*induct-tac t rule: T-induct, simp-all del: const, simp*)

lemma *tFold-foldr1-hConc*:

$assoc\ c \implies tFold\ a\ c\ (foldr1\ hConc\ ts) = foldr1\ c\ (map1\ (tFold\ a\ c)\ ts)$
by (*induct-tac ts rule: neList-append-induct, simp-all*)

lemma *tFold-LRunit*[*simp*]:

$\llbracket \bigwedge h\ t . LRunit\ c\ (a\ h\ t); assoc\ c \rrbracket \implies$
 $tFold\ a\ c\ t = a$ *arbitrary arbitrary*

apply (*induct-tac t rule: T-induct, simp-all*)

apply (*rule LRunit-equal [of c], simp-all add: LRunit-left LRunit-right*)

done

consts *tFold0* :: ($'c \Rightarrow 'c \Rightarrow 'c$) \Rightarrow ($'c, 'u$) $T \Rightarrow 'c$

defs *tFold0-def*: $tFold0 == tFold\ (\% h\ t . h)$

lemma *tFold0-addH*[*simp*]: $tFold0\ c\ (addH\ h\ t) = h$

by (*unfold tFold0-def, simp*)

lemma *tFold0-hConc*[*simp*]:

$assoc\ c \implies tFold0\ c\ (hConc\ t1\ t2) = c\ (tFold0\ c\ t1)\ (tFold0\ c\ t2)$

by (*unfold tFold0-def, simp*)

4.2 tFoldC

consts *tFoldC* :: ($'h \Rightarrow 'c$) \Rightarrow ($'c \Rightarrow 'c \Rightarrow 'c$) \Rightarrow ($'h, 'u$) $T \Rightarrow 'c$

defs *tFoldC-def*: $tFoldC\ a == tFold\ (\% h\ t . a\ h)$

4.3 hHead

constdefs

$hHead :: ('h, 't)\ T \Rightarrow ('h * 't)$

$hHead == tFold\ Pair\ const$

lemma *hHead-addH*[*simp*]: $hHead\ (addH\ h\ t) = (h, t)$

by (*simp add: hHead-def*)

lemma *hHead-uncurry-addH*[*simp*]: $hHead\ (uncurry\ addH\ p) = p$

by (*simp add: hHead-def*)

lemma *hHead-hCons*[*simp*]: $hHead\ (hCons\ p\ t) = p$

by (*simp add: hHead-def*)

lemma *hHead-hConc*[*simp*]: $hHead\ (hConc\ t1\ t2) = hHead\ t1$

by (*simp add: hHead-def*)

constdefs

$hLength :: ('h, 't)\ T \Rightarrow nat$

$hLength == tFold\ (\% h\ t . 1)\ (\% x\ y . x + y)$

lemma *hLength-addH*[*simp*]: $hLength\ (addH\ h\ t) = 1$

by (*simp add: hLength-def*)

lemma *hLength-hCons[simp]*: $hLength (hCons p t) = Suc (hLength t)$
by (*simp add: hLength-def*)

lemma *hLength-hConc[simp]*: $hLength (hConc t1 t2) = hLength t1 + hLength t2$
apply (*subgoal-tac assoc ((op +) :: nat \Rightarrow nat \Rightarrow nat)*)
prefer 2
apply (*rule assoc-intro, simp*)
apply (*simp add: hLength-def*)
done

lemma *hLength-pos[simp]*: $0 < hLength t$
by (*induct-tac t rule: T-induct, simp-all*)

4.4 The “Cons View”

constdefs

$$\begin{aligned}
 hUnCons0 &:: ((h * 't) + ((h * 't) * (h, 't) T)) * (h, 't) T \Rightarrow \\
 &((h * 't) + ((h * 't) * (h, 't) T)) * (h, 't) T \Rightarrow \\
 &((h * 't) + ((h * 't) * (h, 't) T)) * (h, 't) T \\
 hUnCons0 p1 p2 &== (case fst p1 of \\
 &Inl p \Rightarrow Inr (p, snd p2) \\
 &| Inr (p, t) \Rightarrow Inr (p, hConc t (snd p2)) \\
 &, hConc (snd p1) (snd p2) \\
 &)
 \end{aligned}$$

lemma *snd-hUnCons0-p*: $snd (hUnCons0 p1 p2) = hConc (snd p1) (snd p2)$
by (*simp add: hUnCons0-def*)

lemma *fst-hUnCons0-p*: $fst (hUnCons0 p1 p2) = (case fst p1 of$
 $Inl p \Rightarrow Inr (p, snd p2)$
 $| Inr (p, t) \Rightarrow Inr (p, hConc t (snd p2)))$
by (*simp add: hUnCons0-def*)

lemma *snd-hUnCons0[simp]*: $snd (hUnCons0 (r1, t1) (r2, t2)) = hConc t1 t2$
by (*simp add: snd-hUnCons0-p*)

lemma *assoc-hUnCons0[simp]*: *assoc hUnCons0*
apply (*rule assoc-intro*)
apply (*simp add: hUnCons0-def*)
apply (*split sum.split, rule conjI*)
apply (*intro strip, simp*)
apply (*intro strip, simp*)
apply (*split sum.split, rule conjI*)
apply (*intro strip*)
apply (*case-tac b, simp*)
apply (*intro strip, case-tac b, simp*)
apply (*rotate-tac -2, drule sym, simp*)
done

constdefs

$hUnCons1 :: ('h, 't) T \Rightarrow (('h * 't) + (('h * 't) * ('h, 't) T)) * ('h, 't) T$
 $hUnCons1 == tFold (\% h t . (Inl (h, t), addH h t)) hUnCons0$

lemma $hUnCons1\text{-addH}[simp]$: $hUnCons1 (addH h t) = (Inl (h, t), addH h t)$
by ($simp$ add : $hUnCons1\text{-def}$)

lemma $hUnCons1\text{-hCons}[simp]$: $hUnCons1 (hCons p t) = (Inr (p, t), hCons p t)$
apply ($simp$ add : $hUnCons1\text{-def}$ $hUnCons0\text{-def}$)
apply ($induct\text{-tac}$ t $rule$: $T\text{-induct}$)
apply ($simp$ add : $hCons$)
apply ($simp$ add : $snd\text{-}hUnCons0\text{-}p$ $hCons$)
done

lemma $hUnCons1\text{-hConc}[simp]$:
 $hUnCons1 (hConc t1 t2) = hUnCons0 (hUnCons1 t1) (hUnCons1 t2)$
by ($simp$ add : $hUnCons1\text{-def}$)

lemma $snd\text{-}hUnCons1[simp]$: $snd (hUnCons1 t) = t$
by ($induct\text{-tac}$ t $rule$: $T\text{-induct}$, $simp\text{-all}$ add : $snd\text{-}hUnCons0\text{-}p$)

constdefs

$hUnCons :: ('h, 't) T \Rightarrow (('h * 't) + (('h * 't) * ('h, 't) T))$
 $hUnCons t == fst (hUnCons1 t)$

lemma $hUnCons\text{-addH}[simp]$: $hUnCons (addH h t) = Inl (h, t)$
by ($simp$ add : $hUnCons\text{-def}$)

lemma $hUnCons\text{-hCons}[simp]$: $hUnCons (hCons p t) = Inr (p, t)$
by ($simp$ add : $hUnCons\text{-def}$)

lemma $hUnCons\text{-hConc}[simp]$: $hUnCons (hConc t1 t2) = (case hUnCons t1 of$
 $Inl p \Rightarrow Inr (p, t2)$
 $| Inr (p, t) \Rightarrow Inr (p, hConc t t2))$
apply ($simp$ add : $hUnCons\text{-def}$ $fst\text{-}hUnCons0\text{-}p$)
apply ($subst$ $snd\text{-}hUnCons1$, $simp$)
done

consts

$tFoldr0 :: ('h \Rightarrow 't \Rightarrow 'r) * ('h \Rightarrow 't \Rightarrow 'r \Rightarrow 'r) * ('h, 't) T \Rightarrow 'r$

lemma $tFoldr0\text{-measure}[simp]$:
 $\forall h t t'. (\exists ta. hUnCons t = Inr ((h, ta), t')) \longrightarrow hLength t' < hLength t$
apply ($rule$ $allI$)
apply ($rule$ $allI$)
apply ($rule\text{-tac}$ $x=h$ **in** $spec$)
apply ($induct\text{-tac}$ t $rule$: $T\text{-induct}$, $auto$)
apply ($case\text{-tac}$ $hUnCons t1$, $simp$)
apply ($case\text{-tac}$ b , $simp$)

apply (*erule conjE*, *rotate-tac -1*, *drule sym*, *simp*)
done

recdef *tFoldr0* *measure* (% (a,f,t) . *hLength* t)
tFoldr0-def: *tFoldr0* (a,f,t) = (case *hUnCons* t of
 Inl (h,t) \Rightarrow a h t
 | *Inr* ((h,t),t') \Rightarrow f h t (*tFoldr0* (a,f,t')))

constdefs

tFoldr :: ('h \Rightarrow 't \Rightarrow 'r) \Rightarrow ('h \Rightarrow 't \Rightarrow 'r \Rightarrow 'r) \Rightarrow ('h,'t) T \Rightarrow 'r
tFoldr a f == % t . *tFoldr0* (a,f,t)

lemma *tFoldr-addH*[*simp*]: *tFoldr* a f (*addH* h t) = a h t
by (*simp* *add*: *tFoldr-def tFoldr0-def*)

lemma *tFoldr-hCons*[*simp*]: *tFoldr* a f (*hCons* p t) = *uncurry* f p (*tFoldr* a f t)
by (*case-tac* p, *simp* *add*: *tFoldr-def tFoldr0-def*)

lemma *tFoldr-hConc*:

tFoldr a f (*hConc* t1 t2) = *tFoldr* (% h t . f h t (*tFoldr* a f t2)) f t1
apply (*induct-tac* t1 *rule*: *T-cons-induct*, *simp*)
apply (*subst* *hCons-p* [*THEN* *sym*], *simp-all*)
done

lemma *hConc-via-tFoldr*:

hConc t1 t2 = *tFoldr* (% h t . *hCons* (h,t) t2) (*curry* *hCons*) t1
by (*induct-tac* t1 *rule*: *T-cons-induct*, *simp-all*, *simp* *add*: *hCons*)

lemma *tFold-via-tFoldr*:

tFold a c t = *tFoldr* a (% h t r . c (a h t) r) t
by (*induct-tac* t *rule*: *T-cons-induct*, *simp-all*)

4.5 Mapping

constdefs

hMap :: ('h1 \Rightarrow 'h2) \Rightarrow ('h1,'t) T \Rightarrow ('h2,'t) T
hMap f == *tFold* (*addH* o f) (*hConc*)
tMap :: ('t1 \Rightarrow 't2) \Rightarrow ('h,'t1) T \Rightarrow ('h,'t2) T
tMap f == *tFold* (% h t . *addH* h (f t)) (*hConc*)

lemma *tMap-addH*[*simp*]: *tMap* f (*addH* h t) = *addH* h (f t)
by (*simp* *add*: *tMap-def*)

lemma *tMap-hConc*[*simp*]: *tMap* f (*hConc* t1 t2) = *hConc* (*tMap* f t1) (*tMap* f t2)
by (*simp* *add*: *tMap-def*)

lemma *tMap-hCons*[*simp*]: *tMap* f (*hCons* p t) = *hCons* (*fst* p, f (*snd* p)) (*tMap* f t)

by (*simp add: tMap-def hCons*)

lemma *hMap-addH[simp]*: $hMap\ f\ (addH\ h\ t) = addH\ (f\ h)\ t$
by (*simp add: hMap-def*)

lemma *hMap-hConc[simp]*: $hMap\ f\ (hConc\ t1\ t2) = hConc\ (hMap\ f\ t1)\ (hMap\ f\ t2)$
by (*simp add: hMap-def*)

lemma *hMap-hCons[simp]*: $hMap\ f\ (hCons\ p\ t) = hCons\ (f\ (fst\ p),\ snd\ p)\ (hMap\ f\ t)$
by (*simp add: hMap-def hCons*)

lemma *tFold-tMap[simp]*:
 $assoc\ c \implies tFold\ a\ c\ (tMap\ f\ t) = tFold\ (\% h\ t0 . a\ h\ (f\ t0))\ c\ t$
by (*induct-tac t rule: T-induct, simp-all*)

lemma *tMap-tFold-hConc*:
 $tMap\ f\ (tFold\ a\ hConc\ t) = tFold\ (\lambda h0\ t0 . tMap\ f\ (a\ h0\ t0))\ hConc\ t$
by (*induct-tac t rule: T-induct, simp-all*)

lemma *tMap-tMap*: $tMap\ f\ (tMap\ g\ t) = tMap\ (f\ o\ g)\ t$
by (*induct-tac t rule: T-induct, simp-all*)

lemma *tMap-const-tMap[simp]*:
 $tMap\ (const\ t1)\ (tMap\ f\ t) = tMap\ (const\ t1)\ t$
by (*subst tMap-tMap, simp add: comp-def*)

lemma *tMap-CONST-tMap[simp]*:
 $tMap\ (\lambda x . t1)\ (tMap\ f\ t) = tMap\ (const\ t1)\ t$
by (*subst tMap-tMap, simp add: comp-def*)

4.6 Headers and Subtables

constdefs

$headers :: ('h, 't)\ T \Rightarrow 'h\ neList$
 $headers == tFold\ (\% h\ t . singleton\ h)\ append$
 $subtables :: ('h, 't)\ T \Rightarrow 't\ neList$
 $subtables == tFold\ (\% h\ t . singleton\ t)\ append$

lemma *headers-addH[simp]*: $headers\ (addH\ h\ t) = singleton\ h$
by (*simp add: headers-def*)

lemma *headers-hCons[simp]*: $headers\ (hCons\ p\ t) = cons1\ (fst\ p)\ (headers\ t)$
by (*simp add: headers-def cons1*)

lemma *headers-hConc[simp]*: $headers\ (hConc\ t1\ t2) = append\ (headers\ t1)\ (headers\ t2)$
by (*simp add: headers-def*)

lemma *subtables-addH[simp]*: $\text{subtables } (\text{addH } h \ t) = \text{singleton } t$
by (*simp add: subtables-def*)

lemma *subtables-hCons[simp]*: $\text{subtables } (\text{hCons } p \ t) = \text{cons1 } (\text{snd } p) (\text{subtables } t)$
by (*simp add: subtables-def cons1*)

lemma *subtables-hConc[simp]*: $\text{subtables } (\text{hConc } t1 \ t2) = \text{append } (\text{subtables } t1) (\text{subtables } t2)$
by (*simp add: subtables-def*)

lemma *headers-eq-singleton-0*:
 $\text{headers } t = \text{singleton } h \longrightarrow (\exists t0 . t = \text{addH } h \ t0)$
apply (*induct-tac t rule: T-cons-induct, auto*)
apply (*rule-tac x=t0 in exI, drule singleton-inj, simp*)
done

lemma *headers-eq-singleton[dest?]*:
 $\text{headers } t = \text{singleton } h \implies (\exists t0 . t = \text{addH } h \ t0)$
by (*rule headers-eq-singleton-0 [THEN mp]*)

lemma *headers-eq-cons1-0*:
 $\text{headers } t1 = \text{cons1 } h \ hs \longrightarrow$
 $(\exists t0 \ t2 . t1 = \text{hCons } (h, t0) \ t2 \ \& \ \text{headers } t2 = hs)$
apply (*induct-tac t1 rule: T-cons-induct, simp-all*)
apply (*intro strip*)
apply (*drule cons1-inj, erule conjE, simp*)
apply (*rule-tac x=snd p in exI*)
apply (*rule-tac x=t2 in exI*)
apply (*rotate-tac -2, drule sym, simp*)
done

lemma *headers-eq-cons1[dest?]*:
 $\text{headers } t1 = \text{cons1 } h \ hs \implies$
 $(\exists t0 \ t2 . t1 = \text{hCons } (h, t0) \ t2 \ \& \ \text{headers } t2 = hs)$
by (*rule headers-eq-cons1-0 [THEN mp]*)

lemma *headers-tMap[simp]*: $\text{headers } (\text{tMap } f \ t) = \text{headers } t$
by (*induct-tac t rule: T-induct, simp-all*)

lemma *headers-tFold-hConc[simp]*:
 $\text{headers } (\text{tFold } a \ \text{hConc } t) = \text{tFold } (\% \ h \ t . \ \text{headers } (a \ h \ t)) \ \text{append } t$
by (*induct-tac t rule: T-induct, simp-all*)

lemma *tMap-CONST--headers*:
 $\text{tMap } (\lambda \ x . \ c) \ t = \text{foldr1 } \text{hConc } (\text{map1 } (\lambda \ h . \ \text{addH } h \ c) (\text{headers } t))$
by (*induct-tac t rule: T-induct, simp-all*)

lemma *tFold-CONST-hConc-headers*:
 $tFold (\lambda h (t :: unit). c t) hConc t = foldr1 hConc (map1 (\lambda h . c ())) (headers t)$
by (*induct-tac t rule: T-induct, simp-all*)

4.7 Regular Skeletons

constdefs
 $regSkelStep :: ('t \Rightarrow 's\ option) \Rightarrow ('h, 't) T \Rightarrow ('h\ neList * 's)\ option$
 $regSkelStep\ rs\ t == option-map (\% s . (headers\ t,\ s))$
 $(tFold (\% h\ t . rs\ t)\ optEq\ t)$

lemma *regSkelStep-addH[simp]*:
 $regSkelStep\ rs\ (addH\ h\ t) = option-map (\% s . (singleton\ h,\ s)) (rs\ t)$
by (*simp add: regSkelStep-def*)

lemma *regSkelStep-hConc[simp]*:
 $regSkelStep\ rs\ (hConc\ t1\ t2) =$
 $optThen (regSkelStep\ rs\ t1) (\lambda p1 .$
 $optThen (regSkelStep\ rs\ t2) (\lambda p2 .$
 $if\ snd\ p1 = snd\ p2$
 $then\ Some (append (fst\ p1) (fst\ p2), snd\ p1)$
 $else\ None)$
apply (*simp add: regSkelStep-def*)
apply (*cases (tFold (\lambda h. rs) optEq t1), simp add: optEq-N*)
apply (*cases (tFold (\lambda h. rs) optEq t2), simp*)
apply *simp*
done

lemma *regSkelStep-hCons[simp]*:
 $regSkelStep\ rs\ (hCons\ p\ t) =$
 $optThen (rs (snd\ p))$
 $(\lambda hs2. optThen (regSkelStep\ rs\ t)$
 $(\lambda p2. if\ hs2 = snd\ p2$
 $then\ Some (cons1 (fst\ p) (fst\ p2), hs2) else\ None))$
apply (*simp add: hCons cons1*)
apply (*cases rs (snd p), simp-all add: cons1*)
apply (*cases regSkelStep rs t, simp-all add: cons1*)
done

constdefs
 $regSkelOuter1 :: ('h, 't) T \Rightarrow 'h\ neList\ option$
 $regSkelOuter1\ t == Some (headers\ t)$

lemma *regSkelOuter1-addH[simp]*:
 $regSkelOuter1 (addH\ h\ t) = Some (singleton\ h)$
by (*simp add: regSkelOuter1-def*)

lemma *regSkelOuter1-hConc[simp]*:

regSkelOuter1 (hConc t1 t2) = Some (append (headers t1) (headers t2))
by (*simp add: regSkelOuter1-def*)

lemma *regSkelOuter1-hCons[simp]*:
regSkelOuter1 (hCons p t) = Some (cons1 (fst p) (headers t))
by (*simp add: hCons cons1*)

constdefs
*regSkelOuter2 :: ('h1,('h2,'t) T) T \Rightarrow ('h1 neList * ('h2 neList)) option*
regSkelOuter2 t == regSkelStep regSkelOuter1 t

lemma *regSkelOuter2-addH[simp]*:
regSkelOuter2 (addH h t) = Some (singleton h, headers t)
by (*simp add: regSkelOuter2-def regSkelOuter1-def*)

lemma *regSkelOuter2-hCons[simp]*:
regSkelOuter2 (hCons p t2) = optThen (regSkelOuter2 t2)
($\lambda p2$. if headers (snd p) = snd p2
then Some (cons1 (fst p) (fst p2), snd p2) else None)
by (*case-tac regSkelOuter2 t2, simp-all add: regSkelOuter2-def regSkelOuter1-def*)

lemma *regSkelOuter2-hConc[simp]*:
regSkelOuter2 (hConc t1 t2) = optThen (regSkelOuter2 t1)
($\lambda p1$. optThen (regSkelOuter2 t2)
($\lambda p2$. if snd p1 = snd p2
then Some (append (fst p1) (fst p2), snd p1) else None))
by (*simp add: regSkelOuter2-def*)

lemma *regSkelOuter2-addH-eq-Some*:
regSkelOuter2 (addH h t) = Some (hs1, hs2) \implies
hs1 = singleton h & hs2 = headers t
by (*simp add: regSkelOuter1-def regSkelOuter2-def*)

lemma *regSkelOuter2-hConc-eq-Some*:
regSkelOuter2 (hConc t1 t2) = Some (hs1, hs2) \implies
EX hs1a hs1b . append hs1a hs1b = hs1 &
regSkelOuter2 t1 = Some (hs1a, hs2) &
regSkelOuter2 t2 = Some (hs1b, hs2)

apply *simp*
apply (*drule optThen-result-Some, erule exE, erule conjE*)
apply (*drule optThen-result-Some, erule exE, erule conjE*)
apply (*split split-if-asm*)
apply (*rule-tac x=fst y in exI*)
apply (*rule-tac x=fst ya in exI*)
apply *simp*
apply (*erule conjE, rule conjI*)
apply (*rotate-tac 2, drule sym, simp*)
apply (*rotate-tac -1, drule sym, simp*)
apply *fastsimp*

done

lemma *regSkelOuter2-hCons-eq-Some*:

$regSkelOuter2 (hCons p t2) = Some (hs1, hs2) \implies$
 $EX h1 t1 hs1b . p = (h1, t1) \ \&$
 $hs1 = cons1 h1 hs1b \ \&$
 $headers t1 = hs2 \ \&$
 $regSkelOuter2 t2 = Some (hs1b, hs2)$

apply (*simp only: cons1 hCons*)

apply (*drule regSkelOuter2-hConc-eq-Some*)

apply (*erule exE, erule exE, erule conjE, erule conjE*)

apply (*rule-tac x=fst p in exI*)

apply (*rule-tac x=snd p in exI*)

apply (*rule-tac x=hs1b in exI*)

apply *simp*

done

lemma *regSkelOuter2-eq-Some-0*:

$ALL hs1 hs2 . regSkelOuter2 t = Some (hs1, hs2) \longrightarrow headers t = hs1$

apply (*induct-tac t rule: T-induct, simp-all*)

apply (*intro strip*)

apply (*erule exE, drule optThen-result-Some, simp*)

apply (*erule exE, erule exE, erule conjE, drule optThen-result-Some, simp*)

apply (*erule exE, erule exE, erule conjE, simp*)

apply (*split split-if-asm, simp-all*)

done

lemma *regSkelOuter2-eq-Some[simp]*:

$regSkelOuter2 t = Some (hs1, hs2) \implies headers t = hs1$

by (*insert regSkelOuter2-eq-Some-0, auto*)

lemma *regSkelOuter2-tMap-const[simp]*:

$regSkelOuter2 (tMap (const t2) t1) = Some (headers t1, headers t2)$

by (*induct-tac t1 rule: T-induct, simp-all*)

lemma *regSkelOuter2-tMap-CONST[simp]*:

$regSkelOuter2 (tMap (\% x . t2) t1) = Some (headers t1, headers t2)$

by (*induct-tac t1 rule: T-induct, simp-all*)

lemma *regSkelOuter2-tMap-h-0*:

$\llbracket \bigwedge t. headers (h t) = headers t; \bigwedge t. h (h t) = h t;$

$\bigwedge t1 t2. h (hConc t1 t2) = hConc (h t1) (h t2)$

$\rrbracket \implies$

$ALL hs1 hs2 .$

$regSkelOuter2 t = Some (hs1, hs2) \longrightarrow$

$regSkelOuter2 (tMap h t) = Some (hs1, hs2)$

apply (*induct-tac t rule: T-cons-induct, simp-all*)

apply (*intro strip*)

apply (*drule optThen-result-Some, erule exE, erule conjE*)

```

apply (split split-if-asm) prefer 2 apply simp
apply (case-tac p, case-tac y, simp)
done

```

```

lemma regSkelOuter2-tMap-h[simp]:
  [ [  $\bigwedge t. \text{headers } (h \ t) = \text{headers } t; \bigwedge t. h \ (h \ t) = h \ t;$ 
     $\bigwedge t1 \ t2. h \ (h\text{Conc } t1 \ t2) = h\text{Conc } (h \ t1) \ (h \ t2);$ 
     $\text{regSkelOuter2 } t = \text{Some } (hs1, hs2)$ 
  ]  $\implies$ 
   $\text{regSkelOuter2 } (t\text{Map } h \ t) = \text{Some } (hs1, hs2)$ 
by (insert regSkelOuter2-tMap-h-0 [of h t], auto)

```

4.8 Two-Dimensional Regular Tables

constdefs

```

regularOuter2 :: ('h1,('h2,'t) T) T  $\Rightarrow$  bool
regularOuter2 t == option False (% x . True) (regSkelOuter2 t)

```

```

lemma regularOuter2I[simp,intro]:
   $\text{regSkelOuter2 } t = \text{Some } (hs1, hs2) \implies \text{regularOuter2 } t$ 
by (simp add: regularOuter2-def)

```

```

lemma regularOuter2[simp,dest]:
   $\text{regularOuter2 } t \implies \text{EX } hs1 \ hs2 . \text{regSkelOuter2 } t = \text{Some } (hs1, hs2)$ 
by (case-tac regSkelOuter2 t, simp-all add: regularOuter2-def)

```

```

typedef (regT2) ('h1,'h2,'t) regT2 =
  { t :: ('h1,('h2,'t) T) T . regularOuter2 t }
apply (rule-tac x=addH arbitrary (addH arbitrary arbitrary) in exI)
apply (simp add: regularOuter2-def)
done

```

```

lemma regT2[simp,intro!]:  $\text{regSkelOuter2 } t = \text{Some } (hs1, hs2) \implies t \in \text{regT2}$ 
by (simp add: regT2-def, fast)

```

consts

```

regSkel2 :: ('h1,'h2,'t) regT2  $\Rightarrow$  'h1 neList * 'h2 neList
reg2dim1 :: ('h1,'h2,'t) regT2  $\Rightarrow$  'h1 neList
reg2dim2 :: ('h1,'h2,'t) regT2  $\Rightarrow$  'h2 neList

```

defs

```

regSkel2-def: regSkel2 t == option arbitrary id (regSkelOuter2 (Rep-regT2 t))

```

```

lemma regSkel2[simp]:
   $\text{regSkelOuter2 } t = \text{Some } (hs1, hs2) \implies \text{regSkel2 } (\text{Abs-regT2 } t) = (hs1, hs2)$ 
apply (simp add: regSkel2-def regT2)
apply (subst Abs-regT2-inverse, fast)
apply simp
done

```


defs

reg2dim1-def: $\text{reg2dim1} == \text{fst} \circ \text{regSkel2}$
reg2dim2-def: $\text{reg2dim2} == \text{snd} \circ \text{regSkel2}$

lemma *reg2dim1[simp]*:

$\text{regSkelOuter2 } t = \text{Some } (hs1, hs2) \implies \text{reg2dim1 } (\text{Abs-regT2 } t) = hs1$
by (*auto simp add: reg2dim1-def dest: regSkel2*)

lemma *reg2dim2[simp]*:

$\text{regSkelOuter2 } t = \text{Some } (hs1, hs2) \implies \text{reg2dim2 } (\text{Abs-regT2 } t) = hs2$
by (*auto simp add: reg2dim2-def dest: regSkel2*)

4.9 List Interface

constdefs

tList :: $('h, 't) T \Rightarrow ('h * 't) \text{neList}$
tList == $\text{tFold } (\% h t . \text{singleton } (h,t)) \text{append}$
tOfList :: $('h * 't) \text{neList} \Rightarrow ('h, 't) T$
tOfList == $\text{foldr1 } h\text{Conc } o \text{map1 } (\text{uncurry } \text{addH})$
zipT :: $'h \text{neList} \Rightarrow 't \text{neList} \Rightarrow ('h, 't) T$
zipT *hs ts* == $\text{foldr1 } h\text{Conc } (\text{zipWith1 } \text{addH } hs ts)$

lemma *tList-addH[simp]*: $\text{tList } (\text{addH } h t) = \text{singleton } (h,t)$
by (*simp add: tList-def*)

lemma *tList-hConc[simp]*: $\text{tList } (h\text{Conc } t1 t2) = \text{append } (\text{tList } t1) (\text{tList } t2)$
by (*simp add: tList-def*)

lemma *tList-hCons[simp]*: $\text{tList } (h\text{Cons } p t2) = \text{cons1 } p (\text{tList } t2)$
by (*simp add: hCons cons1*)

lemma *tOfList-singleton[simp]*: $\text{tOfList } (\text{singleton } (h,t)) = \text{addH } h t$
by (*simp add: tOfList-def*)

lemma *tOfList-append[simp]*: $\text{tOfList } (\text{append } t1 t2) = h\text{Conc } (\text{tOfList } t1) (\text{tOfList } t2)$
by (*simp add: tOfList-def*)

lemma *tOfList-cons1[simp]*: $\text{tOfList } (\text{cons1 } p ps) = h\text{Cons } p (\text{tOfList } ps)$
by (*simp add: tOfList-def cons1 hCons*)

4.10 ZipWith

constdefs

tZipWith :: $((('h1 * 't1) \Rightarrow ('h2 * 't2) \Rightarrow ('h * 't)) \Rightarrow ('h1, 't1) T \Rightarrow ('h2, 't2) T \Rightarrow ('h, 't) T)$
tZipWith *f t1 t2* ==
 $\text{foldr1 } h\text{Conc } (\text{map1 } (\text{uncurry } \text{addH}) (\text{zipWith1 } f (\text{tList } t1) (\text{tList } t2)))$

lemma *tZipWith-A-A[simp]*:
 $tZipWith\ f\ (addH\ h1\ t1)\ (addH\ h2\ t2) = uncurry\ addH\ (f\ (h1,t1)\ (h2,t2))$
by (*simp add: tZipWith-def*)

lemma *tZipWith-A-C[simp]*:
 $tZipWith\ f\ (addH\ h1\ t1)\ (hCons\ p2\ t2) = uncurry\ addH\ (f\ (h1,t1)\ p2)$
by (*simp add: tZipWith-def*)

lemma *tZipWith-C-A[simp]*:
 $tZipWith\ f\ (hCons\ p1\ t1)\ (addH\ h2\ t2) = uncurry\ addH\ (f\ p1\ (h2,t2))$
by (*simp add: tZipWith-def*)

lemma *tZipWith-C-C[simp]*:
 $tZipWith\ f\ (hCons\ p1\ t1)\ (hCons\ p2\ t2) = hCons\ (f\ p1\ p2)\ (tZipWith\ f\ t1\ t2)$

proof –

have $tZipWith\ f\ (hCons\ p1\ t1)\ (hCons\ p2\ t2) = foldr1\ hCons$
 $(map1\ (\lambda p.\ addH\ (fst\ p)\ (snd\ p))$
 $(zipWith1\ f\ (append\ (singleton\ p1)\ (tList\ t1))$
 $(append\ (singleton\ p2)\ (tList\ t2))))$
by (*simp add: tZipWith-def hCons*)
also have $\dots = hCons\ (f\ p1\ p2)\ (tZipWith\ f\ t1\ t2)$
by (*simp add: tZipWith-def hCons cons1 [THEN sym]*)
finally show *?thesis* .

qed

lemma *tZipWith-A[simp]*:
 $tZipWith\ f\ (addH\ h1\ t1)\ t2 = uncurry\ addH\ (f\ (h1,t1)\ (hHead\ t2))$
by (*induct-tac t2 rule: T-cons-induct, simp-all*)

lemma *tZipWith-C[simp]*:
 $hHead\ (tZipWith\ f\ (hCons\ p1\ t1)\ t2) = f\ p1\ (hHead\ t2)$
by (*induct-tac t2 rule: T-cons-induct, simp-all*)

lemma *tZipWith--A[simp]*:
 $tZipWith\ f\ t1\ (addH\ h2\ t2) = uncurry\ addH\ (f\ (hHead\ t1)\ (h2,t2))$
by (*induct-tac t1 rule: T-cons-induct, simp-all*)

lemma *tZipWith--C[simp]*:
 $hHead\ (tZipWith\ f\ t1\ (hCons\ p2\ t2)) = f\ (hHead\ t1)\ p2$
by (*induct-tac t1 rule: T-cons-induct, simp-all*)

lemma *tZipWith-assoc-0[simp]*:
 $assoc\ f \implies$
 $ALL\ t1\ t2\ t3.\ tZipWith\ f\ (tZipWith\ f\ t1\ t2)\ t3 = tZipWith\ f\ t1\ (tZipWith\ f\ t2\ t3)$
apply (*rule T-cons-inductA*)
apply (*rule T-cons-inductA, simp-all*)
apply (*rule T-cons-inductA, simp-all*)
apply (*rule T-cons-inductA, simp-all*)

done

lemma *assoc-tZipWith[simp]*: $\text{assoc } f \implies \text{assoc } (tZipWith f)$
by (*rule assoc-intro, simp*)

4.11 Collapsing

constdefs

$\text{collapse} :: ('h1 \Rightarrow 't1 \Rightarrow ('h2, 't2) T) \Rightarrow ('h1, 't1) T \Rightarrow ('h2, 't2) T$
 $\text{collapse } f == tFold f hConc$

lemma *collapse-addH[simp]*: $\text{collapse } f (\text{addH } h t) = f h t$
by (*simp add: collapse-def*)

lemma *collapse-hConc[simp]*:

$\text{collapse } f (hConc t1 t2) = hConc (\text{collapse } f t1) (\text{collapse } f t2)$

by (*simp add: collapse-def*)

constdefs

$\text{collapse2} :: ('h1 \Rightarrow 'h2 \Rightarrow 'h3) \Rightarrow ('h1, ('h2, 't2) T) T \Rightarrow ('h3, 't2) T$
 $\text{collapse2 } f == \text{collapse } (\% h1 . hMap (f h1))$

lemma *collapse2-addH[simp]*: $\text{collapse2 } f (\text{addH } h t) = hMap (f h) t$
by (*simp add: collapse2-def*)

lemma *collapse2-hConc[simp]*:

$\text{collapse2 } f (hConc t1 t2) = hConc (\text{collapse2 } f t1) (\text{collapse2 } f t2)$

by (*simp add: collapse2-def*)

theorem *collapse2-tFold2*:

$\llbracket \text{assoc } c1; \text{assoc } c2;$
 $\bigwedge h1 h2 x . a1 h1 (a2 h2 x) = a3 (f h1 h2) x;$
 $\bigwedge h x y . a1 h (c2 x y) = c1 (a1 h x) (a1 h y)$
 $\rrbracket \implies$

$tFold a1 c1 (tMap (tFold a2 c2) t) = tFold a3 c1 (\text{collapse2 } f t)$

apply (*induct-tac t rule: T-induct, simp*)

apply (*induct-tac t0 rule: T-induct, simp-all*)

done

theorem *collapse2-tFold2-wrapped*:

$\llbracket \text{assoc } c1; \text{assoc } c2; \text{assoc } c3;$
 $\bigwedge h1 h2 x . w1 (a1 h1 (a2 h2 x)) = w3 (a3 (f h1 h2) x);$
 $\bigwedge h x y . w1 (a1 h (c2 x y)) = c4 (w1 (a1 h x)) (w1 (a1 h y));$
 $\bigwedge x y . w1 (c1 x y) = c5 (w1 x) (w1 y);$

(* These swapped equations look more intuitive, but destroy simplification:

$\bigwedge x y . w3 (c3 x y) = c4 (w3 x) (w3 y);$
 $\bigwedge x y . w3 (c3 x y) = c5 (w3 x) (w3 y)$

*)

$\bigwedge x y . c4 (w3 x) (w3 y) = w3 (c3 x y);$

```

     $\bigwedge x y . c5 (w3 x) (w3 y) = w3 (c3 x y)$ 
  ]  $\implies$ 
  w1 (tFold a1 c1 (tMap (tFold a2 c2) t)) = w3 (tFold a3 c3 (collapse2 f t))
apply (induct-tac t rule: T-induct, simp)
apply (induct-tac t0 rule: T-induct, simp-all del: tFold-tMap)
done

```

4.12 Compression

lemma *hCompress-0*:

```

  [ assoc c;
     $\bigwedge h1 h2 t1 t2 . c (a h1 t1) (a h2 t2) = a (c' h1 h2) (c'' t1 t2)$ 
  ]  $\implies$ 
  tFold a c (hCons (h1,t1) (hCons (h2,t2) t)) =
  tFold a c (hCons (c' h1 h2, c'' t1 t2) t)
by (simp del: assoc add: assoc [THEN sym])

```

consts

```

dim1addH :: ('h  $\Rightarrow$  'c  $\Rightarrow$  'v)  $\Rightarrow$  ('h  $\Rightarrow$  ('c,'u) T  $\Rightarrow$  'v)

```

defs

```

dim1addH-def: dim1addH a h == tFold (% c u . a h c) arbitrary

```

lemma *dim1addH[simp]*: *dim1addH a h (cell c) = a h c*

by (simp add: dim1addH-def cell-def)

lemma *hCompress-cells*:

```

  [ assoc c;
     $\bigwedge h1 h2 g1 g2 . c (a h1 g1) (a h2 g2) = a (c' h1 h2) (c'' g1 g2)$ 
  ]  $\implies$ 
  tFold (dim1addH a) c (hCons (h1,cell g1) (hCons (h2,cell g2) t)) =
  tFold (dim1addH a) c (hCons (c' h1 h2, cell (c'' g1 g2)) t)
by (simp del: assoc add: assoc [THEN sym])

```

4.13 Elementary Transformations

The titles of the subsections here are the structural elementary transformations as listed in [?].

The lemmas show the corresponding general properties for the first dimension; for other dimensions; the corresponding properties can be obtained using theorem *tTranspose-tFold2*.

4.13.1 Permuting two (-1) -slices with their corresponding header entries

lemma *hCommute*:

```

  [ assoc c; commutative c ]  $\implies$ 
  tFold a c (hConc t1 t2) = tFold a c (hConc t2 t1)

```

by (*simp*, *erule commutative*)

4.13.2 Deleting a (-1) -slice with “false” in the corresponding header entry

Since Zucker also does not allow empty tables, this deletion is only possible for tables that are in the range of *hConc*.

lemma *hDelLeftUnitHeader*:

$$\llbracket \text{assoc } c; \bigwedge t1\ b . c\ (a\ h1\ t1)\ b = b \rrbracket \implies \\ tFold\ a\ c\ (hConc\ (addH\ h1\ t1)\ t) = tFold\ a\ c\ t$$

by *simp*

lemma *hDelLeftUnitHeader-hCons*:

$$\llbracket \text{assoc } c; \bigwedge t1\ b . c\ (a\ h1\ t1)\ b = b \rrbracket \implies \\ tFold\ a\ c\ (hCons\ (h1,t1)\ t) = tFold\ a\ c\ t$$

by *simp*

lemma *hDelRightUnitHeader*:

$$\llbracket \text{assoc } c; \bigwedge t1\ b . c\ b\ (a\ h1\ t1) = b \rrbracket \implies \\ tFold\ a\ c\ (hConc\ t\ (addH\ h1\ t1)) = tFold\ a\ c\ t$$

by *simp*

4.13.3 Deleting a principal slice with only “false” entries from an inverted table

lemma *hDelLeftUnitSubtable*:

$$\llbracket \text{assoc } c; \bigwedge h1\ b . c\ (a\ h1\ t1)\ b = b \rrbracket \implies \\ tFold\ a\ c\ (hConc\ (addH\ h1\ t1)\ t) = tFold\ a\ c\ t$$

by *simp*

lemma *hDelLeftUnitSubtable-hCons*:

$$\llbracket \text{assoc } c; \bigwedge h1\ b . c\ (a\ h1\ t1)\ b = b \rrbracket \implies \\ tFold\ a\ c\ (hCons\ (h1,t1)\ t) = tFold\ a\ c\ t$$

by *simp*

lemma *hDelRightUnitSubtable*:

$$\llbracket \text{assoc } c; \bigwedge h1\ b . c\ b\ (a\ h1\ t1) = b \rrbracket \implies \\ tFold\ a\ c\ (hConc\ t\ (addH\ h1\ t1)) = tFold\ a\ c\ t$$

by *simp*

4.13.4 Splitting a principal slice by “splitting a disjunction” in the corresponding header in an inverted table

This is strange: the corresponding header of a principal slice in an inverted table is a value header!

Another thing to keep in mind here is that splitting a condition header in a normal or inverted table may turn a proper table into an improper table.

lemma *hSplitHeader*:

$\llbracket \text{assoc } c; \bigwedge h1\ h2\ t. a\ (c'\ h1\ h2)\ t = c\ (a\ h1\ t)\ (a\ h2\ t) \rrbracket \implies$
 $tFold\ a\ c\ (addH\ (c'\ h1\ h2)\ t) =$
 $tFold\ a\ c\ (hConc\ (addH\ h1\ t)\ (addH\ h2\ t))$

by *simp*

4.13.5 Combining two or more principal slices with the same value header entry into a single slice in an inverted table

lemma *hCombineEqualHeaders*:

$\llbracket \text{assoc } c; \bigwedge h\ t1\ t2. c\ (a\ h\ t1)\ (a\ h\ t2) = a\ h\ (c'\ t1\ t2) \rrbracket \implies$
 $tFold\ a\ c\ (hConc\ (addH\ h\ t1)\ (addH\ h\ t2)) =$
 $tFold\ a\ c\ (addH\ h\ (c'\ t1\ t2))$

by *simp*

end

5 Functions Interacting Directly with the Second Table Dimension

theory *Tables2 = Tables*:

5.1 Construction in the Second Dimension

constdefs

$addH2 :: 'h2 \Rightarrow ('h1, 't)\ T \Rightarrow ('h1, ('h2, 't)\ T)\ T$
 $addH2\ h2 == tFold\ (\% h1\ t. addH\ h1\ (addH\ h2\ t))\ hConc$

lemma *addH2-addH[simp]*: $addH2\ h2\ (addH\ h1\ t) = addH\ h1\ (addH\ h2\ t)$
by (*simp add: addH2-def*)

lemma *addH2-hConc[simp]*: $addH2\ h2\ (hConc\ t1\ t2) = hConc\ (addH2\ h2\ t1)\ (addH2\ h2\ t2)$
by (*simp add: addH2-def*)

lemma *addH2-hCons[simp]*: $addH2\ h2\ (hCons\ p\ t) = hCons\ (fst\ p,\ addH\ h2\ (snd\ p))\ (addH2\ h2\ t)$
by (*simp add: addH2-def hCons*)

lemma *hHead-addH2[simp]*: $hHead\ (addH2\ h\ t) = (fst\ (hHead\ t),\ addH\ h\ (snd\ (hHead\ t)))$

apply (*simp add: addH2-def*)
apply (*induct-tac t rule: T-cons-induct*)
apply *simp-all*
done

lemma *addH2-tMap*: $addH2\ h = tMap\ (addH\ h)$
by (*simp add: addH2-def tMap-def*)

lemma *tFold-addH2[simp]*:
 $assoc\ c \implies tFold\ a\ c\ (addH2\ h\ t) = tFold\ (\lambda ha\ t0.\ a\ ha\ (addH\ h\ t0))\ c\ t$
by (*simp add: addH2-tMap*)

lemma *tFold-tFold-addH2*:
 $\llbracket\ assoc\ c2\ \rrbracket \implies$
 $tFold\ a2\ c2\ (tMap\ (tFold\ a1\ c1)\ (addH2\ h\ t)) =$
 $tFold\ (\lambda h0\ t0.\ a2\ h0\ (a1\ h\ t0))\ c2\ t$
by (*induct-tac t rule: T-induct, simp-all*)

lemma *tMap-tFold-addH2*:
 $assoc\ c1 \implies$
 $tMap\ (tFold\ a1\ c1)\ (addH2\ h\ t) = tFold\ (\lambda h0\ t0.\ addH\ h0\ (a1\ h\ t0))\ hConc\ t$
by (*induct-tac t rule: T-induct, simp-all*)

constdefs

$vConc0 :: ('h1 * ('h2, 't) T) \Rightarrow ('h1 * ('h2, 't) T) \Rightarrow ('h1 * ('h2, 't) T)$
 $vConc0 == \% p1\ p2.\ (fst\ p1, hConc\ (snd\ p1)\ (snd\ p2))$

lemma *vConc0[simp]*: $vConc0\ (h1,t1)\ (h2,t2) = (h1, hConc\ t1\ t2)$
by (*simp add: vConc0-def*)

lemma *fst-vConc0[simp]*: $fst\ (vConc0\ p1\ p2) = fst\ p1$
by (*simp add: vConc0-def*)

lemma *snd-vConc0[simp]*: $snd\ (vConc0\ p1\ p2) = hConc\ (snd\ p1)\ (snd\ p2)$
by (*simp add: vConc0-def*)

lemma *vConc0-assoc[simp]*: $vConc0\ (vConc0\ p1\ p2)\ p3 = vConc0\ p1\ (vConc0\ p2\ p3)$

proof –

have $vConc0\ (vConc0\ p1\ p2)\ p3 = vConc0\ (vConc0\ (fst\ p1, snd\ p1)\ (fst\ p2, snd\ p2))\ (fst\ p3, snd\ p3)$ **by** *simp*

also have $\dots = vConc0\ (fst\ p1, snd\ p1)\ (vConc0\ (fst\ p2, snd\ p2)\ (fst\ p3, snd\ p3))$ **by** (*simp only: vConc0 hConc-assoc*)

also have $\dots = vConc0\ p1\ (vConc0\ p2\ p3)$ **by** *simp*

finally show *?thesis* .

qed

lemma *assoc-vConc0[simp]*: $assoc\ vConc0$
by (*rule assoc-intro, simp*)

constdefs

$vConc :: ('h1, ('h2, 't) T) T \Rightarrow ('h1, ('h2, 't) T) T \Rightarrow ('h1, ('h2, 't) T) T$
 $vConc == tZipWith\ vConc0$

lemma *assoc-vConc[simp]*: $assoc\ vConc$

by (*simp add: vConc-def*)

lemma *vConc-addH[*simp*]*:

$vConc (addH h t1) t2 = uncurry addH (vConc0 (h, t1) (hHead t2))$

by (*simp add: vConc-def*)

lemma *hHead-vConc-hCons[*simp*]*:

$hHead (vConc (hCons p1 t1) t2) = vConc0 p1 (hHead t2)$

by (*simp add: vConc-def*)

lemma *vConc--addH[*simp*]*:

$vConc t1 (addH h t2) = uncurry addH (vConc0 (hHead t1) (h, t2))$

by (*simp add: vConc-def*)

lemma *hHead-vConc--hCons[*simp*]*:

$hHead (vConc t1 (hCons p2 t2)) = vConc0 (hHead t1) p2$

by (*simp add: vConc-def*)

lemma *vConc-hCons-hCons[*simp*]*:

$vConc (hCons p1 t1) (hCons p2 t2) = hCons (vConc0 p1 p2) (vConc t1 t2)$

by (*simp add: vConc-def*)

constdefs

$vCons :: ('h2 * ('h1, 't) T) \Rightarrow ('h1, ('h2, 't) T) T \Rightarrow ('h1, ('h2, 't) T) T$
 $vCons p t == vConc (uncurry addH2 p) t$

lemma *vCons: vCons (h,t0) t = vConc (addH2 h t0) t*

by (*simp add: vCons-def*)

lemma *vCons-addH[*simp*]*:

$vCons p (addH h t) = addH (fst (hHead (uncurry addH2 p)))$
 $(hConc (snd (hHead (uncurry addH2 p)))) t$

by (*simp add: vCons-def*)

lemma *hHead-vCons-hCons[*simp*]*:

$hHead (vCons p1 (hCons p2 t2)) = vConc0 (hHead (uncurry addH2 p1)) p2$

by (*simp add: vCons-def*)

lemma *vCons-hCons-hCons[*simp*]*:

$vCons (h, hCons p1 t1) (hCons p2 t2) =$
 $hCons (vConc0 (fst p1, addH h (snd p1)) p2) (vCons (h, t1) t2)$

by (*simp add: vCons-def vConc-def*)

lemma *headers-addH2[*simp*]*: *headers (addH2 h t) = headers t*

by (*induct-tac t rule: T-induct, simp-all*)

lemma *headers-vCons-0*:

ALL t p . headers (snd p) = headers t \longrightarrow headers (vCons p t) = headers t
apply (*rule T-cons-inductA, simp-all*)

apply (*intro strip*, *drule headers-eq-singleton*)
apply (*erule exE*, *simp*)
apply (*intro strip*, *drule headers-eq-cons1*)
apply (*erule exE*, *erule exE*, *erule conjE*, *simp*)
done

lemma *headers-vCons[simp]*:
 $headers\ (snd\ p) = headers\ t \implies headers\ (vCons\ p\ t) = headers\ t$
apply (*insert headers-vCons-0*)
apply (*drule-tac x=t in spec*)
apply (*drule-tac x=p in spec*)
apply *simp*
done

lemma *tMap-h-vConc-0*:

$$\begin{aligned} & \llbracket \bigwedge t . headers\ (h\ t) = headers\ t; \\ & \quad \bigwedge t . h\ (h\ t) = h\ t; \\ & \quad \bigwedge t1\ t2 . h\ (hConc\ t1\ t2) = hConc\ (h\ t1)\ (h\ t2) \\ & \rrbracket \implies \\ & \quad ALL\ t2\ hs1\ hs2a\ hs2b . \\ & \quad \quad regSkelOuter2\ t1 = Some\ (hs1,\ hs2a) \longrightarrow \\ & \quad \quad regSkelOuter2\ t2 = Some\ (hs1,\ hs2b) \longrightarrow \\ & \quad \quad tMap\ h\ (vConc\ t1\ t2) = vConc\ (tMap\ h\ t1)\ (tMap\ h\ t2) \end{aligned}$$

apply (*induct-tac t1 rule: T-cons-induct, simp-all*)
apply (*rule allI*)
apply (*induct-tac t2 rule: T-cons-induct, simp-all*)
apply (*rule allI*)
apply (*induct-tac t2a rule: T-cons-induct, simp-all*)
apply (*intro strip, case-tac p, simp*)
apply (*erule exE, drule optThen-result-Some, erule exE, erule conjE*)
apply (*split split-if-asm*) **prefer** 2 **apply** *simp*
apply (*drule optThen-result-Some, erule exE, erule conjE*)
apply (*split split-if-asm*) **prefer** 2 **apply** *simp*
apply (*case-tac y, case-tac ya, case-tac pa, simp*)
apply (*erule conjE*)
apply (*rotate-tac -2, drule sym, simp, drule cons1-inj, erule conjE, simp*)
done

lemma *tMap-h-vConc[simp]*:

$$\begin{aligned} & \llbracket \bigwedge t . headers\ (h\ t) = headers\ t; \\ & \quad \bigwedge t . h\ (h\ t) = h\ t; \\ & \quad \bigwedge t1\ t2 . h\ (hConc\ t1\ t2) = hConc\ (h\ t1)\ (h\ t2); \\ & \quad \quad regSkelOuter2\ t1 = Some\ (hs1,\ hs2a); \\ & \quad \quad regSkelOuter2\ t2 = Some\ (hs1,\ hs2b) \\ & \rrbracket \implies \\ & \quad tMap\ h\ (vConc\ t1\ t2) = vConc\ (tMap\ h\ t1)\ (tMap\ h\ t2) \end{aligned}$$

by (*insert tMap-h-vConc-0 [of h t1], auto*)

lemma *tMap-const-vConc-0*:

```

ALL t2 hs1a hs1b hs2a hs2b .
  regSkelOuter2 t1 = Some (hs1a, hs2a) →
  regSkelOuter2 t2 = Some (hs1b, hs2b) →
  hs1a = hs1b →
  tMap (const t) (vConc t1 t2) = tMap (const t) t1
apply (induct-tac t1 rule: T-cons-induct, simp)
apply (rule allI)
apply (induct-tac t2a rule: T-cons-induct, simp-all)
apply (intro strip, erule exE)
apply (drule optThen-result-Some, erule exE, erule conjE)
apply (case-tac p, case-tac y, simp)
apply (split split-if-asm, simp, simp)
apply (intro strip, erule exE, erule exE)
apply (drule optThen-result-Some, erule exE, erule conjE)
apply (drule optThen-result-Some, erule exE, erule conjE)
apply (case-tac p, case-tac y, simp)
apply (case-tac pa, case-tac ya, simp)
apply (split split-if-asm, simp)
apply (split split-if-asm, simp)
apply (erule conjE, rotate-tac -2, drule sym, simp)
apply (drule cons1-inj, erule conjE, simp-all)
done

```

```

lemma tMap-const-vConc[simp]:
  [ regSkelOuter2 t1 = Some (hs1a, hs2a);
    regSkelOuter2 t2 = Some (hs1b, hs2b);
    hs1a = hs1b
  ] ⇒
  tMap (const t) (vConc t1 t2) = tMap (const t) t1
by (insert tMap-const-vConc-0 [of t1 t], auto)

```

```

lemma tMap-const-hConc:
  tMap (const (hConc t1 t2)) t = vConc (tMap (const t1) t) (tMap (const t2) t)
apply (cut-tac t2.0=t1 and t1.0=t in regSkelOuter2-tMap-const)
apply (cut-tac t2.0=t2 and t1.0=t in regSkelOuter2-tMap-const)
apply (erule mp1)
apply (erule mp1)
apply (induct-tac t rule: T-cons-induct, simp-all)
done

```

5.2 Regular Skeletons and the Second Dimension

```

lemma regSkelOuter1-addH2[simp]:
  regSkelOuter1 (addH2 h t) = Some (headers t)
by (simp add: regSkelOuter1-def)

```

```

lemma regSkelStep-addH2:
  regSkelStep rs (addH2 h t) = option-map (Pair (headers t)) (tFold (λh. rs) optEq
    (addH2 h t))

```

by (*simp add: regSkelStep-def*)

lemma *regSkelStep-Some-0*:

ALL hs . regSkelStep rs t = Some (hs,r) \longrightarrow headers t = hs

apply (*induct-tac t rule: T-cons-induct, simp-all*)

apply (*intro strip*)

apply (*drule optThen-result-Some*)

apply (*erule exE, erule conjE*)

apply (*drule optThen-result-Some*)

apply (*erule exE, erule conjE*)

apply (*case-tac y = snd ya, simp-all*)

apply (*erule conjE*)

apply (*rotate-tac -1, drule sym, simp*)

apply (*drule-tac x=fst ya in spec*)

apply (*rotate-tac 3, drule sym, simp*)

done

lemma *regSkelStep-Some*:

regSkelStep rs t = Some (hs,r) \implies headers t = hs

apply (*insert regSkelStep-Some-0 [of rs t r]*)

apply (*drule-tac x=hs in spec, simp*)

done

lemma *regSkelOuter2-addH2[*simp*]*:

regSkelOuter2 (addH2 h t) = Some (headers t, singleton h)

by (*induct-tac t rule: T-induct, simp-all*)

lemma *regSkelOuter2-eq-Some*:

regSkelOuter2 t = Some (hs1,hs2) \implies headers t = hs1

by (*simp add: regSkelStep-Some regSkelOuter2-def*)

lemma *regSkelOuter2-vCons-0*:

ALL t p . headers (snd p) = headers t \longrightarrow

regSkelOuter2 (vCons p t) =

optThen (tFold ($\lambda h t . \text{Some (headers t)}$)) optEq (uncurry addH2 p)) ($\lambda hs1 .$

optThen (regSkelOuter2 t) ($\lambda p2 . \text{Some (headers t, append hs1 (snd p2))}$))

apply (*rule T-cons-inductA*)

Case 1: *t=addH*

apply (*simp add: regSkelOuter2-def*)

apply (*intro strip*)

apply (*drule headers-eq-singleton*)

apply (*erule exE*)

apply (*simp add: regSkelOuter1-def regSkelOuter2-def cons1 singleton-def append-def neList-def Abs-neList-inverse*)

Case 2: *t=hCons*

apply (*intro strip*)

apply (*simp add: regSkelOuter2-def*)

apply (*drule headers-eq-cons1*)
apply (*erule exE, erule exE, erule conjE*)
apply (*subgoal-tac vCons pa = vCons (fst pa, hCons (fst p, t0) t2a)*)
prefer 2
apply (*drule sym, simp*)
apply (*rotate-tac -1, drule sym, rotate-tac -1, erule subst*)
apply *simp*
apply (*case-tac regSkelStep regSkelOuter1 t2, simp*)
apply (*simp add: regSkelOuter1-def*)
apply (*case-tac regSkelOuter1 (snd p), simp*)

Case 2.1: *regSkelOuter1 (snd p) = None*

apply (*rule append-inj2-neq-conv*)
apply (*drule-tac x=fst pa in spec*)
apply (*drule-tac x=t2a in spec*)
apply (*simp add: regSkelOuter1-def*)

Case 2.2: *regSkelOuter1 (snd p) = Some aa*

apply *simp*
apply (*rule conjI*)

Case 2.2.1: *aa = snd a*

apply (*intro strip, rule append-inj2-conv*)
apply (*simp add: regSkelOuter1-def*)

Case 2.2.2: *aa ≠ snd a*

apply (*intro strip, rule append-inj2-neq-conv*)
apply (*simp add: regSkelOuter1-def*)
done

lemma *regSkelOuter2-vCons:*

headers (snd p) = headers t \implies
regSkelOuter2 (vCons p t) =
*optThen (tFold ($\lambda h t . \text{Some (headers t)}$)) optEq (uncurry addH2 p)) ($\lambda hs1 .$
*optThen (regSkelOuter2 t) ($\lambda p2 . \text{Some (headers t, append hs1 (snd p2))}$))**

apply (*insert regSkelOuter2-vCons-0*)
apply (*drule-tac x=t in spec*)
apply (*drule-tac x=p in spec*)
apply *simp*
done

lemma *regSkelOuter2-Some-vCons:*

$\llbracket \text{regSkelOuter2 } t = \text{Some (hs1, hs2)}; \text{headers (snd p) = headers } t \rrbracket \implies$
regSkelOuter2 (vCons p t) = Some (hs1, cons1 (fst p) hs2)
by (*simp add: regSkelOuter2-vCons cons1*)

lemma *regSkelOuter2-vConc-0:*

ALL t1 t2 hs1a hs2a hs1b hs2b . headers t1 = headers t2 \longrightarrow

```

    regSkelOuter2 t1 = Some (hs1a, hs2a) →
    regSkelOuter2 t2 = Some (hs1b, hs2b) →
    hs1a = hs1b →
    regSkelOuter2 (vConc t1 t2) = Some (hs1a, append hs2a hs2b)
  apply (rule T-cons-inductA)
  apply (rule T-cons-inductA)
  apply simp
  apply (intro strip)
  apply (simp add: regSkelOuter1-def regSkelOuter2-def)
  apply (rule T-cons-inductA)
  apply simp
  apply simp
  apply (intro strip, drule cons1-inj, erule conjE)
  apply (drule-tac x=t2a in spec)
  apply (simp add: regSkelOuter1-def)
  apply (case-tac regSkelOuter2 t2, simp)
  apply (case-tac regSkelOuter2 t2a, simp-all)
  apply (case-tac headers (snd p) = snd a, simp-all)
  apply (case-tac headers (snd pa) = snd aa, simp-all)
  apply (erule conjE, erule conjE)
  apply (drule-tac x=fst a in spec)
  apply (drule-tac x=hs2a in spec)
  apply (drule mp, fastsimp)
  apply (drule-tac x=hs2b in spec)
  apply (drule mp)
  apply (subgoal-tac cons1 (fst pa) (fst a) = cons1 (fst pa) (fst aa))
  prefer 2
  apply bestsimp
  apply (drule cons1-inj, erule conjE, fastsimp)
  apply simp
done

```

```

lemma regSkelOuter2-vConc[simp]:
  [| regSkelOuter2 t1 = Some (hs1a, hs2a);
    regSkelOuter2 t2 = Some (hs1b, hs2b);
    hs1a = hs1b
  |] ⇒ regSkelOuter2 (vConc t1 t2) = Some (hs1a, append hs2a hs2b)
  apply (insert regSkelOuter2-vConc-0)
  apply (drule-tac x=t1 in spec)
  apply (drule-tac x=t2 in spec)
  apply (drule-tac x=hs1a in spec)
  apply (drule-tac x=hs2a in spec)
  apply (drule-tac x=hs1b in spec)
  apply (drule-tac x=hs2b in spec)
  apply (frule regSkelOuter2-eq-Some)
  apply (rotate-tac 1, frule regSkelOuter2-eq-Some, simp)
done

```

lemma *headers-vConc*[*simp*]:
 \llbracket *regSkelOuter2* *t1* = *Some* (*hs1*,*hs2a*);
regSkelOuter2 *t2* = *Some* (*hs1*,*hs2b*)
 $\rrbracket \implies$ *headers* (*vConc* *t1* *t2*) = *headers* *t1*
apply (*frule-tac* *t1.0=t1* **and** *t2.0=t2* **in** *regSkelOuter2-vConc*, *assumption*)
apply *simp*
apply (*rotate-tac* -1 , *drule* *regSkelOuter2-eq-Some*)
apply (*drule* *regSkelOuter2-eq-Some* [*THEN sym*], *simp*)
done

lemma *vConc-hConc-0*:
 ALL *t3 h1 v1 v2 t2 t4 h2* .
regSkelOuter2 *t1* = *Some* (*h1*,*v1*) \longrightarrow
regSkelOuter2 *t2* = *Some* (*h2*,*v1*) \longrightarrow
regSkelOuter2 *t3* = *Some* (*h1*,*v2*) \longrightarrow
regSkelOuter2 *t4* = *Some* (*h2*,*v2*) \longrightarrow
vConc (*hConc* *t1* *t2*) (*hConc* *t3* *t4*) = *hConc* (*vConc* *t1* *t3*) (*vConc* *t2* *t4*)
apply (*induct-tac* *t1* *rule*: *T-cons-induct*, *simp add*: *regSkelOuter1-def*)
apply (*rule allI*)
apply (*induct-tac* *t3* *rule*: *T-cons-induct*, *simp add*: *regSkelOuter1-def*)
apply (*intro strip*, *drule* *singleton-inj*)
apply (*simp add*: *hCons-p* [*THEN sym*])

1.2

apply (*intro strip*, *simp*)
apply (*drule* *optThen-result-Some*, *erule* *exE*, *erule* *conjE*)
apply (*simp add*: *regSkelOuter1-def*)
apply (*case-tac* *p*, *simp*)
apply (*split split-if-asm*, *simp*, *simp*)

2

apply (*rule allI*)
apply (*erule* *mp1*)
apply (*induct-tac* *t3* *rule*: *T-cons-induct*, *simp add*: *regSkelOuter1-def*)
apply (*intro strip*)
apply (*drule* *optThen-result-Some*, *erule* *exE*, *erule* *conjE*)
apply (*split split-if-asm*, *simp*)
apply (*erule* *conjE*, *rotate-tac* -2 , *drule* *sym*, *simp*, *simp*)

2.2

apply (*intro strip*, *simp add*: *regSkelOuter1-def del*: *vConc-hCons-hCons*)
apply (*drule* *optThen-result-Some*, *erule* *exE*, *erule* *conjE*)+
apply (*case-tac* *y*)
apply (*case-tac* *ya*)
apply (*case-tac* *p*)
apply (*case-tac* *pa*)
apply (*simp del*: *vConc-hCons-hCons*)
apply (*split split-if-asm*, *simp*)

```

apply (split split-if-asm, simp)
apply (erule conjE, rotate-tac -2, drule sym, simp)
apply (drule cons1-inj, erule conjE, simp)
apply simp
apply simp
done

```

lemma *vConc-hConc*:

```

  [| regSkelOuter2 t1 = Some (h1,v1);
    regSkelOuter2 t2 = Some (h2,v1);
    regSkelOuter2 t3 = Some (h1,v2);
    regSkelOuter2 t4 = Some (h2,v2) |]  $\implies$ 
  vConc (hConc t1 t2) (hConc t3 t4) = hConc (vConc t1 t3) (vConc t2 t4)
by (insert vConc-hConc-0 [of t1], auto)

```

5.3 Table Transposition

constdefs

```

  tTranspose :: ('h1, ('h2, 't) T) T  $\Rightarrow$  ('h2, ('h1, 't) T) T
  tTranspose == tFold addH2 vConc

```

lemma *tTranspose-addH[simp]*: *tTranspose (addH h t) = addH2 h t*
by (*simp add: tTranspose-def*)

lemma *tTranspose-hConc[simp]*: *tTranspose (hConc t1 t2) = vConc (tTranspose t1) (tTranspose t2)*
by (*simp add: tTranspose-def*)

lemma *tTranspose-hCons[simp]*: *tTranspose (hCons p t) = vCons p (tTranspose t)*
by (*simp add: tTranspose-def hCons vCons-def*)

lemma *regSkelOuter2-tTranspose-via-T-cons-induct*:

```

  ALL hs1 hs2 . regSkelOuter2 t = Some (hs1,hs2)  $\longrightarrow$ 
  regSkelOuter2 (tTranspose t) = Some (hs2,hs1)
apply (induct-tac t rule: T-cons-induct)
apply (simp add: regSkelOuter1-def regSkelOuter2-def)
apply (induct-tac t0 rule: T-cons-induct, simp, simp)
apply (subst tTranspose-hCons)
apply (intro strip)
apply (simp only: regSkelOuter2-hCons)
apply (case-tac regSkelOuter2 t2, simp, simp)
apply (split split-if-asm,simp-all, erule conjE)
apply (case-tac a, simp)
apply (frule-tac p=p in regSkelOuter2-Some-vCons)
apply (simp add: regSkelOuter2-def)
apply (drule regSkelStep-Some, simp, simp)
done

```

lemma *regSkelOuter2-tTranspose-via-T-induct*:
 $ALL\ hs1\ hs2 . regSkelOuter2\ t = Some\ (hs1,hs2) \longrightarrow$
 $regSkelOuter2\ (tTranspose\ t) = Some\ (hs2,hs1)$
apply (*induct-tac t rule: T-induct*)
apply (*simp add: regSkelOuter1-def regSkelOuter2-def*)
apply (*induct-tac t0 rule: T-induct, simp, simp*)
apply (*subst tTranspose-hConc*)
apply (*intro strip, simp only: regSkelOuter2-hConc*)
apply (*case-tac regSkelOuter2 t1, simp, simp*)
apply (*drule optThen-result-Some, erule exE, erule conjE*)
apply (*case-tac snd a = snd y, simp-all*)
apply (*erule conjE*)
apply (*drule-tac x=fst a in spec*)
apply (*drule-tac x=hs2 in spec*)
apply (*drule mp, fastsimp*)
apply (*drule-tac x=fst y in spec*)
apply (*drule-tac x=hs2 in spec*)
apply (*drule mp, fastsimp*)
apply (*rotate-tac -4, drule sym*)
apply (*simp add: regSkelOuter1-def*)
done

theorem *regSkelOuter2-tTranspose*:
 $regSkelOuter2\ t = Some\ (hs1,hs2) \implies$
 $regSkelOuter2\ (tTranspose\ t) = Some\ (hs2,hs1)$
by (*insert regSkelOuter2-tTranspose-via-T-induct [of t], auto*)

lemma *tFold-tFold-vConc-0*:
 $\llbracket\ assoc\ c1; assoc\ c2;$
 $\bigwedge\ x\ y\ z . a2\ x\ (c1\ y\ z) = c1\ (a2\ x\ y)\ (a2\ x\ z);$
 $\bigwedge\ x1\ x2\ y1\ y2 . c2\ (c1\ x1\ y1)\ (c1\ x2\ y2) = c1\ (c2\ x1\ x2)\ (c2\ y1\ y2)$
 $\rrbracket \implies$
 $ALL\ t2\ hs1a\ hs2a\ hs1b\ hs2b .$
 $regSkelOuter2\ t1 = Some\ (hs1a,hs2a) \longrightarrow$
 $regSkelOuter2\ t2 = Some\ (hs1b,hs2b) \longrightarrow$
 $hs1a = hs1b \longrightarrow$
 $tFold\ a2\ c2\ (tMap\ (tFold\ a1\ c1)\ (vConc\ t1\ t2)) =$
 $c1\ (tFold\ a2\ c2\ (tMap\ (tFold\ a1\ c1)\ t1))$
 $(tFold\ a2\ c2\ (tMap\ (tFold\ a1\ c1)\ t2))$
apply (*induct-tac t1 rule: T-cons-induct*)
apply (*rule T-cons-inductA*)

Case 1.1: $t1=addH, t2=addH$

apply (*intro strip*)
apply (*simp add: regSkelOuter1-def regSkelOuter2-def*)
apply (*erule conjE, erule conjE*)
apply (*subgoal-tac singleton h = singleton ha, drule singleton-inj, simp*)
apply *fastsimp*

Case 1.2: $t1=addH, t2=hCons$

apply (*intro strip*)
apply (*drule regSkelOuter2-hCons-eq-Some*)
apply (*erule exE, erule exE, erule exE, erule conjE, erule conjE, erule conjE*)
apply (*simp add: regSkelOuter2-def*)

Case 2: $t1=hCons$

apply (*rule T-cons-inductA*)

Case 2.1: $t1=hCons, t2=addH$

apply (*intro strip*)
apply (*drule regSkelOuter2-hCons-eq-Some*)
apply (*erule exE, erule exE, erule exE, erule conjE, erule conjE, erule conjE*)
apply (*drule regSkelOuter2-addH-eq-Some, erule conjE*)
apply *simp*

Case 2.2: $t1=hCons, t2=hCons$

apply (*intro strip*)
apply (*drule regSkelOuter2-hCons-eq-Some*)
apply (*erule exE, erule exE, erule exE, erule conjE, erule conjE, erule conjE*)
apply (*drule regSkelOuter2-hCons-eq-Some*)
apply (*erule exE, erule exE, erule exE, erule conjE, erule conjE, erule conjE*)
apply *simp*
apply (*drule cons1-inj, erule conjE, simp*)
done

lemma *tFold-tFold-vConc*:

\llbracket *assoc c1; assoc c2;*
 $\bigwedge x y z . a2 x (c1 y z) = c1 (a2 x y) (a2 x z);$
 $\bigwedge x1 x2 y1 y2 . c2 (c1 x1 y1) (c1 x2 y2) = c1 (c2 x1 x2) (c2 y1 y2);$
regSkelOuter2 t1 = Some (hs1a,hs2a);
regSkelOuter2 t2 = Some (hs1b,hs2b);
 $hs1a = hs1b$
 $\rrbracket \implies$
 $tFold a2 c2 (tMap (tFold a1 c1) (vConc t1 t2)) =$
 $c1 (tFold a2 c2 (tMap (tFold a1 c1) t1))$
 $(tFold a2 c2 (tMap (tFold a1 c1) t2))$

by (*insert tFold-tFold-vConc-0 [of c1 c2 a2 t1 a1], auto*)

lemma *tTranspose-tFold2-0*:

\llbracket *assoc c1; assoc c2;*
 $\bigwedge x y z . a1 x (a2 y z) = a2 y (a1 x z);$
 $\bigwedge x y z . a1 x (c2 y z) = c2 (a1 x y) (a1 x z);$
 $\bigwedge x y z . a2 x (c1 y z) = c1 (a2 x y) (a2 x z);$
 $\bigwedge x1 x2 y1 y2 . c2 (c1 x1 y1) (c1 x2 y2) = c1 (c2 x1 x2) (c2 y1 y2)$
 $\rrbracket \implies$
 $ALL hs1 hs2 . regSkelOuter2 t = Some (hs1,hs2) \longrightarrow$
 $tFold a2 c2 (tMap (tFold a1 c1) (tTranspose t)) =$
 $tFold a1 c1 (tMap (tFold a2 c2) t)$

apply (*induct-tac t rule: T-induct*)

```

apply (induct-tac t0 rule: T-induct)
apply (intro strip, simp)
apply (intro strip, simp add: regSkelOuter1-def regSkelOuter2-def)
apply (intro strip)
apply (drule regSkelOuter2-hConc-eq-Some)
apply (erule exE, erule exE, erule conjE, erule conjE)
apply (simp del: tFold-tMap)
apply (subst tFold-tFold-vConc)
  apply (assumption, assumption, simp, simp)
  apply (erule regSkelOuter2-tTranspose)
  apply (erule regSkelOuter2-tTranspose)
  apply simp
apply simp
done

```

theorem *tTranspose-tFold2*:

```

[[ assoc c1; assoc c2;
   $\wedge x y z . a1 x (a2 y z) = a2 y (a1 x z);$ 
   $\wedge x y z . a1 x (c2 y z) = c2 (a1 x y) (a1 x z);$ 
   $\wedge x y z . a2 x (c1 y z) = c1 (a2 x y) (a2 x z);$ 
   $\wedge x1 x2 y1 y2 . c2 (c1 x1 y1) (c1 x2 y2) = c1 (c2 x1 x2) (c2 y1 y2);$ 
  regSkelOuter2 t = Some (hs1,hs2)
]]  $\implies$ 
tFold a2 c2 (tMap (tFold a1 c1) (tTranspose t)) =
tFold a1 c1 (tMap (tFold a2 c2) t)
by (insert tTranspose-tFold2-0 [of c1 c2 a1 a2 t], auto)

```

lemma *tFold-tFold-vConc-gen-0*:

```

[[ assoc c1; assoc c2; assoc c3;
   $\wedge x y z . a3 x (c2 y z) = c1 (a1 x y) (a1 x z);$ 
   $\wedge x1 x2 y1 y2 . c3 (c1 x1 y1) (c1 x2 y2) = c1 (c1 x1 x2) (c1 y1 y2)$ 
]]  $\implies$ 
ALL t2 hs1a hs2a hs1b hs2b .
regSkelOuter2 t1 = Some (hs1a,hs2a)  $\longrightarrow$ 
regSkelOuter2 t2 = Some (hs1b,hs2b)  $\longrightarrow$ 
hs1a = hs1b  $\longrightarrow$ 
tFold a3 c3 (tMap (tFold a2 c2) (vConc t1 t2)) =
  c1 (tFold a1 c1 (tMap (tFold a2 c2) t1))
  (tFold a1 c1 (tMap (tFold a2 c2) t2))

```

```

apply (induct-tac t1 rule: T-cons-induct)
apply (rule T-cons-inductA)

```

Case 1.1: *t1=addH, t2=addH*

```

apply (intro strip)
apply (simp add: regSkelOuter1-def regSkelOuter2-def)
apply (erule conjE, erule conjE)
apply (subgoal-tac singleton h = singleton ha, drule singleton-inj, simp)
apply fastsimp

```

Case 1.2: *t1=addH, t2=hCons*

apply (*intro strip*)
apply (*drule regSkelOuter2-hCons-eq-Some*)
apply (*erule exE, erule exE, erule exE, erule conjE, erule conjE, erule conjE*)
apply (*simp add: regSkelOuter2-def*)

Case 2: $t1=hCons$

apply (*rule T-cons-inductA*)

Case 2.1: $t1=hCons, t2=addH$

apply (*intro strip*)
apply (*drule regSkelOuter2-hCons-eq-Some*)
apply (*erule exE, erule exE, erule exE, erule conjE, erule conjE, erule conjE*)
apply (*drule regSkelOuter2-addH-eq-Some, erule conjE*)
apply *simp*

Case 2.2: $t1=hCons, t2=hCons$

apply (*intro strip*)
apply (*drule regSkelOuter2-hCons-eq-Some*)
apply (*erule exE, erule exE, erule exE, erule conjE, erule conjE, erule conjE*)
apply (*drule regSkelOuter2-hCons-eq-Some*)
apply (*erule exE, erule exE, erule exE, erule conjE, erule conjE, erule conjE*)
apply *simp*
apply (*drule cons1-inj, erule conjE, simp*)
done

lemma *tFold-tFold-vConc-gen:*

\llbracket *assoc c1; assoc c2; assoc c3;*
 $\wedge x y z . a3 x (c2 y z) = c1 (a1 x y) (a1 x z);$
 $\wedge x1 x2 y1 y2 . c3 (c1 x1 y1) (c1 x2 y2) = c1 (c1 x1 x2) (c1 y1 y2);$
regSkelOuter2 t1 = Some (hs1a,hs2a);
regSkelOuter2 t2 = Some (hs1b,hs2b);
hs1a = hs1b
 $\rrbracket \implies$
 $tFold a3 c3 (tMap (tFold a2 c2) (vConc t1 t2)) =$
 $c1 (tFold a1 c1 (tMap (tFold a2 c2) t1))$
 $(tFold a1 c1 (tMap (tFold a2 c2) t2))$

by (*insert tFold-tFold-vConc-gen-0 [of c1 c2 c3 a3 a1 t1], auto*)

lemma *tFold-tFold-vConc-wrapped-0:*

\llbracket *assoc c1; assoc c2; assoc c3;*
 $\wedge h x y . w1 (a3 h (c2 x y)) = w2 (c1 (a1 h x) (a1 h y));$
 $\wedge x y . w1 (c3 x y) = c4 (w1 x) (w1 y);$
 $\wedge x1 x2 y1 y2 . c4 (w2 (c1 x1 x2)) (w2 (c1 y1 y2)) =$
 $w2 (c1 (c1 x1 y1) (c1 x2 y2))$
 $\rrbracket \implies$
 $ALL t2 hs1a hs2a hs1b hs2b .$
 $regSkelOuter2 t1 = Some (hs1a,hs2a) \longrightarrow$

```

regSkelOuter2 t2 = Some (hs1b,hs2b) →
hs1a = hs1b →
w1 (tFold a3 c3 (tMap (tFold a2 c2) (vConc t1 t2))) =
  w2 (c1 (tFold a1 c1 (tMap (tFold a2 c2) t1))
      (tFold a1 c1 (tMap (tFold a2 c2) t2)))
apply (induct-tac t1 rule: T-cons-induct)
apply (rule T-cons-inductA)

Case 1.1: t1=addH, t2=addH
apply (intro strip)
apply (simp add: regSkelOuter1-def regSkelOuter2-def)
apply (erule conjE, erule conjE)
apply (subgoal-tac singleton h = singleton ha, drule singleton-inj, simp)
apply fastsimp

Case 1.2: t1=addH, t2=hCons
apply (intro strip)
apply (drule regSkelOuter2-hCons-eq-Some)
apply (erule exE, erule exE, erule exE, erule conjE, erule conjE, erule conjE)
apply (simp add: regSkelOuter2-def)

Case 2: t1=hCons
apply (rule T-cons-inductA)

Case 2.1: t1=hCons, t2=addH
apply (intro strip)
apply (drule regSkelOuter2-hCons-eq-Some)
apply (erule exE, erule exE, erule exE, erule conjE, erule conjE, erule conjE)
apply (drule regSkelOuter2-addH-eq-Some, erule conjE)
apply simp

Case 2.2: t1=hCons, t2=hCons
apply (intro strip)
apply (drule regSkelOuter2-hCons-eq-Some)
apply (erule exE, erule exE, erule exE, erule conjE, erule conjE, erule conjE)
apply (drule regSkelOuter2-hCons-eq-Some)
apply (erule exE, erule exE, erule exE, erule conjE, erule conjE, erule conjE)
apply simp
apply (drule cons1-inj, erule conjE, simp)
done

lemma tFold-tFold-vConc-wrapped:
[[ assoc c1; assoc c2; assoc c3;
  ∧ h x y . w1 (a3 h (c2 x y)) = w2 (c1 (a1 h x) (a1 h y));
  ∧ x y . w1 (c3 x y) = c4 (w1 x) (w1 y);
  ∧ x1 x2 y1 y2 . c4 (w2 (c1 x1 x2)) (w2 (c1 y1 y2)) =
    w2 (c1 (c1 x1 y1) (c1 x2 y2));
  regSkelOuter2 t1 = Some (hs1,hs2a);
  regSkelOuter2 t2 = Some (hs1,hs2b)
]]

```

$\mathbb{I} \Rightarrow$
 $w1 (tFold\ a3\ c3\ (tMap\ (tFold\ a2\ c2)\ (vConc\ t1\ t2))) =$
 $w2 (c1 (tFold\ a1\ c1\ (tMap\ (tFold\ a2\ c2)\ t1))$
 $(tFold\ a1\ c1\ (tMap\ (tFold\ a2\ c2)\ t2)))$
by (*insert tFold-tFold-vConc-wrapped-0 [of c1 c2 c3 w1 a3 w2 a1 c4 t1], auto*)

lemma *tTranspose-tFold2-gen-0:*

\mathbb{I} *assoc c1; assoc c2; assoc c3; assoc c4; assoc c5;*
 $\wedge x\ y\ z . w1 (a1\ x\ (a2\ y\ z)) = w2 (a3\ y\ (a4\ x\ z));$
 $\wedge x\ y . w1 (c1\ x\ y) = c5 (w1\ x) (w1\ y);$
 $\wedge x\ y . w2 (c3\ x\ y) = c5 (w2\ x) (w2\ y);$
 $\wedge x\ y\ z . w1 (a1\ x\ (c2\ y\ z)) = c5 (w1 (a1\ x\ y)) (w1 (a1\ x\ z));$
 $\wedge h\ x\ y . w2 (a3\ h\ (c4\ x\ y)) = c5 (w2 (a3\ h\ x)) (w2 (a3\ h\ y));$
 $\wedge x1\ x2\ y1\ y2 . c5 (c5\ x1\ x2) (c5\ y1\ y2) = c5 (c5\ x1\ y1) (c5\ x2\ y2)$
 $(*$
 $\wedge x\ y\ z . a2\ x (c1\ y\ z) = c1 (a2\ x\ y) (a2\ x\ z);$
 $*)$
 $\mathbb{I} \Rightarrow$
 $ALL\ hs1\ hs2 . regSkelOuter2\ t = Some\ (hs1,hs2) \longrightarrow$
 $w2 (tFold\ a3\ c3\ (tMap\ (tFold\ a4\ c4)\ (tTranspose\ t))) =$
 $w1 (tFold\ a1\ c1\ (tMap\ (tFold\ a2\ c2)\ t))$
apply (*induct-tac t rule: T-induct*)
apply (*induct-tac t0 rule: T-induct*)
apply (*intro strip, simp*)
apply (*intro strip, simp add: regSkelOuter1-def regSkelOuter2-def*)
apply (*intro strip*)
apply (*drule regSkelOuter2-hConc-eq-Some*)
apply (*erule exE, erule exE, erule conjE, erule conjE*)
apply (*simp del: tFold-tMap*)
apply (*subst tFold-tFold-vConc-wrapped [of c3 c4 c3 w2 a3 w2 a3 c5]*)
apply (*assumption, assumption, assumption*)
apply *simp*
apply *simp*
apply *simp*
apply (*erule regSkelOuter2-tTranspose*)
apply (*erule regSkelOuter2-tTranspose*)
apply *simp*
done

lemma *tTranspose-tFold2-gen:*

\mathbb{I} *assoc c1; assoc c2; assoc c3; assoc c4; assoc c5;*
 $\wedge x\ y\ z . w1 (a1\ x\ (a2\ y\ z)) = w2 (a3\ y\ (a4\ x\ z));$
 $\wedge x\ y . w1 (c1\ x\ y) = c5 (w1\ x) (w1\ y);$
 $\wedge x\ y . w2 (c3\ x\ y) = c5 (w2\ x) (w2\ y);$
 $\wedge x\ y\ z . w1 (a1\ x\ (c2\ y\ z)) = c5 (w1 (a1\ x\ y)) (w1 (a1\ x\ z));$
 $\wedge h\ x\ y . w2 (a3\ h\ (c4\ x\ y)) = c5 (w2 (a3\ h\ x)) (w2 (a3\ h\ y));$
 $\wedge x1\ x2\ y1\ y2 . c5 (c5\ x1\ x2) (c5\ y1\ y2) = c5 (c5\ x1\ y1) (c5\ x2\ y2);$
 $regSkelOuter2\ t = Some\ (hs1,hs2)$
 $\mathbb{I} \Rightarrow$

$w2 (tFold\ a3\ c3\ (tMap\ (tFold\ a4\ c4)\ (tTranspose\ t))) =$
 $w1 (tFold\ a1\ c1\ (tMap\ (tFold\ a2\ c2)\ t))$
by (*insert tTranspose-tFold2-gen-0 [of c1 c2 c3 c4 c5 w1 a1 a2 w2 a3 a4 t], auto*)

end

6 Inversion of Normal Tables

theory *Inversion* = *Tables2*:

6.1 One-Dimensional Inversion

consts

inverse1 :: ('a,('b,'c) T) T \Rightarrow ('b,('a,'c) T) T

defs

inverse1-def: *inverse1* == *tFold addH2 hConc*

lemma *inverse1-addH[simp]*:

inverse1 (*addH* h t) = *addH2* h t

by (*unfold inverse1-def, simp*)

lemma *inverse1-hConc[simp]*:

inverse1 (*hConc* t1 t2) = *hConc* (*inverse1* t1) (*inverse1* t2)

by (*unfold inverse1-def, simp*)

lemma *tFold-tFold-inverse1*:

\llbracket *assoc* c2; *assoc* c4;

\wedge h2 h1 t . a2 h2 (a1 h1 t) = a3 h1 (a4 h2 t);

\wedge h t u . a3 h (c4 t u) = c2 (a3 h t) (a3 h u)

$\rrbracket \Rightarrow$

tFold a2 c2 (*tMap* (*tFold* a1 c1) (*inverse1* t)) =

tFold a3 c2 (*tMap* (*tFold* a4 c4) t)

apply (*induct-tac* t rule: *T-induct*)

apply (*induct-tac* t0 rule: *T-induct, simp*)

apply (*induct-tac* t1 rule: *T-induct, simp*)

apply *simp*

apply (*subgoal-tac* *tFold* (λ ha t0. a3 h (a4 ha t0)) c2 t1a = a3 h (*tFold* a4 c4 t1a), *simp*)

apply (*induct-tac* t1a rule: *T-induct, simp, simp*)

original induction, second case:

apply (*simp del: tFold-tMap*)

done

lemma *tFold-tFold0-inverse1*:

\llbracket *assoc* c2; *assoc* c4;

\wedge h t u . a2 (c4 t u) h = c2 (a2 t h) (a2 u h)

$\square \implies$
 $tFold\ a2\ c2\ (tMap\ (tFold0\ c1)\ (inverse1\ t)) =$
 $tFold\ (flip\ a2)\ c2\ (tMap\ (tFold0\ c4)\ t)$
by (*unfold tFold0-def, rule tFold-tFold-inverse1, simp-all*)

6.2 spread1

consts *spread1* :: 'a neList \Rightarrow 'b \Rightarrow ('a, 'b) T

defs *spread1-def*: *spread1* *hs* *t* == *foldr1* *hConc* (*map1* ($\lambda h.$ *addH* *h* *t*) *hs*)

lemma *regSkelOuter2-spread1*[*simp*]:

regSkelOuter2 (*spread1* *hs* *t*) = *Some* (*hs*, *headers* *t*)

apply (*subst* *neList-append-induct* [*of* % *hs* . *regSkelOuter2* (*spread1* *hs* *t*) = *Some* (*hs*, *headers* *t*)], *simp-all*)

WHy does standard induction not work here?!

apply (*unfold* *spread1-def*, *simp-all*)

done

lemma *spread1-singleton*[*simp*]: *spread1* (*singleton* *h*) *t* = *addH* *h* *t*

by (*simp* *add*: *spread1-def*)

lemma *spread1-append*[*simp*]:

spread1 (*append* *hs1* *hs2*) *t* = *hConc* (*spread1* *hs1* *t*) (*spread1* *hs2* *t*)

by (*simp* *add*: *spread1-def*)

lemma *spread1-tFold-append*:

spread1 (*tFold* *f* *append* *t1*) *t* = *tFold* (% *h0* *t0* . *spread1* (*f* *h0* *t0*) *t*) *hConc* *t1*

by (*induct-tac* *t1* *rule*: *T-induct*, *simp-all*)

lemma *spread1-hConc-0*:

ALL *t1* . *spread1* *hs* (*hConc* *t1* *t2*) = *vConc* (*spread1* *hs* *t1*) (*spread1* *hs* *t2*)

apply (*induct-tac* *hs* *rule*: *neList-append-induct*, *simp*)

apply *simp*

apply (*rule* *allI*)

apply (*subst* *vConc-hConc*)

apply *simp*

apply (*rule* *conjI*, *simp*, *simp*)

apply *simp*

apply *simp*

apply *simp*

apply *simp*

done

lemma *spread1-hConc*:

spread1 *hs* (*hConc* *t1* *t2*) = *vConc* (*spread1* *hs* *t1*) (*spread1* *hs* *t2*)

by (*insert* *spread1-hConc-0* [*of* *hs* *t2*], *auto*)

lemma *tMap-f-spread1*:

$tMap\ f\ (spread1\ hs\ t) = spread1\ hs\ (f\ t)$
by (*induct-tac hs rule: neList-append-induct, simp-all*)

lemma *tFold-spread1*:
 $assoc\ c \implies tFold\ a\ c\ (spread1\ hs\ t) = foldr1\ c\ (map1\ (\% h . a\ h\ t)\ hs)$
by (*induct-tac hs rule: neList-append-induct, simp-all*)

6.3 spread2

This was used in the previous definition of *dconc*.

consts

spread2 :: 'a neList \Rightarrow ('b,'c) T \Rightarrow ('b,('a,'c) T) T

defs

spread2-def: $spread2\ hs\ t == foldr1\ vConc\ (map1\ (\% h . addH2\ h\ t)\ hs)$

lemma *spread2-singleton[simp]*:
 $spread2\ (singleton\ h)\ t = addH2\ h\ t$
by (*unfold spread2-def, simp*)

lemma *spread2-append[simp]*:
 $spread2\ (append\ hs1\ hs2)\ t = vConc\ (spread2\ hs1\ t)\ (spread2\ hs2\ t)$
by (*unfold spread2-def, simp*)

lemma *spread2-cons1[simp]*:
 $spread2\ (cons1\ h\ hs)\ t = vCons\ (h,t)\ (spread2\ hs\ t)$
by (*simp add: cons1 vCons*)

lemma *headers-spread2[simp]*:
 $headers\ (spread2\ hs\ t) = headers\ t$
by (*induct-tac hs rule: neList-cons-induct, simp-all*)

lemma *regSkelOuter2-spread2-0*:
 $ALL\ hs2 . regSkelOuter2\ (spread2\ hs2\ t) = Some\ (headers\ t,hs2)$
apply (*induct-tac t rule: T-cons-induct, simp*)
apply (*rule allI, induct-tac hs2 rule: neList-append-induct*)
apply *simp*
apply *simp*
apply (*rule allI*)
apply (*erule mp1*)
apply (*rule-tac x=t2 in spec*)
apply (*induct-tac hs2 rule: neList-cons-induct*)
apply (*intro strip, simp*)

$hs2 = cons1\ x\ xs$

apply (*intro strip, simp*)
apply (*drule-tac x=xa in spec*)
apply *simp*
apply (*subst regSkelOuter2-vCons, simp*)

apply (*simp add: cons1*)
done

lemma *regSkelOuter2-spread2[simp]*:
 $\text{regSkelOuter2 } (\text{spread2 } hs2 \ t) = \text{Some } (\text{headers } t, hs2)$
by (*rule regSkelOuter2-spread2-0 [THEN spec]*)

lemma *spread2-hConc-0*:
 $\text{ALL } t1 \ t2 \ . \ \text{spread2 } hs2 \ (\text{hConc } t1 \ t2) = \text{hConc } (\text{spread2 } hs2 \ t1) \ (\text{spread2 } hs2 \ t2)$
apply (*induct-tac hs2 rule: neList-append-induct*)
apply *simp*
apply *simp*
apply (*intro strip*)
apply (*subst vConc-hConc*)
apply *simp*
apply (*rule conjI*)
apply *simp*
apply *simp*
apply *simp*
apply *simp*
apply *simp*
apply *simp*
done

lemma *spread2-hConc[simp]*:
 $\text{spread2 } hs2 \ (\text{hConc } t1 \ t2) = \text{hConc } (\text{spread2 } hs2 \ t1) \ (\text{spread2 } hs2 \ t2)$
by (*insert spread2-hConc-0, auto*)

lemma *spread2-addH*:
 $\text{spread2 } hs \ (\text{addH } h \ t) = \text{addH } h \ (\text{foldr1 } \text{hConc } (\text{map1 } (\% \ h2 \ . \ \text{addH } h2 \ t) \ hs))$
by (*induct-tac hs rule: neList-append-induct, simp-all*)

lemma *spread2-tFold-hConc*:
 $\text{spread2 } hs2 \ (\text{tFold } a \ \text{hConc } \ t) = \text{tFold } (\% \ h \ t \ . \ \text{spread2 } hs2 \ (a \ h \ t)) \ \text{hConc } \ t$
by (*induct-tac t rule: T-induct, simp-all*)

lemma *tMap-tFold-spread2*:
 $\text{assoc } c \implies$
 $\text{tMap } (\text{tFold } a \ c) \ (\text{spread2 } hs \ t) =$
 $\text{tMap } (\% \ t0 \ . \ \text{foldr1 } c \ (\text{map1 } (\lambda h2. \ a \ h2 \ t0) \ hs)) \ t$
by (*induct-tac t rule: T-induct, simp-all add: spread2-addH map1-contract-comp tFold-foldr1-hConc o-def*)

6.4 delH1

consts *delH1* :: ('a,'b) T ⇒ 'b
defs *delH1-def*: *delH1* == *tFold* (% h t . t) *const*

lemma *delH1-hConc*[simp]:
 $delH1 (hConc t1 t2) = delH1 t1$
by (*simp add: delH1-def*)

lemma *delH1-addH*[simp]:
 $delH1 (addH h t) = t$
by (*simp add: delH1-def*)

lemma *delH1-hCons*[simp]:
 $delH1 (hCons (h,t0) t) = t0$
by (*simp add: hCons*)

lemma *delH1-tFold-hConc*[simp]:
 $delH1 (tFold f hConc t) = delH1 (uncurry f (hHead t))$
by (*induct-tac t rule: T-induct, simp-all*)

lemma *headers-delH1--reg2-0*:
 $ALL hs1 hs2 . regSkelOuter2 t1 = Some (hs1, hs2) \longrightarrow headers (delH1 t1) = hs2$
apply (*induct-tac t1 rule: T-induct, simp-all*)
apply (*intro strip*)
apply (*drule optThen-result-Some, erule exE, erule conjE*)
apply (*drule optThen-result-Some, erule exE, erule conjE*)
apply (*split split-if-asm, simp-all*)
apply (*case-tac y, case-tac ya, simp*)
done

lemma *headers-delH1--reg2*[simp]:
 $regSkelOuter2 t1 = Some (hs1, hs2) \Longrightarrow headers (delH1 t1) = hs2$
by (*insert headers-delH1--reg2-0, auto*)

lemma *delH1-addH2*[simp]: $delH1 (addH2 h t) = addH h (delH1 t)$
by (*induct-tac t rule: T-induct, simp-all*)

lemma *delH1-vConc-0*:
 $ALL hs1a hs1b hs2a hs2b .$
 $regSkelOuter2 t1 = Some (hs1a, hs1b) \longrightarrow$
 $regSkelOuter2 t2 = Some (hs2a, hs2b) \longrightarrow$
 $delH1 (vConc t1 t2) = hConc (delH1 t1) (delH1 t2)$
apply (*induct-tac t1 rule: T-cons-induct, simp-all*)
apply (*induct-tac t2 rule: T-cons-induct, simp-all*)
apply (*intro strip, erule exE, erule exE*)
apply (*drule optThen-result-Some, erule exE, erule conjE*)
apply (*split split-if-asm, simp-all*)
apply (*case-tac p, case-tac y, simp*)
apply (*induct-tac t2 rule: T-cons-induct, simp-all*)
apply (*intro strip, erule exE, erule exE*)
apply (*drule optThen-result-Some, erule exE, erule conjE*)
apply (*split split-if-asm, simp-all*)

apply (*case-tac p, case-tac y, simp*)
apply (*intro strip, erule exE, erule exE, erule exE, erule exE*)
apply *simp*
apply (*case-tac p, case-tac pa, simp*)
done

lemma *delH1-vConc*:

\llbracket *regSkelOuter2 t1 = Some (hs1a, hs1b);*
 \llbracket *regSkelOuter2 t2 = Some (hs2a, hs2b)*
 $\rrbracket \implies$
 $delH1 (vConc t1 t2) = hConc (delH1 t1) (delH1 t2)$
by (*insert delH1-vConc-0, auto*)

lemma *delH1-tMap-const[simp]*: $delH1 (tMap (const t2) t1) = t2$
by (*induct-tac t1 rule: T-induct, simp-all*)

lemma *delH1-tMap-CONST[simp]*: $delH1 (tMap (\% x . t2) t1) = t2$
by (*induct-tac t1 rule: T-induct, simp-all*)

6.5 delH2

consts *delH2* :: $(a, (b, c) T) T \Rightarrow (a, c) T$

defs *delH2-def*: $delH2 == tMap delH1$

lemma *delH2-hConc[simp]*:

$delH2 (hConc t1 t2) = hConc (delH2 t1) (delH2 t2)$
by (*simp add: delH2-def*)

lemma *delH2-addH[simp]*:

$delH2 (addH h t) = addH h (delH1 t)$
by (*simp add: delH2-def*)

lemma *delH2-hCons[simp]*:

$delH2 (hCons (h, t0) t) = hCons (h, delH1 t0) (delH2 t)$
by (*simp add: hCons*)

lemma *delH2-addH2[simp]*:

$delH2 (addH2 h t) = t$
by (*induct-tac t rule: T-induct, simp-all*)

lemma *headers-delH2[simp]*: $headers (delH2 t) = headers t$

by (*induct-tac t rule: T-induct, simp-all*)

lemma *delH2-vConc-0*:

ALL hs1 hs2a t2 hs2b .
 $regSkelOuter2 t1 = Some (hs1, hs2a) \longrightarrow$
 $regSkelOuter2 t2 = Some (hs1, hs2b) \longrightarrow$
 $delH2 (vConc t1 t2) = delH2 t1$

```

apply (induct-tac t1 rule: T-cons-induct, simp-all)
apply (rule allI)
apply (rule impI)
apply (rule allI)
apply (erule exE)+
apply (drule optThen-result-Some, erule exE, erule conjE, simp)
apply (split split-if-asm, simp-all)
apply (erule conjE, case-tac y, simp)
apply (induct-tac t2a rule: T-cons-induct, simp-all)
apply (rule impI, rotate-tac -1, drule sym, simp)
apply (intro strip)
apply (erule exE)+
apply (drule optThen-result-Some, erule exE, erule conjE, simp)
apply (split split-if-asm, simp-all)
apply (erule conjE, case-tac p, case-tac pa, case-tac ya, simp)
apply (drule-tac x=t2b in spec, simp)
apply (rotate-tac 2, drule sym, simp, drule cons1-inj, simp)
done

```

```

lemma delH2-vConc[simp]:
  [| regSkelOuter2 t1 = Some (hs1, hs2a);
    regSkelOuter2 t2 = Some (hs1, hs2b) |] ==>
    delH2 (vConc t1 t2) = delH2 t1
by (insert delH2-vConc-0 [of t1], auto)

```

```

lemma delH2-spread2[simp]:
  delH2 (spread2 hs t) = t
apply (induct-tac t rule: T-induct, simp-all)
apply (induct-tac hs rule: neList-append-induct, simp-all)
apply (subst delH2-vConc, simp-all)
apply (rule conjI, simp-all)
apply (rule conjI, simp-all)
done

```

```

lemma delH2-tMap-const:
  delH2 (tMap (const t2) t1) = tMap (const (delH1 t2)) t1
by (induct-tac t1 rule: T-induct, simp-all)

```

```

lemma delH2-tMap-CONST:
  delH2 (tMap (% x . t2) t1) = tMap (const (delH1 t2)) t1
by (induct-tac t1 rule: T-induct, simp-all)

```

6.6 Third Dimension Operators

```

consts addH3 :: 'c => ('a, ('b, 'd) T) T => ('a, ('b, ('c, 'd) T) T) T

```

```

defs addH3-def: addH3 == tMap o addH2

```

```

lemma delH1-addH3[simp]: delH1 (addH3 h t) = addH2 h (delH1 t)

```

by (*unfold addH3-def*, *induct-tac t rule: T-induct*, *simp-all*)

consts *delH3* :: ('a, ('b, ('c, 'd) T) T) T ⇒ ('a, ('b, 'd) T) T

defs *delH3-def*: *delH3* == *tMap delH2*

6.7 Slimming Operators

These operators slim down a selected dimension in the regular table skeleton to a singleton containing a supplied entry *u*.

consts *slimH1* :: 'a ⇒ ('c, 'b) T ⇒ ('a, 'b) T

defs *slimH1-def*[*simp*]: *slimH1 h1 t* == *addH h1 (delH1 t)*

lemma *tFold-slimH1*:

assoc c ⇒ *tFold a c (slimH1 u t) = tFold (% h0 t0 . a u t0) const t*

apply (*induct-tac t rule: T-induct*, *simp-all del: const*)

apply (*subst const*)

apply (*simp (no-asm-simp) del: const*)

done

lemma *headers-slimH1*[*simp*]: *headers (slimH1 u t) = singleton u*

by *simp*

consts *slimH2* :: 'b ⇒ ('a, ('d, 'c) T) T ⇒ ('a, ('b, 'c) T) T

defs *slimH2-def*[*simp*]: *slimH2 h2 t* == *addH2 h2 (delH2 t)*

lemma *slimH2-tMap*: *slimH2 h2 = tMap (slimH1 h2)*

by (*rule ext*, *induct-tac x rule: T-induct*, *simp-all*)

lemma *slimH2-slimH2*[*simp*]: *slimH2 u (slimH2 v t) = slimH2 u t*

by *simp*

lemma *slimH2-addH*[*simp*]: *slimH2 u (addH h t) = addH h (slimH1 u t)*

by *simp*

lemma *slimH2-hConc*[*simp*]:

slimH2 u (hConc t1 t2) = hConc (slimH2 u t1) (slimH2 u t2)

by *simp*

lemma *slimH2-as-tFold*[*simp*]:

tFold (λh t. addH h (slimH1 u t)) hConc t = slimH2 u t

by (*induct-tac t rule: T-induct*, *simp-all*)

lemma *slimH2-addH2*[*simp*]: *slimH2 u (addH2 h2 t) = addH2 u t*

by *simp*

lemma *headers-slimH2*[*simp*]: *headers (slimH2 u t) = headers t*

by (*induct-tac t rule: T-induct, simp-all*)

lemma *regSkelOuter2-slimH2[simp]*:

regSkelOuter2 (slimH2 u t) = Some (headers t, singleton u)

by (*induct-tac t rule: T-induct, simp-all*)

lemma *slimH2-tFold-hConc[simp]*:

slimH2 u (tFold f hConc t) = tFold (% h t . slimH2 u (f h t)) hConc t

by (*induct-tac t rule: T-induct, simp-all del: slimH2-def*)

lemma *slimH2-tFold-addH2-hConc[simp]*:

slimH2 u (tFold (% h . addH2 v) hConc t) = tFold (% h . addH2 u) hConc t

by (*induct-tac t rule: T-induct, simp-all del: slimH2-def*)

lemma *del1-slimH2-isConst[simp]*:

delH1 (slimH2 u (t1 :: ('a, ('b, unit) T) T)) = delH1 (slimH2 u t2)

by *simp*

lemma *tFold-slimH2*:

assoc c \implies tFold a c (slimH2 u t) = tFold (% h0 t0 . a h0 (slimH1 u t0)) c t

by (*induct-tac t rule: T-induct, simp-all*)

consts *slimH3* :: 'c \Rightarrow ('a, ('b, ('e, 'd) T) T) T \Rightarrow ('a, ('b, ('c, 'd) T) T) T

defs *slimH3-def*: *slimH3 == tMap \circ slimH2*

lemma *slimH3-expand*: *slimH3 u t = addH3 u (delH3 t)*

by (*induct-tac t rule: T-induct, simp-all add: delH3-def addH3-def slimH3-def*)

lemma *slimH3-addH[simp]*: *slimH3 u (addH h t) = addH h (slimH2 u t)*

by (*simp add: slimH3-def del: slimH2-def*)

lemma *slimH3-hConc[simp]*: *slimH3 u (hConc t1 t2) = hConc (slimH3 u t1) (slimH3 u t2)*

by (*simp add: slimH3-def del: slimH2-def*)

lemma *spread1-slimH2*: *spread1 hs (slimH2 u t) = slimH3 u (spread1 hs t)*

by (*induct-tac hs rule: neList-append-induct, simp-all del: slimH2-def slimH3-def*)

lemma *tFold-slimH3*:

assoc c \implies tFold a c (slimH3 u t) = tFold (% h0 t0 . a h0 (slimH2 u t0)) c t

by (*induct-tac t rule: T-induct, simp-all del: slimH2-def slimH3-def*)

lemma *tMap-f-slimH3*:

tMap f (slimH3 u t) = tMap (% t0 . f (slimH2 u t0)) t

apply (*unfold slimH3-def*)

apply (*unfold comp-def*)

apply (*subst tMap-tMap*)

apply (*unfold comp-def*)

apply (*rule refl*)
done

6.8 Right-Updating Horizontal Concatenation

consts $hConcU :: ('c \Rightarrow ('a, 'b) T) \Rightarrow ('a, 'b) T \Rightarrow 'c \Rightarrow ('a, 'b) T$

defs $hConcU\text{-def}[simp]: hConcU\ u\ t1\ t2 == hConc\ t1\ (u\ t2)$

lemma $hConcU\text{-assoc}[simp]:$

$\llbracket \bigwedge t . u\ (u\ t) = u\ t;$
 $\bigwedge t1\ t2 . u\ (hConc\ t1\ t2) = hConc\ (u\ t1)\ (u\ t2)$
 $\rrbracket \Longrightarrow assoc\ (hConcU\ u)$

by (*rule assoc-intro, simp*)

lemma $U\text{-}hConcU[simp]:$

$\llbracket \bigwedge t . u\ (u\ t) = u\ t;$
 $\bigwedge t1\ t2 . u\ (hConc\ t1\ t2) = hConc\ (u\ t1)\ (u\ t2)$
 $\rrbracket \Longrightarrow u\ (hConcU\ u\ t1\ t2) = hConc\ (u\ t1)\ (u\ t2)$

by *simp*

lemma $U\text{-}tFold\text{-}hConcU[simp]:$

$\llbracket \bigwedge t . u\ (u\ t) = u\ t;$
 $\bigwedge t1\ t2 . u\ (hConc\ t1\ t2) = hConc\ (u\ t1)\ (u\ t2)$
 $\rrbracket \Longrightarrow u\ (tFold\ f\ (hConcU\ u)\ t) = tFold\ (\% h0\ t0 . u\ (f\ h0\ t0))\ hConc\ t$

by (*induct-tac t rule: T-induct, simp-all del: hConcU-def*)

consts $hConcSH2 :: 'b \Rightarrow ('a, ('b, 'c) T) T \Rightarrow ('a, ('d, 'c) T) T \Rightarrow ('a, ('b, 'c) T) T$

defs $hConcSH2\text{-def}: hConcSH2\ u == hConcU\ (slimH2\ u)$

lemma $hConcSH2\text{-assoc}[simp]: assoc\ (hConcSH2\ u)$

by (*rule assoc-intro, simp add: hConcSH2-def*)

lemma $slimH2\text{-}hConcSH2[simp]:$

$slimH2\ u\ (hConcSH2\ u\ t1\ t2) = hConc\ (slimH2\ u\ t1)\ (slimH2\ u\ t2)$

by (*unfold hConcSH2-def, simp del: slimH2-def*)

lemma $slimH2\text{-}tFold\text{-}hConcSH2[simp]:$

$slimH2\ u\ (tFold\ f\ (hConcSH2\ u)\ t) = tFold\ (\% h\ t . slimH2\ u\ (f\ h\ t))\ hConc\ t$

by (*induct-tac t rule: T-induct, simp-all del: slimH2-def*)

6.9 Diagonal Table Concatenation

Diagonal concatenation allocates the two last arguments as blocks on the main diagonal and fills the remainder with “empty” tables created using the first argument.

Since we shall use this as combinator of a table fold, we need its associativity, but that holds only on tables that are regular in their outer two dimensions.

For this reason we define diagonal concatenation on the subtype $regT2$.

```
consts dConc :: (('b,'c) T ⇒ ('b,'c) T) ⇒
           ('a,'b,'c) regT2 ⇒
           ('a,'b,'c) regT2 ⇒
           ('a,'b,'c) regT2
```

```
defs dConc-def: dConc h t1 t2 ==
  let t1' = Rep-regT2 t1;
      t2' = Rep-regT2 t2
  in Abs-regT2
     (hConc (vConc t1' (tMap (const (h (delH1 t2')))) t1')
            (vConc (tMap (const (h (delH1 t1')))) t2') t2'
          )
    )
```

```
lemma dConc[simp]:
  [[ regSkelOuter2 t1 = Some (cs1, hs1);
    regSkelOuter2 t2 = Some (cs2, hs2)
  ]] ⇒
  dConc h (Abs-regT2 t1) (Abs-regT2 t2)
  = Abs-regT2
    (hConc (vConc t1 (tMap (const (h (delH1 t2)))) t1)
           (vConc (tMap (const (h (delH1 t1)))) t2) t2
    )
```

```
apply (unfold dConc-def, simp only: Let-def regT2-def)
apply (subst Abs-regT2-inverse, simp add: regT2-def, fast)+
apply simp
done
```

```
lemma dConc1:
  [[ regSkelOuter2 (Rep-regT2 t1) = Some (cs1, hs1);
    regSkelOuter2 (Rep-regT2 t2) = Some (cs2, hs2)
  ]] ⇒
  dConc h t1 t2
  = Abs-regT2
    (hConc (vConc (Rep-regT2 t1) (tMap (const (h (delH1 (Rep-regT2 t2))))
      (Rep-regT2 t1)))
           (vConc (tMap (const (h (delH1 (Rep-regT2 t1))))
      (Rep-regT2 t2)))
      (Rep-regT2 t2))
    )
apply (frule-tac h=h and t2.0=Rep-regT2 t2 in dConc, assumption)
apply (simp only: Rep-regT2-inverse)
done
```

```
lemma regSkelOuter2-dConcDef[simp]:
  [[ ∧ t . headers (h t) = headers t;
    regSkelOuter2 t1 = Some (cs1, hs1);
    regSkelOuter2 t2 = Some (cs2, hs2)
  ]] ⇒
```



```

regSkelOuter2
  (hConc (vConc t1 (tMap (const (h (delH1 t2))) t1))
    (vConc (tMap (const (h (delH1 t1))) t2) t2))
= Some (append cs1 cs2, append hs1 hs2)
apply (rotate-tac -1, frule regSkelOuter2-eq-Some)
apply (rotate-tac -2, frule regSkelOuter2-eq-Some)
apply (cut-tac t1.0=t1 and t2.0=tMap (const (h (delH1 t2))) t1 in regSkelOuter2-vConc)
  apply assumption
  apply (simp (no-asm-simp))
  apply (rule conjI)
  apply (simp (no-asm-simp))
  apply (simp (no-asm-simp))
  apply (simp (no-asm-simp))
apply (rotate-tac -1, frule regSkelOuter2-eq-Some)
apply (cut-tac t1.0=tMap (const (h (delH1 t1))) t2 and t2.0=t2 in regSkelOuter2-vConc)
  apply (simp (no-asm-simp))
  apply (rule conjI)
  apply (simp (no-asm-simp))
  apply (simp (no-asm-simp))
  apply assumption
  apply (simp (no-asm-simp))
apply (rotate-tac -1, frule regSkelOuter2-eq-Some)
apply (simp (no-asm-simp))
apply (subst headers-delH1--reg2, assumption)
apply (subst headers-delH1--reg2, assumption)
apply (simp (no-asm-simp))
done

lemma regSkelOuter2-dConc[simp]:
  [ [  $\wedge t . \text{headers } (h t) = \text{headers } t;$ 
    regSkelOuter2 t1 = Some (cs1, hs1);
    regSkelOuter2 t2 = Some (cs2, hs2)
  ]  $\implies$ 
  regSkelOuter2 (Rep-regT2 (dConc h (Abs-regT2 t1) (Abs-regT2 t2)))
  = Some (append cs1 cs2, append hs1 hs2)
apply (subst dConc)
  apply assumption
  apply assumption
apply (subst Abs-regT2-inverse, rule regT2)
apply (subst regSkelOuter2-dConcDef)
  apply simp
  apply assumption
  apply assumption
  apply simp
  apply (rule conjI, simp, simp)
apply (subst regSkelOuter2-dConcDef)
  apply simp
  apply assumption
  apply assumption

```

apply *simp*
done

lemma *dConc2[*simp*]*:

$\llbracket \bigwedge t . \text{headers } (h \ t) = \text{headers } t;$
 $\text{regSkelOuter2 } (\text{Rep-regT2 } t1) = \text{Some } (cs1, hs1);$
 $\text{regSkelOuter2 } (\text{Rep-regT2 } t2) = \text{Some } (cs2, hs2)$
 $\rrbracket \implies$
 $\text{Rep-regT2 } (dConc \ h \ t1 \ t2)$
 $= hConc \ (vConc \ (\text{Rep-regT2 } t1) \ (tMap \ (\text{const } (h \ (\text{delH1 } (\text{Rep-regT2 } t2))))$
 $(\text{Rep-regT2 } t1)))$
 $\quad (vConc \ (tMap \ (\text{const } (h \ (\text{delH1 } (\text{Rep-regT2 } t1)))) \ (\text{Rep-regT2 } t2))$
 $(\text{Rep-regT2 } t2))$
apply (*subst dConc1*)
apply *assumption+*
apply (*subst Abs-regT2-inverse, rule regT2*)
apply (*subst regSkelOuter2-dConcDef*)
apply *simp*
apply *assumption*
apply *assumption*
apply *simp*
apply (*rule conjI, simp, simp*)
apply (*simp del: const*)
done

lemma *regSkelOuter2-dConc1[*simp*]*:

$\llbracket \bigwedge t . \text{headers } (h \ t) = \text{headers } t;$
 $\text{regSkelOuter2 } (\text{Rep-regT2 } t1) = \text{Some } (cs1, hs1);$
 $\text{regSkelOuter2 } (\text{Rep-regT2 } t2) = \text{Some } (cs2, hs2)$
 $\rrbracket \implies$
 $\text{regSkelOuter2 } (\text{Rep-regT2 } (dConc \ h \ t1 \ t2))$
 $= \text{Some } (\text{append } cs1 \ cs2, \text{append } hs1 \ hs2)$
apply (*subst dConc2*)
apply *simp*
apply *assumption*
apply *assumption*
apply (*subst regSkelOuter2-dConcDef*)
apply *simp*
apply *assumption*
apply *assumption*
apply *simp*
done

lemma *headers-dConc[*simp*]*:

$\llbracket \bigwedge t . \text{headers } (h \ t) = \text{headers } t$
 $\rrbracket \implies$
 $\text{headers } (\text{Rep-regT2 } (dConc \ h \ t1 \ t2)) = \text{append } (\text{headers } (\text{Rep-regT2 } t1)) \ (\text{headers } (\text{Rep-regT2 } t2))$
apply (*cut-tac x=t1 in Rep-regT2*)

```

apply (cut-tac x=t2 in Rep-regT2)
apply (simp add: regT2-def)
apply (drule regularOuter2, erule exE, erule exE)
apply (drule regularOuter2, erule exE, erule exE)
apply (subst dConc2)
  apply simp
  apply assumption
  apply assumption
apply simp
apply (subst headers-vConc)
  apply simp
  apply (rule conjI)
  apply simp
  apply simp
  apply simp
  apply (rule conjI)
  apply (rule sym, simp)
  apply simp
apply (subst headers-vConc)
  apply simp
  apply (rule conjI)
  apply simp
  apply simp
  apply (simp del: const)
apply (simp del: const)
done

```

lemma reg2dim2-dConc[simp]:

```

[[  $\bigwedge t . \text{headers } (h \ t) = \text{headers } t;$ 
  regSkelOuter2 (Rep-regT2 t1) = Some (cs1, hs1);
  regSkelOuter2 (Rep-regT2 t2) = Some (cs2, hs2)
]]  $\implies$ 
reg2dim2 (dConc h t1 t2) = append hs1 hs2
apply (cut-tac h=h and t1.0=Rep-regT2 t1 and t2.0=Rep-regT2 t2 in regSkelOuter2-dConc)
  apply (simp (no-asm-simp))
  apply assumption+
apply (simp only: Rep-regT2-inverse)
apply (subst Rep-regT2-inverse [THEN sym])
apply (rule reg2dim2)
apply assumption
done

```

lemma cs-dConc[simp]:

```

[[  $\bigwedge t . \text{headers } (h \ t) = \text{headers } t;$ 
   $\bigwedge t . h \ (h \ t) = h \ t;$ 
   $\bigwedge t1 \ t2 . h \ (hConc \ t1 \ t2) = hConc \ (h \ t1) \ (h \ t2);$ 
  regSkelOuter2 (Rep-regT2 t1) = Some (cs1, hs1);
  regSkelOuter2 (Rep-regT2 t2) = Some (cs2, hs2)
]]  $\implies$ 

```

```

    h (delH1 (Rep-regT2 (dConc h t1 t2))) =
    hConc (h (delH1 (Rep-regT2 t1)))
      (h (delH1 (Rep-regT2 t2)))
apply (cut-tac h=h and t1.0=Rep-regT2 t1 and t2.0=Rep-regT2 t2 in regSkelOuter2-dConc)
  apply (simp (no-asm-simp))
  apply assumption+
apply (subst dConc1)
  apply assumption+
apply (subst Abs-regT2-inverse, rule regT2)
  apply (subst regSkelOuter2-dConcDef)
  apply simp
  apply assumption
  apply assumption
  apply (simp (no-asm-simp))
  apply (rule conjI)
  apply (simp (no-asm-simp))
  apply (simp (no-asm-simp))
apply (simp (no-asm-simp))
apply (subst delH1-vConc)
  apply (simp (no-asm-simp))
  apply (rule conjI)
  apply (simp (no-asm-simp))
  apply (simp (no-asm-simp))
  apply (simp (no-asm-simp))
  apply (rule conjI)
  apply (rule sym, simp (no-asm-simp))
  apply (simp (no-asm-simp))
apply (simp del: const)
done

```

lemma *dConc-assoc*[simp]:

```

  [ [  $\bigwedge t . \text{headers } (h t) = \text{headers } t;$ 
    [  $\bigwedge t . h (h t) = h t;$ 
      [  $\bigwedge t1 t2 . h (hConc t1 t2) = hConc (h t1) (h t2)$ 
        ]  $\implies$  assoc (dConc h)
      ]
    ]
apply (rule assoc-intro)
apply (cut-tac x=x in Rep-regT2)
apply (cut-tac x=y in Rep-regT2)
apply (cut-tac x=z in Rep-regT2)
apply (simp add: regT2-def)
apply (drule regularOuter2, erule exE, erule exE)
apply (drule regularOuter2, erule exE, erule exE)
apply (drule regularOuter2, erule exE, erule exE)
apply (cut-tac h=h and t1.0=Rep-regT2 x and t2.0=Rep-regT2 y in regSkelOuter2-dConc)
  apply (simp (no-asm-simp))
  apply assumption+
apply (cut-tac h=h and t1.0=Rep-regT2 y and t2.0=Rep-regT2 z in regSkelOuter2-dConc)
  apply (simp (no-asm-simp))
  apply assumption+

```

```

apply (simp only: Rep-regT2-inverse)
apply (cut-tac h=h and t1.0=Rep-regT2 (dConc h x y) and t2.0=Rep-regT2 z
in regSkelOuter2-dConc)
  apply (simp (no-asm-simp))
  apply assumption+
apply (cut-tac h=h and t1.0=Rep-regT2 x and t2.0=Rep-regT2 (dConc h y z)
in regSkelOuter2-dConc)
  apply (simp (no-asm-simp))
  apply assumption+
apply (simp only: Rep-regT2-inverse)
apply (rotate-tac -1, frule regSkelOuter2-eq-Some)
apply (rotate-tac -1, frule regSkelOuter2-eq-Some)
apply (rotate-tac -1, frule regSkelOuter2-eq-Some)
apply (rotate-tac -1, frule regSkelOuter2-eq-Some)
apply (rotate-tac -1, frule regSkelOuter2-eq-Some)
apply (rotate-tac -1, frule regSkelOuter2-eq-Some)
apply (rotate-tac -1, frule regSkelOuter2-eq-Some)
apply (rotate-tac -1, frule regSkelOuter2-eq-Some)
apply (cut-tac h=h and t1.0=dConc h x y and t2.0=z in dConc1)
  apply assumption
  apply assumption
apply (cut-tac h=h and t1.0=x and t2.0=dConc h y z in dConc1)
  apply assumption
  apply assumption
apply (simp (no-asm-simp))
apply (subst cs-dConc [of h])
  apply (simp (no-asm-simp))
  apply (simp (no-asm-simp))
  apply (simp (no-asm-simp))
  apply assumption
  apply assumption
apply (subst cs-dConc [of h])
  apply (simp (no-asm-simp))
  apply (simp (no-asm-simp))
  apply (simp (no-asm-simp))
  apply assumption
  apply assumption
apply (subst dConc2 [of h])
  apply (simp (no-asm-simp))
  apply assumption
  apply assumption
apply (subst dConc2 [of h])
  apply (simp (no-asm-simp))
  apply assumption
  apply assumption

```

Lots of preparations

```

apply (frule-tac t=Rep-regT2 x in reg2dim2)
apply (frule-tac t=Rep-regT2 y in reg2dim2)
apply (frule-tac t=Rep-regT2 z in reg2dim2)

```

```

apply (simp only: Rep-regT2-inverse)
apply (cut-tac t1.0=Rep-regT2 y and t2.0=(tMap (const (h (delH1 (Rep-regT2 z)))) (Rep-regT2 y)) in regSkelOuter2-vConc)
  apply assumption
  apply (simp (no-asm-simp))
  apply (rule conjI)
  apply (simp (no-asm-simp))
  apply (simp (no-asm-simp))
  apply (rule sym, simp (no-asm-simp))
apply (rotate-tac -1, frule regSkelOuter2-eq-Some)
apply (cut-tac t1.0=(tMap (const (h (delH1 (Rep-regT2 y)))) (Rep-regT2 z)) and t2.0=Rep-regT2 z in regSkelOuter2-vConc)
  apply (simp (no-asm-simp))
  apply (rule conjI)
  apply (simp (no-asm-simp))
  apply (simp (no-asm-simp))
  apply assumption
  apply (simp (no-asm-simp))
apply (rotate-tac -1, frule regSkelOuter2-eq-Some)
apply (cut-tac t1.0=Rep-regT2 x and t2.0=(tMap (const (h (delH1 (Rep-regT2 y)))) (Rep-regT2 x)) in regSkelOuter2-vConc)
  apply assumption
  apply (simp (no-asm-simp))
  apply (rule conjI)
  apply (simp (no-asm-simp))
  apply (simp (no-asm-simp))
  apply (rule sym, simp (no-asm-simp))
apply (rotate-tac -1, frule regSkelOuter2-eq-Some)
apply (cut-tac t1.0=(tMap (const (h (delH1 (Rep-regT2 x)))) (Rep-regT2 y)) and t2.0=Rep-regT2 y in regSkelOuter2-vConc)
  apply (simp (no-asm-simp))
  apply (rule conjI)
  apply (simp (no-asm-simp))
  apply (simp (no-asm-simp))
  apply assumption
  apply (simp (no-asm-simp))
apply (rotate-tac -1, frule regSkelOuter2-eq-Some)

```

end of preparations

The remainder is rewriting to normal form with *vConc-hConc*.

```

apply (simp (no-asm-simp) del: const)
apply (subst vConc-hConc)
  apply (simp (no-asm-simp))
  apply (rule conjI)
  apply (simp (no-asm-simp))
  apply (simp (no-asm-simp))
  apply (simp (no-asm-simp))
  apply assumption
  apply (simp (no-asm-simp))

```

```

    apply (subst headers-delH1--reg2, assumption)
    apply (subst headers-delH1--reg2, assumption)
  apply (simp (no-asm-simp))
apply (subst vConc-hConc)
  apply assumption
  apply (simp (no-asm-simp) del: const)
  apply (rule conjI)
  apply (simp (no-asm-simp))
    apply (subst headers-delH1--reg2, assumption)
    apply (subst headers-delH1--reg2, assumption)
  apply (simp (no-asm-simp))
  apply (simp (no-asm-simp) del: const)
  apply (simp (no-asm-simp) del: const)
apply (simp (no-asm-simp) del: const)
apply (subst tMap-const-vConc)
  apply (simp (no-asm-simp) del: const)
  apply (rule conjI)
  apply (simp (no-asm-simp))
    apply (subst headers-delH1--reg2, assumption)
  apply (simp (no-asm-simp))
  apply assumption
  apply (simp (no-asm-simp))
apply (subst tMap-const-vConc)
  apply (simp (no-asm-simp) del: const)
  apply (rule conjI)
  apply (simp (no-asm-simp))
  apply (simp (no-asm-simp))
  apply (simp (no-asm-simp) del: const)
  apply (rule conjI)
  apply (simp (no-asm-simp))
    apply (subst headers-delH1--reg2, assumption)
  apply (simp (no-asm-simp))
  apply (rule sym, simp (no-asm-simp))
apply (subst tMap-const-vConc)
  apply (simp (no-asm-simp) del: const)
  apply (rule conjI)
  apply (simp (no-asm-simp))
    apply (subst headers-delH1--reg2, assumption)
  apply (simp (no-asm-simp))
  apply assumption
  apply (simp (no-asm-simp))
apply (subst tMap-const-vConc)
  apply (simp (no-asm-simp) del: const)
  apply (rule conjI)
  apply (simp (no-asm-simp))
  apply (simp (no-asm-simp))
  apply (simp (no-asm-simp) del: const)
  apply (rule conjI)
  apply (simp (no-asm-simp))

```

```

  apply (subst headers-delH1--reg2, assumption)
  apply (simp (no-asm-simp))
  apply (rule sym, simp (no-asm-simp))
  apply (subst tMap-const-hConc)+
  apply (simp (no-asm-simp) del: const)
done

```

6.10 Lifting of Inversion Combinators to the Next Dimension

The initial inversion operator is *addH2*, as used in the definition of *inverse1*. The first argument *u* is the contents of “empty” cells created by inversion. The second argument *h* converts *u* into an update function for *dConc*, and the third argument *g* is the inversion operator of the next-lower dimension.

consts *invLift0* ::

```

'a ⇒
('a ⇒ ('e,'d) T ⇒ ('e,'d) T) ⇒
('a ⇒ 'b ⇒ ('c,'d) T) ⇒
('a ⇒ ('e,'b) T ⇒ ('c,'e,'d) regT2)

```

defs *invLift0-def*:

```

invLift0 u h g h1 ==
  tFold (% h2 t2 . Abs-regT2 (addH2 h2 (g h1 t2))) (dConc (h u))

```

lemma *invLift0-addH[simp]*:

```

invLift0 u h g h1 (addH h2 t2) = Abs-regT2 (addH2 h2 (g h1 t2))

```

by (*unfold invLift0-def, simp*)

lemma *invLift0-addH1[simp]*:

```

Rep-regT2 (invLift0 u h g h1 (addH h2 t2)) = addH2 h2 (g h1 t2)

```

apply *simp*

apply (*subst Abs-regT2-inverse, rule regT2, simp*)

apply (*rule conjI, simp-all*)

done

lemma *reg2dim2-invLift0-addH[simp]*:

```

reg2dim2 (invLift0 u h g h1 (addH h2 t2)) = singleton h2

```

by *simp*

lemma *cs-invLift0-addH[simp]*:

```

[[ ∧ u v t . h u (g v t) = g u t ]] ==>

```

```

  h u (delH2 (Rep-regT2 (invLift0 u h g h1 (addH h2 t2)))) = g u t2

```

apply *simp*

apply (*subst Abs-regT2-inverse, rule regT2, simp*)

apply (*rule conjI, simp-all*)

done

lemma *invLift0-hConc[simp]*:

$\llbracket \bigwedge t . \text{headers } (h \ u \ t) = \text{headers } t;$
 $\bigwedge t . h \ u \ (h \ u \ t) = h \ u \ t;$
 $\bigwedge t1 \ t2 . h \ u \ (h\text{Conc } t1 \ t2) = h\text{Conc } (h \ u \ t1) \ (h \ u \ t2)$
 $\rrbracket \implies \text{invLift0 } u \ h \ g \ h1 \ (h\text{Conc } t1 \ t2) =$
 $d\text{Conc } (h \ u) \ (\text{invLift0 } u \ h \ g \ h1 \ t1) \ (\text{invLift0 } u \ h \ g \ h1 \ t2)$
by (*unfold invLift0-def, simp*)

lemma *reg2dim2-invLift0*:

$\llbracket \bigwedge t . \text{headers } (h \ u \ t) = \text{headers } t;$
 $\bigwedge t . h \ u \ (h \ u \ t) = h \ u \ t;$
 $\bigwedge t1 \ t2 . h \ u \ (h\text{Conc } t1 \ t2) = h\text{Conc } (h \ u \ t1) \ (h \ u \ t2);$
 $\bigwedge x \ t . \text{headers } (g \ x \ t) = \text{headers } t \rrbracket \implies$
 $\text{reg2dim2 } (\text{invLift0 } u \ h \ g \ h1 \ t) = \text{headers } t$
apply (*unfold invLift0-def, induct-tac t rule: T-induct, simp-all*)
apply (*fold invLift0-def*)
apply (*cut-tac x=invLift0 u h g h1 t1 in Rep-regT2*)
apply (*cut-tac x=invLift0 u h g h1 t2 in Rep-regT2*)
apply (*subst reg2dim2-dConc*)
apply *simp*
apply (*simp add: regT2-def*)
apply (*drule regularOuter2, erule exE, erule exE, simp*)
apply (*rule conjI*)
apply (*rule sym, erule regSkelOuter2-eq-Some*)
apply (*drule reg2dim2*)
apply (*rule sym, assumption*)
apply (*simp add: regT2-def*)
apply (*rotate-tac -1, drule regularOuter2, erule exE, erule exE, simp*)
apply (*rule conjI*)
apply (*rule sym, erule regSkelOuter2-eq-Some*)
apply (*drule reg2dim2*)
apply (*rule sym, assumption*)
apply (*simp add: Rep-regT2-inverse*)
done

lemma *regSkelOuter2-invLift0-EX*:

$\llbracket \bigwedge t . \text{headers } (h \ u \ t) = \text{headers } t;$
 $\bigwedge t . h \ u \ (h \ u \ t) = h \ u \ t;$
 $\bigwedge t1 \ t2 . h \ u \ (h\text{Conc } t1 \ t2) = h\text{Conc } (h \ u \ t1) \ (h \ u \ t2);$
 $\bigwedge x \ y \ t . \text{headers } (g \ x \ t) = \text{headers } (g \ y \ t) \rrbracket \implies$
 $EX \ hs1 .$
 $\text{regSkelOuter2 } (\text{Rep-regT2 } (\text{invLift0 } u \ h \ g \ h1 \ t)) = \text{Some } (hs1, \text{headers } t)$
apply (*unfold invLift0-def, induct-tac t rule: T-induct, simp-all*)
apply (*subst Abs-regT2-inverse, rule regT2, simp*)
apply (*rule conjI*)
apply *simp*
apply *simp*
apply *simp*
apply (*fold invLift0-def*)
apply (*erule exE, erule exE*)

```

apply (subst regSkelOuter2-dConc1)
  apply simp
  apply assumption
  apply assumption
apply (rule-tac x=append hs1 hs1a in exI, simp)
done

```

```

lemma cs-invLift0-hConc[simp]:
   $\llbracket \bigwedge t . \text{headers } (h \ u \ t) = \text{headers } t;$ 
   $\bigwedge t . h \ u \ (h \ u \ t) = h \ u \ t;$ 
   $\bigwedge t1 \ t2 . h \ u \ (hConc \ t1 \ t2) = hConc \ (h \ u \ t1) \ (h \ u \ t2)$ 
   $\rrbracket \implies$ 
   $h \ u \ (\text{delH1 } (\text{Rep-regT2 } (\text{invLift0 } u \ h \ g \ h1 \ (hConc \ t1 \ t2)))) =$ 
   $hConc \ (h \ u \ (\text{delH1 } (\text{Rep-regT2 } (\text{invLift0 } u \ h \ g \ h1 \ t1))))$ 
   $(h \ u \ (\text{delH1 } (\text{Rep-regT2 } (\text{invLift0 } u \ h \ g \ h1 \ t2))))$ 
apply (cut-tac x=invLift0 u h g h1 t1 in Rep-regT2)
apply (cut-tac x=invLift0 u h g h1 t2 in Rep-regT2)
apply (simp add: regT2-def)
apply (drule regularOuter2, erule exE, erule exE)
apply (drule regularOuter2, erule exE, erule exE)
apply (subst cs-dConc [of h u])
  apply simp
  apply simp
  apply simp
  apply simp
  apply (rule conjI)
  apply simp
  apply simp
  apply simp
  apply (rule conjI)
  apply simp
  apply simp
apply simp
done

```

```

consts invLift ::
  'a  $\Rightarrow$ 
  ('a  $\Rightarrow$  ('e,'d) T  $\Rightarrow$  ('e,'d) T)  $\Rightarrow$ 
  ('a  $\Rightarrow$  'b  $\Rightarrow$  ('c,'d) T)  $\Rightarrow$ 
  ('a  $\Rightarrow$  ('e,'b) T  $\Rightarrow$  ('c,('e,'d) T) T)

```

```

defs invLift-def: invLift u h g h1 t == Rep-regT2 (invLift0 u h g h1 t)

```

```

lemma invLift-addH[simp]:
   $\llbracket \bigwedge x \ y \ t . \text{headers } (g \ x \ t) = \text{headers } (g \ y \ t) \rrbracket \implies$ 
   $\text{invLift } u \ h \ g \ h1 \ (\text{addH } h2 \ t2) = \text{addH2 } h2 \ (g \ h1 \ t2)$ 
apply (unfold invLift-def, simp)
apply (subst Abs-regT2-inverse, simp-all add: regT2-def)
apply (rule regularOuter2I, auto)

```

done

lemma *headers-invLift0-invariant*:

$\llbracket \bigwedge t . \text{headers } (h \ u \ t) = \text{headers } t;$
 $\bigwedge t . h \ u \ (h \ u \ t) = h \ u \ t;$
 $\bigwedge t1 \ t2 . h \ u \ (h\text{Conc } t1 \ t2) = h\text{Conc } (h \ u \ t1) \ (h \ u \ t2);$
 $\bigwedge x \ y \ t . \text{headers } (g \ x \ t) = \text{headers } (g \ y \ t) \rrbracket \implies$
 $\text{headers } (\text{Rep-regT2 } (\text{invLift0 } u \ h \ g \ x \ t)) =$
 $\text{headers } (\text{Rep-regT2 } (\text{invLift0 } u \ h \ g \ y \ t))$

apply (*induct-tac t rule: T-induct*)

apply (*subst invLift0-addH1*)

apply (*subst invLift0-addH1*)

apply *simp*

apply *simp*

done

The following lemma justifies the build-up for higher-dimensional inversion functions.

lemma *headers-invLift-invariant*:

$\llbracket \bigwedge t . \text{headers } (h \ u \ t) = \text{headers } t;$
 $\bigwedge t . h \ u \ (h \ u \ t) = h \ u \ t;$
 $\bigwedge t1 \ t2 . h \ u \ (h\text{Conc } t1 \ t2) = h\text{Conc } (h \ u \ t1) \ (h \ u \ t2);$
 $\bigwedge x \ y \ t . \text{headers } (g \ x \ t) = \text{headers } (g \ y \ t) \rrbracket \implies$
 $\text{headers } (\text{invLift } u \ h \ g \ x \ t) = \text{headers } (\text{invLift } u \ h \ g \ y \ t)$

apply (*unfold invLift-def*)

apply (*cut-tac x=invLift0 u h g x t in Rep-regT2*)

apply (*cut-tac x=invLift0 u h g y t in Rep-regT2*)

apply (*unfold regT2-def, simp*)

apply (*drule regularOuter2, erule exE, erule exE*)

apply (*drule regularOuter2, erule exE, erule exE*)

apply (*drule regSkelOuter2-eq-Some*)

apply (*drule regSkelOuter2-eq-Some*)

apply (*rule headers-invLift0-invariant*)

apply *simp-all*

done

lemma *invLift-hConc[simp]*:

$\llbracket \bigwedge t . \text{headers } (h \ u \ t) = \text{headers } t;$
 $\bigwedge t . h \ u \ (h \ u \ t) = h \ u \ t;$
 $\bigwedge t1 \ t2 . h \ u \ (h\text{Conc } t1 \ t2) = h\text{Conc } (h \ u \ t1) \ (h \ u \ t2);$
 $\bigwedge x \ y \ t . \text{headers } (g \ x \ t) = \text{headers } (g \ y \ t);$
 $\text{regSkelOuter2 } (\text{invLift } u \ h \ g \ h1 \ t1) = \text{Some } (hs1a, hs2a);$
 $\text{regSkelOuter2 } (\text{invLift } u \ h \ g \ h1 \ t2) = \text{Some } (hs1b, hs2b)$
 $\rrbracket \implies$
 $\text{invLift } u \ h \ g \ h1 \ (h\text{Conc } t1 \ t2) =$
 $h\text{Conc } (v\text{Conc } (\text{invLift } u \ h \ g \ h1 \ t1)$
 $\quad (\text{tMap } (\text{const } (h \ u \ (\text{delH1 } (\text{invLift } u \ h \ g \ h1 \ t2)))) (\text{invLift } u \ h \ g \ h1$
 $t1)))$
 $\quad (v\text{Conc } (\text{tMap } (\text{const } (h \ u \ (\text{delH1 } (\text{invLift } u \ h \ g \ h1 \ t1)))) (\text{invLift } u \ h \ g \ h1$
 $t1)))$

$t2))$
 $(invLift\ u\ h\ g\ h1\ t2))$
apply $(unfold\ invLift-def,\ simp)$
apply $(subst\ dConc2\ [THEN\ sym],\ auto)$
done

lemma *regSkelOuter2-invLift-EX*:
 $\llbracket \bigwedge t . headers\ (h\ u\ t) = headers\ t;$
 $\bigwedge t . h\ u\ (h\ u\ t) = h\ u\ t;$
 $\bigwedge t1\ t2 . h\ u\ (hConc\ t1\ t2) = hConc\ (h\ u\ t1)\ (h\ u\ t2);$
 $\bigwedge x\ y\ t . headers\ (g\ x\ t) = headers\ (g\ y\ t) \rrbracket \implies$
 $EX\ hs1 . regSkelOuter2\ (invLift\ u\ h\ g\ h1\ t) = Some\ (hs1,\ headers\ t)$
apply $(unfold\ invLift-def)$
apply $(rule\ regSkelOuter2-invLift0-EX)$
apply *simp-all*
done

lemma *delH1-invLift*:
 $\llbracket \bigwedge t . headers\ (h\ u\ t) = headers\ t;$
 $\bigwedge t . h\ u\ (h\ u\ t) = h\ u\ t;$
 $\bigwedge t1\ t2 . h\ u\ (hConc\ t1\ t2) = hConc\ (h\ u\ t1)\ (h\ u\ t2);$
 $\bigwedge x\ y\ t . headers\ (g\ x\ t) = headers\ (g\ y\ t) \rrbracket \implies$
 $delH1\ (invLift\ u\ h\ g\ v\ t) = tFold\ (\% h0\ t0 . addH\ h0\ (delH1\ (g\ v\ t0)))\ (hConcU$
 $(h\ u))\ t$
apply $(induct-tac\ t\ rule:\ T-induct,\ simp-all\ del:\ const)$
apply $(cut-tac\ u=u\ and\ h=h\ and\ g=g\ and\ h1.0=v\ and\ t=t1\ in\ regSkelOuter2-invLift-EX)$
apply $(simp\ (no-asm-simp))$
apply $(simp\ (no-asm-simp))$
apply $(simp\ (no-asm-simp))$
apply $(simp\ (no-asm-simp))$
apply $(erule\ exE)$
apply $(cut-tac\ u=u\ and\ h=h\ and\ g=g\ and\ h1.0=v\ and\ t=t2\ in\ regSkelOuter2-invLift-EX)$
apply $(simp\ (no-asm-simp))$
apply $(simp\ (no-asm-simp))$
apply $(simp\ (no-asm-simp))$
apply $(simp\ (no-asm-simp))$
apply $(erule\ exE)$
apply $(subst\ invLift-hConc)$
apply $(simp\ (no-asm-simp))$
apply $(simp\ (no-asm-simp))$
apply $(simp\ (no-asm-simp))$
apply $(simp\ (no-asm-simp))$
apply *assumption*
apply *assumption*
apply $(simp\ (no-asm-simp))$
apply $(subst\ delH1-vConc)$
apply *assumption*
apply $(simp\ (no-asm-simp))$

```

  apply (rule conjI, simp (no-asm-simp), simp (no-asm-simp))
apply (subst delH1-tMap-const)
apply (simp (no-asm-simp))
apply (fold hConcU-def)
apply (cut-tac u=h u in hConcU-assoc)
  apply (simp (no-asm-simp))
  apply (simp (no-asm-simp))
apply (simp del: hConcU-def)
done

```

lemma *h-u-delH1-invLift*:

```

[[  $\bigwedge t . \text{headers } (h \ u \ t) = \text{headers } t;$ 
   $\bigwedge t \ x . h \ u \ (h \ x \ t) = h \ u \ t;$ 
   $\bigwedge t1 \ t2 . h \ u \ (hConc \ t1 \ t2) = hConc \ (h \ u \ t1) \ (h \ u \ t2);$ 
   $\bigwedge x \ y \ t . \text{headers } (g \ x \ t) = \text{headers } (g \ y \ t);$ 
   $\bigwedge h0 \ t0 . h \ u \ (\text{addH } h0 \ (\text{delH1 } (g \ v \ t0))) = \text{addH } h0 \ (\text{delH1 } (g \ u \ t0))$ 
]]  $\implies$ 
 $h \ u \ (\text{delH1 } (\text{invLift } u \ h \ g \ v \ t)) = tFold \ (\% \ h0 \ t0 . \text{addH } h0 \ (\text{delH1 } (g \ u \ t0)))$ 
 $hConc \ t$ 
apply (subst delH1-invLift)
  apply (simp (no-asm-simp))
  apply (simp (no-asm-simp))
  apply (simp (no-asm-simp))
  apply (simp (no-asm-simp))
apply (simp (no-asm-simp))
done

```

lemma *headers-h-u-delH1-invLift[simp]*:

```

[[  $\bigwedge t . \text{headers } (h \ u \ t) = \text{headers } t;$ 
   $\bigwedge t \ x . h \ u \ (h \ x \ t) = h \ u \ t;$ 
   $\bigwedge t1 \ t2 . h \ u \ (hConc \ t1 \ t2) = hConc \ (h \ u \ t1) \ (h \ u \ t2);$ 
   $\bigwedge x \ y \ t . \text{headers } (g \ x \ t) = \text{headers } (g \ y \ t);$ 
   $\bigwedge h0 \ t0 . h \ u \ (\text{addH } h0 \ (\text{delH1 } (g \ v \ t0))) = \text{addH } h0 \ (\text{delH1 } (g \ u \ t0))$ 
]]  $\implies$ 
 $\text{headers } (h \ u \ (\text{delH1 } (\text{invLift } u \ h \ g \ v \ t))) = \text{headers } t$ 
apply (subst h-u-delH1-invLift [of h u g v t])
  apply (simp (no-asm-simp))
  apply (simp (no-asm-simp))
  apply (simp (no-asm-simp))
  apply (simp (no-asm-simp))
  apply (simp (no-asm-simp))
apply (simp (no-asm-simp))
apply (subst headers-def)
apply (simp (no-asm-simp))
done

```

lemma *headers-delH1-invLift[simp]*:

```

[[  $\bigwedge t . \text{headers } (h \ u \ t) = \text{headers } t;$ 
   $\bigwedge t \ x . h \ u \ (h \ x \ t) = h \ u \ t;$ 

```

$\wedge t1\ t2 . h\ u\ (hConc\ t1\ t2) = hConc\ (h\ u\ t1)\ (h\ u\ t2);$
 $\wedge x\ y\ t . headers\ (g\ x\ t) = headers\ (g\ y\ t);$
 $\wedge h0\ t0 . h\ u\ (addH\ h0\ (delH1\ (g\ v\ t0))) = addH\ h0\ (delH1\ (g\ u\ t0))$
 $\] \Rightarrow$
 $headers\ (delH1\ (invLift\ u\ h\ g\ v\ t)) = headers\ t$
apply (*cut-tac headers-h-u-delH1-invLift [of h u g v t]*)
apply *simp*
apply (*simp (no-asm-simp)*)
apply (*simp (no-asm-simp)*)
apply (*simp (no-asm-simp)*)
apply (*simp (no-asm-simp)*)
apply (*simp (no-asm-simp)*)
done

lemma *regSkelOuter2-dConcFill*:

$\[\wedge t . headers\ (h\ u\ t) = headers\ t;$
 $\wedge t\ x . h\ u\ (h\ x\ t) = h\ u\ t;$
 $\wedge t1\ t2 . h\ u\ (hConc\ t1\ t2) = hConc\ (h\ u\ t1)\ (h\ u\ t2);$
 $\wedge x\ y\ t . headers\ (g\ x\ t) = headers\ (g\ y\ t);$
 $regSkelOuter2\ (invLift\ u\ h\ g\ h1\ t1) = Some\ (hs1a,\ hs2a);$
 $regSkelOuter2\ (invLift\ u\ h\ g\ h1\ t2) = Some\ (hs1b,\ hs2b)$
 $\] \Rightarrow$
 $regSkelOuter2\ (tMap\ (const\ (h\ u\ (delH1\ (invLift\ u\ h\ g\ h1\ t2))))$
 $\quad (invLift\ u\ h\ g\ h1\ t1)) = Some\ (hs1a,\ hs2b)$
apply (*cut-tac t2.0=h u (delH1 (invLift u h g h1 t2)) and t1.0=invLift u h g h1*
t1 in regSkelOuter2-tMap-const)
apply *simp*
done

lemma *headers-invLift-hConc[simp]*:

$\[\wedge t . headers\ (h\ u\ t) = headers\ t;$
 $\wedge t\ x . h\ u\ (h\ x\ t) = h\ u\ t;$
 $\wedge t1\ t2 . h\ u\ (hConc\ t1\ t2) = hConc\ (h\ u\ t1)\ (h\ u\ t2);$
 $\wedge x\ y\ t . headers\ (g\ x\ t) = headers\ (g\ y\ t);$
 $regSkelOuter2\ (invLift\ u\ h\ g\ h1\ t1) = Some\ (hs1a,\ hs2a);$
 $regSkelOuter2\ (invLift\ u\ h\ g\ h1\ t2) = Some\ (hs1b,\ hs2b)$
 $\] \Rightarrow$
 $headers\ (invLift\ u\ h\ g\ h1\ (hConc\ t1\ t2)) =$
 $append\ (headers\ (invLift\ u\ h\ g\ h1\ t1))$
 $\quad (headers\ (invLift\ u\ h\ g\ h1\ t2))$
apply (*subst invLift-hConc*)
apply (*simp (no-asm-simp)*)
apply (*simp (no-asm-simp)*)
apply (*simp (no-asm-simp)*)
apply (*simp (no-asm-simp)*)
apply *assumption*
apply *assumption*
apply (*cut-tac h=h and u=u and g=g and h1.0=h1 and*
t1.0=t1 and hs1a=hs1a and hs2a=hs2a and

```

      t2.0=t2 and hs1b=hs1b and hs2b=hs2b in regSkelOuter2-dConcFill)
  apply (simp (no-asm-simp))
  apply (simp (no-asm-simp))
  apply (simp (no-asm-simp))
  apply (simp (no-asm-simp))
  apply assumption
  apply assumption
  apply (cut-tac h=h and u=u and g=g and h1.0=h1 and
    t1.0=t2 and hs1a=hs1b and hs2a=hs2b and
    t2.0=t1 and hs1b=hs1a and hs2b=hs2a in regSkelOuter2-dConcFill)
  apply (simp (no-asm-simp))
  apply (simp (no-asm-simp))
  apply (simp (no-asm-simp))
  apply (simp (no-asm-simp))
  apply assumption
  apply assumption
  apply (simp (no-asm-simp))
  apply (subst headers-vConc)
  apply assumption
  apply assumption
  apply (subst headers-vConc)
  apply assumption
  apply assumption
  apply (simp (no-asm-simp))
done

```

lemma *headers-invLift*[*simp*]:

```

[[  $\bigwedge t . \text{headers } (h \ u \ t) = \text{headers } t;$ 
   $\bigwedge t \ x . h \ u \ (h \ x \ t) = h \ u \ t;$ 
   $\bigwedge t1 \ t2 . h \ u \ (hConc \ t1 \ t2) = hConc \ (h \ u \ t1) \ (h \ u \ t2);$ 
   $\bigwedge x \ y \ t . \text{headers } (g \ x \ t) = \text{headers } (g \ y \ t)$ 
]]  $\implies$ 
  headers (invLift u h g h1 t) = tFold (% h2 t2 . headers (g h1 t2)) append t
  apply (induct-tac t rule: T-induct, simp)
  apply (cut-tac u=u and h=h and g=g and h1.0=h1 and t=t1 in regSkelOuter2-invLift-EX)
    apply (simp (no-asm-simp))
    apply (simp (no-asm-simp))
    apply (simp (no-asm-simp))
    apply (simp (no-asm-simp))
    apply (erule exE)
  apply (cut-tac u=u and h=h and g=g and h1.0=h1 and t=t2 in regSkelOuter2-invLift-EX)
    apply (simp (no-asm-simp))
    apply (simp (no-asm-simp))
    apply (simp (no-asm-simp))
    apply (simp (no-asm-simp))
    apply (erule exE)
  apply (subst headers-invLift-hConc)
    apply (simp (no-asm-simp))
    apply (simp (no-asm-simp))

```

```

apply (simp (no-asm-simp))
apply (simp (no-asm-simp))
apply assumption
apply assumption
apply (simp (no-asm-simp))
done

```

lemma *regSkelOuter2-invLift[*simp*]*:

```

[[  $\bigwedge t . \text{headers } (h \ u \ t) = \text{headers } t;$ 
 $\bigwedge t \ x . h \ u \ (h \ x \ t) = h \ u \ t;$ 
 $\bigwedge t1 \ t2 . h \ u \ (h\text{Conc } t1 \ t2) = h\text{Conc } (h \ u \ t1) \ (h \ u \ t2);$ 
 $\bigwedge x \ y \ t . \text{headers } (g \ x \ t) = \text{headers } (g \ y \ t);$ 
 $\bigwedge h0 \ t0 . h \ u \ (\text{addH } h0 \ (\text{delH1 } (g \ h1 \ t0))) = \text{addH } h0 \ (\text{delH1 } (g \ u \ t0))$ 
]]  $\implies$ 

```

```

regSkelOuter2 (invLift u h g h1 t) =

```

```

Some (tFold (% h2 t2 . headers (g h1 t2)) append t, headers t)

```

```

apply (induct-tac t rule: T-induct, simp)

```

```

apply (cut-tac u=u and h=h and g=g and h1.0=h1 and t=t1 in regSkelOuter2-invLift-EX)

```

```

apply (simp (no-asm-simp))

```

```

apply (simp (no-asm-simp))

```

```

apply (simp (no-asm-simp))

```

```

apply (simp (no-asm-simp))

```

```

apply (erule exE)

```

```

apply (cut-tac u=u and h=h and g=g and h1.0=h1 and t=t2 in regSkelOuter2-invLift-EX)

```

```

apply (simp (no-asm-simp))

```

```

apply (simp (no-asm-simp))

```

```

apply (simp (no-asm-simp))

```

```

apply (simp (no-asm-simp))

```

```

apply (erule exE)

```

```

apply (subst invLift-hConc)

```

```

apply (simp (no-asm-simp))

```

```

apply (simp (no-asm-simp))

```

```

apply (simp (no-asm-simp))

```

```

apply (simp (no-asm-simp))

```

```

apply assumption

```

```

apply assumption

```

```

apply (subst regSkelOuter2-hConc)

```

```

apply (cut-tac h=h and u=u and g=g and h1.0=h1 and

```

```

 $t1.0=t1$  and  $hs1a=hs1$  and  $hs2a=headers \ t1$  and

```

```

 $t2.0=t2$  and  $hs1b=hs1a$  and  $hs2b=headers \ t2$  in regSkelOuter2-dConcFill)

```

```

apply (simp (no-asm-simp))

```

```

apply (simp (no-asm-simp))

```

```

apply (simp (no-asm-simp))

```

```

apply (simp (no-asm-simp))

```

```

apply assumption

```

```

apply assumption

```

```

apply (cut-tac h=h and u=u and g=g and h1.0=h1 and

```

```

 $t1.0=t2$  and  $hs1a=hs1a$  and  $hs2a=headers \ t2$  and

```

```

 $t2.0=t1$  and  $hs1b=hs1$  and  $hs2b=headers \ t1$  in regSkelOuter2-dConcFill)

```



```

apply (simp (no-asm-simp))
apply (simp (no-asm-simp))
apply (simp (no-asm-simp))
apply (simp (no-asm-simp))
apply assumption
apply assumption
apply (subst regSkelOuter2-vConc)
apply (simp (no-asm-simp))
apply (rule conjI)
apply (simp (no-asm-simp))
apply (simp (no-asm-simp))
apply assumption
apply (simp (no-asm-simp))
apply (subst regSkelOuter2-vConc)
apply assumption
apply (simp (no-asm-simp))
apply (rule conjI)
apply (simp (no-asm-simp))
apply (simp (no-asm-simp))
apply (simp (no-asm-simp))
apply (simp (no-asm-simp))
done

```

6.11 The Inversion Operator for Two Dimensions

```

consts invOp2 :: 'a ⇒ 'a ⇒ ('b,('c,'d) T) T ⇒ ('c,('b,('a,'d) T) T) T

```

```

defs invOp2-def: invOp2 u == invLift u slimH2 addH2

```

```

lemma delH1-invOp2:

```

```

  delH1 (invOp2 u v t) = tFold (% h t . addH h (slimH1 v t)) (hConcSH2 u) t
apply (unfold invOp2-def)
apply (subst delH1-invLift)
apply (simp (no-asm-simp))
apply (simp (no-asm-simp))
apply (simp (no-asm-simp))
apply (simp (no-asm-simp))
apply (simp del: tFold-hConc hConcU-def slimH2-def)
apply (fold hConcSH2-def)
apply (rule refl)
done

```

```

lemma headers-invOp2[simp]:

```

```

  headers (invOp2 u v t) = tFold (% h . headers) append t
apply (unfold invOp2-def)
apply (subst headers-invLift)
apply (simp (no-asm-simp))
apply (simp (no-asm-simp))
apply (simp (no-asm-simp))

```

```

apply (simp (no-asm-simp))
apply (simp (no-asm-simp))
done

```

```

lemma cs-invOp2[simp]:
  slimH2 u (delH1 (invOp2 u v t)) = tFold (% h t . addH h (slimH1 u t)) hConc t
apply (subst delH1-invOp2)
apply (induct-tac t rule: T-induct, simp-all del: slimH2-def)
done

```

```

lemma regSkelOuter2-invOp2:
  regSkelOuter2 (invOp2 u h1 t) = Some (tFold (% h . headers) append t, headers
t)
apply (cut-tac u=u and h=slimH2 and g=addH2 and h1.0=h1 and t=t in
regSkelOuter2-invLift)
  apply (simp (no-asm-simp))
  apply (simp (no-asm-simp))
  apply (simp (no-asm-simp))
  apply (simp (no-asm-simp))
  apply (simp (no-asm-simp))
apply (fold invOp2-def)
apply (simp (no-asm-simp))
done

```

```

lemma invOp2-hConc[simp]:
  [regSkelOuter2 (invOp2 u h1 t1) = Some (hs1a, hs2a);
  regSkelOuter2 (invOp2 u h1 t2) = Some (hs1b, hs2b)
  ]  $\implies$ 
  invOp2 u h1 (hConc t1 t2) =
  hConc (vConc (invOp2 u h1 t1)
    (spread1 hs1a (slimH2 u t2)))
    (vConc (spread1 hs1b (slimH2 u t1))
      (invOp2 u h1 t2))
apply (unfold invOp2-def, subst invLift-hConc)
  apply (simp (no-asm-simp))
  apply (simp (no-asm-simp))
  apply (simp (no-asm-simp))
  apply (simp (no-asm-simp))
  apply assumption
  apply assumption
apply (fold invOp2-def)
apply (subst cs-invOp2)
apply (subst cs-invOp2)
apply (simp (no-asm-simp))
apply (subst tMap-CONST--headers)
apply (subst tMap-CONST--headers)
apply (subst regSkelOuter2-eq-Some [of invOp2 u h1 t1 hs1a hs2a])
  apply assumption
apply (subst regSkelOuter2-eq-Some [of invOp2 u h1 t2])

```

```

apply assumption
apply (fold slimH2-def)
apply (unfold spread1-def)
apply (fold slimH1-def)
apply (simp (no-asm-simp) only: slimH2-as-tFold)
done

```

6.12 Two-Dimensional Inversion

```

consts inverse2 :: 'a ⇒ ('a,('b,('c,'d) T) T) T
          ⇒ ('c,('b,('a,'d) T) T) T

```

```

defs inverse2-def: inverse2 u == tFold (invLift u slimH2 addH2) hConc

```

```

lemma inverse2-addH[simp]: inverse2 u (addH h t) = invLift u slimH2 addH2 h t
by (unfold inverse2-def, simp)

```

```

lemma inverse2-hConc[simp]:
inverse2 u (hConc t1 t2) = hConc (inverse2 u t1) (inverse2 u t2)
by (unfold inverse2-def, simp)

```

```

lemma tFold-tFold-tFold-inverse2:

```

```

[[ assoc c2; assoc c3; assoc c1; assoc c5; assoc c6;
  ∧ h ha hb t . a4 hb (a5 ha (a6 h t)) = a1 h (a2 ha (a3 hb t));
  ∧ h ha t1 t2 .
    c1 (a1 h (a2 ha t1)) (a1 h (a2 ha t2)) = a1 h (a2 ha (c3 t1 t2));
  (* The next two are for tFold-tFold-vConc: *)
  ∧ x y z . a4 x (c5 y z) = c5 (a4 x y) (a4 x z);
  ∧ x1 x2 y1 y2 . c1 (c5 x1 y1) (c5 x2 y2) = c5 (c1 x1 x2) (c1 y1 y2);
  (* units: *)
  ∧ x . LRunit c1 (a1 u x);
  ∧ x y . LRunit c5 (a5 y (a6 u x));
  ∧ x y z . LRunit c1 (a1 z (a5 y (a6 u x)));
  (* finishing 1.2: *)
  ∧ h x1 x2 . c5 (a1 h x1) (a1 h x2) = a1 h (c2 x1 x2)
]] ⇒
tFold a4 c1 (tMap (tFold a5 c5 o (tMap (tFold a6 c6))) (inverse2 u t)) =
tFold a1 c1 (tMap (tFold a2 c2 o (tMap (tFold a3 c3))) t)

```

```

apply (induct-tac t rule: T-induct)

```

```

1: addH h t0

```

```

apply (induct-tac t0 rule: T-induct)

```

```

1.1: t0 = addH ha t0a

```

```

apply (induct-tac t0a rule: T-induct, simp)

```

```

1.1.2: t0a = hConc t1 t2

```

```

apply simp

```

```

1.2: t0 = hConc t1 t2

```

```

apply (simp del: slimH2-def tFold-tMap)
apply (fold invOp2-def)
apply (cut-tac u=u and h1.0=h and t=t1 in regSkelOuter2-invOp2)
apply (cut-tac u=u and h1.0=h and t=t2 in regSkelOuter2-invOp2)
apply (subst invOp2-hConc)
  apply assumption
  apply assumption
apply (subgoal-tac ALL t . tFold a5 c5 (tMap (tFold a6 c6) t) =
  tFold (λh t0. a5 h (tFold a6 c6 t0)) c5 t)
  prefer 2
  apply (rule allI, simp (no-asm-simp))
apply (simp (no-asm-simp) del: slimH2-def tFold-tMap)
apply (subst tFold-tFold-vConc [of c5 c1 a4])
  apply assumption
  apply assumption
  apply (simp (no-asm-simp))
  apply (simp (no-asm-simp))
  apply (simp (no-asm-simp))
  apply (rule conjI)
  apply (rule refl)
  apply (rule refl)
  apply assumption
  apply (rule refl)
apply (subst tFold-tFold-vConc [of c5 c1 a4])
  apply assumption
  apply assumption
  apply (simp (no-asm-simp))
  apply (simp (no-asm-simp))
  apply assumption
  apply (simp (no-asm-simp) del: slimH2-def)
  apply (rule conjI)
  apply (rule refl)
  apply (rule refl)
  apply (rule refl)
apply (simp (no-asm-simp) only: tMap-f-spread1)
apply (simp (no-asm-simp) only: tFold-slimH2 tFold-slimH1)
apply (simp (no-asm-simp) only: f-tFold-const)
apply (simp (no-asm-simp) only: tFold-spread1)
apply (subgoal-tac (λh. a4 h (tFold (λh0. tFold (λh0a t0. a5 h0 (a6 u t0)) const)
c5 t1)) =
  (λh. a4 h (tFold (λh0. tFold (λh0a t0. a5 arbitrary (a6 u arbitrary)) const) c5
t1))))
  prefer 2
  apply (rule ext)
  apply (rule-tac f=a4 ha in arg-cong)
  apply (rule-tac x=t1 in fun-cong)
  apply (rule-tac x=c5 in fun-cong)
  apply (rule-tac f=tFold in arg-cong)
  apply (rule ext)

```

```

apply (rule fun-cong [of - - const])
apply (rule arg-cong [of - - tFold])
apply (rule ext)
apply (rule ext)
apply (cut-tac c=c5 and F=%x . a5 x (a6 u t0) and x=h0 in LRunit-const)
apply (simp (no-asm-simp))
apply (cut-tac c=c5 and F=%x . a5 arbitrary (a6 u x) and x=t0 in LRunit-const)
apply (simp (no-asm-simp))
apply (simp (no-asm-simp))
apply (subgoal-tac ( $\lambda h. a4 h (tFold (\lambda h0. tFold (\lambda h0a t0. a5 h0 (a6 u t0))) const$ 
c5 t2)) =
( $\lambda h. a4 h (tFold (\lambda h0. tFold (\lambda h0a t0. a5 arbitrary (a6 u arbitrary)) const) c5$ 
t2)))
prefer 2
apply (rule ext)
apply (rule-tac f=a4 ha in arg-cong)
apply (rule-tac x=t2 in fun-cong)
apply (rule-tac x=c5 in fun-cong)
apply (rule-tac f=tFold in arg-cong)
apply (rule ext)
apply (rule fun-cong [of - - const])
apply (rule arg-cong [of - - tFold])
apply (rule ext)
apply (rule ext)
apply (cut-tac c=c5 and F=%x . a5 x (a6 u t0) and x=h0 in LRunit-const)
apply (simp (no-asm-simp))
apply (cut-tac c=c5 and F=%x . a5 arbitrary (a6 u x) and x=t0 in LRunit-const)
apply (simp (no-asm-simp))
apply (simp (no-asm-simp))
apply (rotate-tac -1, simp (no-asm-simp) del: slimH2-def tFold-tMap)
apply (thin-tac ALL t . tFold a5 c5 (tMap (tFold a6 c6) t) =
tFold ( $\lambda h t0. a5 h (tFold a6 c6 t0)$ ) c5 t)
apply (subgoal-tac tFold ( $\lambda h t0. a5 h (tFold a6 c6 t0)$ ) c5 = (% t . tFold a5 c5
(tMap (tFold a6 c6) t)))
prefer 2
apply (rule ext)
apply (subst tFold-tMap)
apply assumption
apply (rule refl)
apply (simp (no-asm-simp) del: tFold-tMap add: o-def)
apply (thin-tac tFold ( $\lambda h t0. a5 h (tFold a6 c6 t0)$ ) c5 = (% t . tFold a5 c5 (tMap
(tFold a6 c6) t)))
apply (simp (no-asm-simp) add: o-def LRunit-left LRunit-right)

2: t = hConc t1 t2

apply (simp del: slimH2-def tFold-tMap)
done

```

lemma *tFold-tFold-tFold0-inverse2*:

```

[[ assoc c2; assoc c3; assoc c1; assoc c5; assoc c6;
  ∧ h ha hb t . a4 hb (a5 ha h) = a1 h (a2 ha hb);
  ∧ h ha t1 t2 .
    c1 (a1 h (a2 ha t1)) (a1 h (a2 ha t2)) = a1 h (a2 ha (c3 t1 t2));
  (* The next two are for tFold-tFold-vConc: *)
  ∧ x y z . a4 x (c5 y z) = c5 (a4 x y) (a4 x z);
  ∧ x1 x2 y1 y2 . c1 (c5 x1 y1) (c5 x2 y2) = c5 (c1 x1 x2) (c1 y1 y2);
  (* units: *)
  ∧ x . LRunit c1 (a1 u x);
  ∧ x y . LRunit c5 (a5 y u);
  ∧ x y z . LRunit c1 (a1 z (a5 y u));
  (* finishing 1.2: *)
  ∧ h x1 x2 . c5 (a1 h x1) (a1 h x2) = a1 h (c2 x1 x2)
]] ⇒
tFold a4 c1 (tMap (tFold a5 c5 o (tMap (tFold0 c6))) (inverse2 u t)) =
tFold a1 c1 (tMap (tFold a2 c2 o (tMap (tFold0 c3))) t)
apply (unfold tFold0-def)
apply (rule tFold-tFold-tFold-inverse2)
apply (simp-all (no-asm-simp))
done

```

Adapting for cells at lowest level:

lemma *tFold-tFold-tFoldC-inverse2*:

```

[[ assoc c1; assoc c2; assoc c3; assoc c5; assoc c6;
  ∧ h ha hb t . a4 hb (a5 ha (a6 h)) = a1 h (a2 ha (a3 hb));
  ∧ h ha t1 t2 .
    c1 (a1 h (a2 ha t1)) (a1 h (a2 ha t2)) = a1 h (a2 ha (c3 t1 t2));
  (* The next two are for tFold-tFold-vConc: *)
  ∧ x y z . a4 x (c5 y z) = c5 (a4 x y) (a4 x z);
  ∧ x1 x2 y1 y2 . c1 (c5 x1 y1) (c5 x2 y2) = c5 (c1 x1 x2) (c1 y1 y2);
  (* units: *)
  ∧ x . LRunit c1 (a1 u x);
  ∧ x y . LRunit c5 (a5 y (a6 u));
  ∧ x y z . LRunit c1 (a1 z (a5 y (a6 u)));
  (* finishing 1.2: *)
  ∧ h x1 x2 . c5 (a1 h x1) (a1 h x2) = a1 h (c2 x1 x2)
]] ⇒
tFold a4 c1 (tMap (tFold a5 c5 o (tMap (tFoldC a6 c6))) (inverse2 u t)) =
tFold a1 c1 (tMap (tFold a2 c2 o (tMap (tFoldC a3 c3))) t)
apply (unfold tFoldC-def)
apply (rule tFold-tFold-tFold-inverse2)
apply assumption+
apply (simp-all (no-asm-simp))
done

```

lemma *tFold-tFold-tFold-inverse2-wrapped*:

```

[[ assoc c1; assoc c2; assoc c3; assoc c4; assoc c5; assoc c6;
  ∧ h ha hb t . w2 (a4 hb (a5 ha (a6 h t))) = w1 (a1 h (a2 ha (a3 hb t)));

```

$\wedge x y . w2 (c4 x y) = c7 (w2 x) (w2 y);$
 $\wedge x y . c7 (w1 x) (w1 y) = w1 (c1 x y);$
 $\wedge x y . w2 (c5 x y) = c8 (w2 x) (w2 y);$
 $\wedge h x y . c7 (w1 (a1 h x)) (w1 (a1 h y)) = w1 (a1 h (c2 x y));$
 $\wedge h ha x y .$
 $c7 (w1 (a1 h (a2 ha x))) (w1 (a1 h (a2 ha y))) =$
 $w1 (a1 h (a2 ha (c3 x y)));$
(The next two are for tFold-tFold-vConc: *)*
 $\wedge x y z . a4 x (c5 y z) = c5 (a4 x y) (a4 x z);$
 $\wedge x1 x2 y1 y2 . c4 (c5 x1 y1) (c5 x2 y2) = c5 (c4 x1 x2) (c4 y1 y2);$
(units: *)*
 $\wedge x . LRunit c8 (w1 (a1 u x));$
 $\wedge x y . LRunit c5 (a5 y (a6 u x));$
 $\wedge x y z . LRunit c4 (a4 z (a5 y (a6 u x)));$
(finishing 1.2: *)*
 $\wedge h x1 x2 . c5 (a1 h x1) (a1 h x2) = a1 h (c2 x1 x2)$
 $\boxed{\implies}$
 $w2 (tFold a4 c4 (tMap (tFold a5 c5 o (tMap (tFold a6 c6))) (inverse2 u t))) =$
 $w1 (tFold a1 c1 (tMap (tFold a2 c2 o (tMap (tFold a3 c3))) t))$
apply (*induct-tac t rule: T-induct*)

1: *addH h t0*

apply (*induct-tac t0 rule: T-induct*)

1.1: *t0 = addH ha t0a*

apply (*induct-tac t0a rule: T-induct, simp (no-asm-simp)*)

1.1.2: *t0a = hConc t1 t2*

apply *simp*

1.2: *t0 = hConc t1 t2*

apply (*simp del: slimH2-def tFold-tMap*)

apply (*fold invOp2-def*)

apply (*cut-tac u=u and h1.0=h and t=t1 in regSkelOuter2-invOp2*)

apply (*cut-tac u=u and h1.0=h and t=t2 in regSkelOuter2-invOp2*)

apply (*subst invOp2-hConc*)

apply *assumption*

apply *assumption*

apply (*subgoal-tac ALL t . tFold a5 c5 (tMap (tFold a6 c6) t) =*
 $tFold (\lambda h t0 . a5 h (tFold a6 c6 t0)) c5 t$)

prefer 2

apply (*rule allI, simp (no-asm-simp)*)

apply (*simp (no-asm-simp) del: slimH2-def tFold-tMap*)

apply (*subst tFold-tFold-vConc [of c5 c4 a4]*)

apply *assumption*

apply *assumption*

apply (*simp (no-asm-simp)*)

apply (*simp (no-asm-simp)*)

apply (*simp (no-asm-simp)*)

```

    apply (rule conjI, rule refl, rule refl)
  apply assumption
  apply (rule refl)
apply (subst tFold-tFold-vConc [of c5 c4 a4])
  apply assumption
  apply assumption
  apply (simp (no-asm-simp))
  apply (simp (no-asm-simp))
  apply assumption
  apply (simp (no-asm-simp) del: slimH2-def)
    apply (rule conjI, rule refl, rule refl)
  apply (rule refl)
apply (simp (no-asm-simp) only: tMap-f-spread1)
apply (simp (no-asm-simp) only: tFold-slimH2 tFold-slimH1)
apply (simp (no-asm-simp) only: f-tFold-const)
apply (subgoal-tac (λh0 . tFold (λh0a t0 . a5 h0 (a6 u t0)) const) = (λ h0 t0 .
a5 arbitrary (a6 u arbitrary)))
  prefer 2
  apply (subgoal-tac (λh0 . tFold (λh0a t0 . a5 h0 (a6 u t0)) const) = (λh0 .
tFold (λ h0 t0 . a5 arbitrary (a6 u arbitrary)) const))
    prefer 2
    apply (rule ext)
    apply (rule-tac x=const in fun-cong)
    apply (rule-tac f=tFold in arg-cong)
    apply (rule ext)
    apply (rule ext)
    apply (cut-tac c=c5 and F=%x . a5 x (a6 u t0) and x=h0 in LRunit-const)
      apply (simp (no-asm-simp))
      apply (cut-tac c=c5 and F=%x . a5 arbitrary (a6 u x) and x=t0 in
LRunit-const)
        apply (simp (no-asm-simp))
        apply (simp (no-asm-simp))
        apply (rotate-tac -1, erule trans)
        apply (rule ext)
        apply (rule ext)
        apply (simp (no-asm-simp))
    apply (rotate-tac -1, erule ssubst)
    apply (simp (no-asm-simp) del: slimH2-def tFold-tMap)
    apply (simp (no-asm-simp) only: tFold-spread1)
    apply (subst foldr1-map1-LRunit)
      apply (simp (no-asm-simp))
      apply assumption
    apply (subst foldr1-map1-LRunit)
      apply (simp (no-asm-simp))
      apply assumption
    apply (thin-tac ALL t . tFold a5 c5 (tMap (tFold a6 c6) t) =
tFold (λh t0. a5 h (tFold a6 c6 t0)) c5 t)
    apply (subgoal-tac tFold (λh t0. a5 h (tFold a6 c6 t0)) c5 = (% t . tFold a5 c5
(tMap (tFold a6 c6) t)))

```



```

prefer 2
apply (rule ext)
apply (subst tFold-tMap)
apply assumption
apply (rule refl)
apply (simp (no-asm-simp) del: tFold-tMap add: LRunit-left LRunit-right)

```

2: $t = hConc\ t1\ t2$

```

apply (simp del: slimH2-def tFold-tMap)
done

```

lemma *tFold-tFold-tFoldC-inverse2-wrapped*:

```

[[ assoc c1; assoc c2; assoc c3; assoc c4; assoc c5; assoc c6;
  ∧ h ha hb t . w2 (a4 hb (a5 ha (a6 h))) = w1 (a1 h (a2 ha (a3 hb)));
  ∧ x y . w2 (c4 x y) = c7 (w2 x) (w2 y);
  ∧ x y . c7 (w1 x) (w1 y) = w1 (c1 x y);
  ∧ x y . w2 (c5 x y) = c8 (w2 x) (w2 y);
  ∧ h x y . c7 (w1 (a1 h x)) (w1 (a1 h y)) = w1 (a1 h (c2 x y));
  ∧ h ha x y .
    c7 (w1 (a1 h (a2 ha x))) (w1 (a1 h (a2 ha y))) =
      w1 (a1 h (a2 ha (c3 x y)));
  (* The next two are for tFold-tFold-vConc: *)
  ∧ x y z . a4 x (c5 y z) = c5 (a4 x y) (a4 x z);
  ∧ x1 x2 y1 y2 . c4 (c5 x1 y1) (c5 x2 y2) = c5 (c4 x1 x2) (c4 y1 y2);
  (* units: *)
  ∧ x . LRunit c8 (w1 (a1 u x));
  ∧ x y . LRunit c5 (a5 y (a6 u));
  ∧ x y z . LRunit c4 (a4 z (a5 y (a6 u)));
  (* finishing 1.2: *)
  ∧ h x1 x2 . c5 (a1 h x1) (a1 h x2) = a1 h (c2 x1 x2)
]] ⇒
w2 (tFold a4 c4 (tMap (tFold a5 c5 o (tMap (tFoldC a6 c6))) (inverse2 u t))) =
w1 (tFold a1 c1 (tMap (tFold a2 c2 o (tMap (tFoldC a3 c3))) t))
apply (unfold tFoldC-def)
apply (rule tFold-tFold-tFold-inverse2-wrapped)
apply assumption+
apply (simp-all (no-asm-simp))
apply simp-all

```

Strange: Why doesn't the first *simp-all*, even without *no-asm-simp*, go through?

done

lemma *wrapped-tFold*:

```

[[ assoc c4; assoc c7; ∧ x y . w2 (c4 x y) = c7 (w2 x) (w2 y) ]] ⇒
w2 (tFold a4 c4 t) = tFold (% h t . w2 (a4 h t)) c7 t
by (induct-tac t rule: T-induct, simp-all)

```

lemma *wrapped-foldr1*:

```

[[ assoc c5; assoc c8; ∧ x y . f (c5 x y) = c8 (f x) (f y) ]] ⇒

```

f (foldr1 c5 xs) = foldr1 c8 (map1 f xs)
by (induct-tac xs rule: neList-append-induct, simp-all)

lemma wrapped2-tFold:

[[assoc c5; assoc c7;
 $\bigwedge h x y . w2 (a4 h (c5 x y)) = c7 (w2 (a4 h x)) (w2 (a4 h y))$]] \implies
 $(\% x . w2 (a4 x (tFold a8 c5 t))) = (\% x . tFold (\% h t . w2 (a4 x (a8 h t)))$
 $c7 t)$
apply (rule ext)
apply (rule wrapped-tFold [of c5 c7])
apply assumption
apply assumption
apply (simp-all (no-asm-simp))
done

lemma tFold-tFold-tFold-inverse2-gen:

[[assoc c1; assoc c2; assoc c3; assoc c4; assoc c5; assoc c6; assoc c7;
 $\bigwedge h ha hb t . w2 (a4 hb (a5 ha (a6 h t))) = w1 (a1 h (a2 ha (a3 hb t)))$;
(* finishing 2, and influencing the next rule: *)
 $\bigwedge x y . c7 (w1 x) (w1 y) = w1 (c1 x y)$;
(* finishing 1.1.2 *)
 $\bigwedge h ha x y .$
 $w1 (c1 (a1 h (a2 ha x)) (a1 h (a2 ha y))) =$
 $w1 (a1 h (a2 ha (c3 x y)))$;
(* The next few are for tFold-tFold-vConc-wrapped: *)
 $\bigwedge h x y . w2 (a4 h (c5 x y)) = w2 (c4 (a4 h x) (a4 h y))$;
 $\bigwedge x y . w2 (c4 x y) = c7 (w2 x) (w2 y)$;
 $\bigwedge x1 x2 y1 y2 . c7 (w2 (c4 x1 x2)) (w2 (c4 y1 y2)) =$
 $w2 (c4 (c4 x1 y1) (c4 x2 y2))$;
(* stronger:
 $\bigwedge x1 x2 y1 y2 . c4 (c4 x1 x2) (c4 y1 y2) =$
 $c4 (c4 x1 y1) (c4 x2 y2)$;
*)
(* units: *)
*)
 $\bigwedge x y . LRunit c5 (a5 y (a6 u x))$;
 $\bigwedge h x y . LRunit c4 (a4 h (a5 y (a6 u x)))$;
*)
 $\bigwedge x . LRunit c7 (w1 (a1 u x))$;
(* finishing 1.2: *)
 $\bigwedge h x y . w1 (c1 (a1 h x) (a1 h y)) = w1 (a1 h (c2 x y))$
]] \implies
 $w2 (tFold a4 c4 (tMap (tFold a5 c5 o (tMap (tFold a6 c6))) (inverse2 u t))) =$
 $w1 (tFold a1 c1 (tMap (tFold a2 c2 o (tMap (tFold a3 c3))) t))$
apply (induct-tac t rule: T-induct)

1: addH h t0

apply (induct-tac t0 rule: T-induct)

1.1: t0 = addH ha t0a

apply (*induct-tac* *t0a* *rule*: *T-induct*, *simp* (*no-asm-simp*))

1.1.2: *t0a* = *hConc* *t1* *t2*

apply *simp*

1.2: *t0* = *hConc* *t1* *t2*

apply (*simp* *del*: *slimH2-def* *tFold-tMap*)

apply (*fold* *invOp2-def*)

apply (*cut-tac* *u=u* **and** *h1.0=h* **and** *t=t1* **in** *regSkelOuter2-invOp2*)

apply (*cut-tac* *u=u* **and** *h1.0=h* **and** *t=t2* **in** *regSkelOuter2-invOp2*)

apply (*subst* *invOp2-hConc*)

apply *assumption*

apply *assumption*

apply (*subgoal-tac* *ALL* *t* . *tFold* *a5* *c5* (*tMap* (*tFold* *a6* *c6*) *t*) =
tFold (λh *t0*. *a5* *h* (*tFold* *a6* *c6* *t0*)) *c5* *t*)

prefer 2

apply (*rule* *allI*, *simp* (*no-asm-simp*))

apply (*simp* (*no-asm-simp*) *del*: *slimH2-def* *tFold-tMap*)

apply (*subst* *tFold-tFold-vConc-wrapped* [*of* *c4* *c5* *c4* *w2* *a4* *w2* *a4* *c7*])

apply *assumption*

apply *assumption*

apply *assumption*

apply (*simp* (*no-asm-simp*))

apply (*simp* (*no-asm-simp*))

apply (*simp* (*no-asm-simp*))

apply (*simp* (*no-asm-simp*))

apply (*rule* *conjI*, *rule* *refl*, *rule* *refl*)

apply *assumption*

apply (*subst* *tFold-tFold-vConc-wrapped* [*of* *c4* *c5* *c4* *w2* *a4* *w2* *a4* *c7*])

apply *assumption*

apply *assumption*

apply *assumption*

apply (*simp* (*no-asm-simp*))

apply (*simp* (*no-asm-simp*))

apply (*simp* (*no-asm-simp*))

apply *assumption*

apply (*simp* (*no-asm-simp*) *del*: *slimH2-def*)

apply (*simp* (*no-asm-simp*) *only*: *spread1-slimH2*)

apply (*simp* (*no-asm-simp*) *only*: *tMap-f-slimH3* *tFold-slimH2* *tFold-slimH1*)

apply (*subst* *tFold-tMap* [*of* - - (*tFold* ($\lambda h0$ *t0*. *a5* *h0* (*tFold* ($\lambda h0$. *a6* *u*) *const* *t0*)) *c5*)]])

apply *assumption*

apply (*subst* *tFold-tMap* [*of* - - (*tFold* ($\lambda h0$ *t0*. *a5* *h0* (*tFold* ($\lambda h0$. *a6* *u*) *const* *t0*)) *c5*)]])

apply *assumption*

apply (*subst* *wrapped-tFold* [*of* *c4* *c7* *w2* (λh *t0*. *a4* *h* (*tFold* ($\lambda h0$ *t0*. *a5* *h0* (*tFold* ($\lambda h0$. *a6* *u*) *const* *t0*)) *c5* *t0*)]])

apply *assumption*

apply *assumption*

```

apply (simp (no-asm-simp))
apply (subst wrapped-tFold [of c4 c7 w2 ( $\lambda h\ t0. a4\ h\ (tFold\ (\lambda h0\ t0. a5\ h0\ (tFold\ (\lambda h0. a6\ u)\ const\ t0))\ c5\ t0))$ )]))
apply assumption
apply assumption
apply (simp (no-asm-simp))
apply (subgoal-tac ( $\lambda h\ t. w2\ (a4\ h\ (tFold\ (\lambda h0\ t0. a5\ h0\ (tFold\ (\lambda h0. a6\ u)\ const\ t0))\ c5\ t)) = (\lambda\ h\ t. w1\ (a1\ u\ arbitrary))$ ))
prefer 2
apply (rule ext)
apply (rule ext)
apply (cut-tac c4.0=c5 and c7.0=c7 and w2.0=%z . w2 (a4 ha z) in wrapped-tFold)
apply assumption
apply assumption
apply (simp (no-asm-simp))
apply (rotate-tac -1, erule ssubst)
apply (simp (no-asm-simp) only: f-tFold-const)
apply (cut-tac c=c7 and a=( $\lambda h. tFold\ (\lambda h0\ t0. w1\ (a1\ u\ (a2\ h\ (a3\ ha\ t0))$ ))) in tFold-LRunit)
apply (simp (no-asm-simp))
apply (simp (no-asm-simp) only: f-tFold-const)
apply (simp (no-asm-simp) only: tFold-const-const)
apply assumption
apply (rotate-tac -1, erule ssubst)
apply (subgoal-tac tFold ( $\lambda h0\ t0. w1\ (a1\ u\ (a2\ arbitrary\ (a3\ ha\ t0))$ )) const = tFold ( $\lambda h0\ t0. w1\ (a1\ u\ arbitrary)$ ) const)
apply (rotate-tac -1, erule ssubst)
apply (simp (no-asm-simp) only: tFold-const-const)
apply (rule-tac x=const in fun-cong)
apply (rule-tac f=tFold in arg-cong)
apply (rule ext)
apply (rule ext)
apply (rule-tac c=c7 and F=%x . w1 (a1 u x) in LRunit-const)
apply (simp (no-asm-simp))
apply (rotate-tac -1, erule ssubst)
apply (simp (no-asm-simp) del: slimH2-def tFold-tMap)
apply (thin-tac ALL t . tFold a5 c5 (tMap (tFold a6 c6) t) = tFold ( $\lambda h\ t0. a5\ h\ (tFold\ a6\ c6\ t0)$ ) c5 t)
apply (subgoal-tac tFold ( $\lambda h\ t0. a5\ h\ (tFold\ a6\ c6\ t0)$ ) c5 = ( $\% t. tFold\ a5\ c5$ ) (tMap (tFold a6 c6) t))
prefer 2
apply (rule ext)
apply (subst tFold-tMap)
apply assumption
apply (rule refl)
apply (simp (no-asm-simp) del: tFold-tMap add: LRunit-left LRunit-right)

```

Now we expand some extra effort to be able to work with a single, wrapped-only unit assumption.

apply (*subgoal-tac* $w1$ ($c1$ ($a1$ h ($tFold$ $a2$ $c2$ ($tMap$ ($tFold$ $a3$ $c3$ $t1$))))))
 $(c1$ ($a1$ u *arbitrary*))
 $(c1$ ($a1$ u *arbitrary*))
 $(a1$ h ($tFold$ $a2$ $c2$ ($tMap$ ($tFold$ $a3$ $c3$ $t2$)))))) =
 $c7$ ($w1$ ($a1$ h ($tFold$ $a2$ $c2$ ($tMap$ ($tFold$ $a3$ $c3$ $t1$))))))
 $(c7$ ($w1$ ($a1$ u *arbitrary*)))
 $(c7$ ($w1$ ($a1$ u *arbitrary*)))
 $(w1$ ($a1$ h ($tFold$ $a2$ $c2$ ($tMap$ ($tFold$ $a3$ $c3$ $t2$))))))
prefer 2
apply (*simp* (*no-asm-simp*))
apply (*rotate-tac* -1, *erule* *trans*)
apply (*subst* *LRunit-left* [of $c7$ $w1$ ($a1$ u *arbitrary*)])
apply (*simp* (*no-asm-simp*))
apply (*subst* *LRunit-left* [of $c7$ $w1$ ($a1$ u *arbitrary*)])
apply (*simp* (*no-asm-simp*))
apply (*simp* (*no-asm-simp*))
2: $t = hConc$ $t1$ $t2$
apply (*simp* *del*: *slimH2-def* *tFold-tMap*)
done

lemma *tFold-tFold-tFoldC-inverse2-gen*:

\llbracket *assoc* $c1$; *assoc* $c2$; *assoc* $c3$; *assoc* $c4$; *assoc* $c5$; *assoc* $c6$; *assoc* $c7$;
 \wedge h ha hb . $w2$ ($a4$ hb ($a5$ ha ($a6$ h))) = $w1$ ($a1$ h ($a2$ ha ($a3$ hb)));
(* *finishing* 2, and *influencing* the next rule: *)
 \wedge x y . $c7$ ($w1$ x) ($w1$ y) = $w1$ ($c1$ x y);
(* *finishing* 1.1.2 *)
 \wedge h ha x y .
 $w1$ ($c1$ ($a1$ h ($a2$ ha x)) ($a1$ h ($a2$ ha y))) =
 $w1$ ($a1$ h ($a2$ ha ($c3$ x y)));
(* *The next few are for* *tFold-tFold-vConc-wrapped*: *)
 \wedge h x y . $w2$ ($a4$ h ($c5$ x y)) = $w2$ ($c4$ ($a4$ h x) ($a4$ h y));
 \wedge x y . $w2$ ($c4$ x y) = $c7$ ($w2$ x) ($w2$ y);
 \wedge $x1$ $x2$ $y1$ $y2$. $c7$ ($w2$ ($c4$ $x1$ $x2$)) ($w2$ ($c4$ $y1$ $y2$)) =
 $w2$ ($c4$ ($c4$ $x1$ $y1$) ($c4$ $x2$ $y2$));
(* *stronger*:
 \wedge $x1$ $x2$ $y1$ $y2$. $c4$ ($c4$ $x1$ $x2$) ($c4$ $y1$ $y2$) =
 $c4$ ($c4$ $x1$ $y1$) ($c4$ $x2$ $y2$);
*)
(* *units*: *)
(*
 \wedge x y . *LRunit* $c5$ ($a5$ y ($a6$ u));
 \wedge h x y . *LRunit* $c4$ ($a4$ h ($a5$ y ($a6$ u)));
*)
 \wedge x . *LRunit* $c7$ ($w1$ ($a1$ u x));
(* *finishing* 1.2: *)
 \wedge h x y . $w1$ ($c1$ ($a1$ h x) ($a1$ h y)) = $w1$ ($a1$ h ($c2$ x y))
 $\rrbracket \implies$
 $w2$ ($tFold$ $a4$ $c4$ ($tMap$ ($tFold$ $a5$ $c5$ o ($tMap$ ($tFoldC$ $a6$ $c6$))) (*inverse2* u t))) =

```

    w1 (tFold a1 c1 (tMap (tFold a2 c2 o (tMap (tFoldC a3 c3))) t))
  apply (unfold tFoldC-def)
  apply (rule tFold-tFold-tFold-inverse2-gen)
  apply assumption+
  apply (simp-all (no-asm-simp))
  apply simp
  apply (subgoal-tac c7 (c7 (w2 x1) (w2 x2)) (c7 (w2 y1) (w2 y2)) =
    c7 (w2 (c4 x1 x2)) (w2 (c4 y1 y2)))
    prefer 2
    apply (simp (no-asm-simp))
  apply simp
done

```

lemma *id-tFold-tFold-tFoldC-inverse2-gen*:

```

[[ assoc c1; assoc c2; assoc c3; assoc c5; assoc c6;
  ∧ h ha hb . a4 hb (a5 ha (a6 h)) = a1 h (a2 ha (a3 hb));
  (* finishing 1.1.2 *)
  ∧ h ha x y .
    c1 (a1 h (a2 ha x)) (a1 h (a2 ha y)) = a1 h (a2 ha (c3 x y));
  (* The next few are for tFold-tFold-vConc-wrapped: *)
  ∧ h x y . a4 h (c5 x y) = c1 (a4 h x) (a4 h y);
  ∧ x1 x2 y1 y2 . c1 (c1 x1 x2) (c1 y1 y2) =
    c1 (c1 x1 y1) (c1 x2 y2);
  (* units: *)
  ∧ x . LRunit c1 (a1 u x);
  (* finishing 1.2: *)
  ∧ h x y . c1 (a1 h x) (a1 h y) = a1 h (c2 x y)
]] ⇒
id (tFold a4 c1 (tMap (tFold a5 c5 o (tMap (tFoldC a6 c6))) (inverse2 u t))) =
id (tFold a1 c1 (tMap (tFold a2 c2 o (tMap (tFoldC a3 c3))) t))
apply (rule tFold-tFold-tFoldC-inverse2-gen)
apply assumption+
apply (simp-all (no-asm-simp))
apply (subgoal-tac a1 h (c2 (a2 ha x) (a2 ha y)) = c1 (a1 h (a2 ha x)) (a1 h (a2
ha y)))
  prefer 2
  apply simp
apply (erule trans)
apply (rotate-tac -1, simp)
done

```

end